## HW6

3.15: The spanning trees are listed below.

3.16: The spanning tree in row $i$ and col $j$ above has the Prüfer code $i j$.
3.22: For each edge $e$ we divide the spanning trees of $G$ into two subsets, those which contain $e$ and those that do not. Clearly a spanning tree of $G$ which does not contain $e$ is also a spanning tree of $G-e$. Conversely, a spanning tree of $G-e$ is a spanning tree of $G$ which does not contain $e$. This is a 1-1 correspondence between spanning trees of $G$ not containing $e$ and spanning trees of $G-e$.

Suppose $T$ is a spanning tree of $G$ which does contain $e=u v$. The graph $T \circ e$ is still connected and has one less edge than vertex so it is a tree and it spans $G \circ e$. The problem is going backwards from a spanning tree of $G \circ e$ to a spanning tree of $G$ which uses $e$ when $G \circ e$ is a multi-graph and not a graph. Let $\bar{e}$ be the new vertex of $G \circ e$ and label each edge incident to $\bar{e}$ with either $u$ or $v$ depending on if it came from an edge incident to $u$ or $v$ respectively.

Suppose $S$ is a spanning tree of the multi-graph $G \circ e$. Construct $T$ from $S$ as follows. Any edge in $S$ not incident to $\bar{e}$ is also an edge of $T$. The edges $w \bar{e}$ in $S$ incident to $\bar{e}$ are labeled $u$ or $v$ which correspond to $w u$ or $w v$ respectively in $T$. Add the edge $e$ to $T$, so $T$ has the correct number of edges for a tree. Furthemore, this construction cannot create a cycle so $T$ is a spanning tree of $G$. This is a $1-1$ correspondence between spanning trees of $G$ containing $e$ and spanning trees of $G \circ e$.

