

Use the recursive definitions of ordinal arithmetic:

$$\alpha + 0 = \alpha, \quad \alpha \cdot 0 = 0 \quad \alpha^0 = 1$$

$$\alpha + (\beta + 1) = (\alpha + \beta) + 1 \quad \alpha \cdot (\beta + 1) = \alpha \cdot \beta + \alpha \quad \alpha^{\beta+1} = \alpha^\beta \cdot \alpha$$

And if  $\beta$  is a non-zero limit ordinal

$$\alpha + \beta = \cup\{\alpha + \gamma : \gamma < \beta\} \quad \alpha \cdot \beta = \cup\{\alpha \cdot \gamma : \gamma < \beta\} \quad \alpha^\beta = \cup\{\alpha^\gamma : \gamma < \beta\}$$

to show the following:

1. For all ordinals  $\alpha, \beta, \gamma, \delta$  :

$$0 + \alpha = \alpha, \text{ and } \omega \leq \alpha \Rightarrow 1 + \alpha = \alpha$$

$$0 < \beta \Rightarrow \alpha < \alpha + \beta$$

$$\alpha \leq \beta \ \& \ \gamma \leq \delta \Rightarrow \alpha + \gamma \leq \beta + \delta$$

$$\alpha \leq \beta \ \& \ \gamma < \delta \Rightarrow \alpha + \gamma < \beta + \delta$$

Show also, that, in general

$$\alpha < \beta \text{ does not imply } \alpha + \delta < \beta + \gamma$$

2. For all ordinals  $\alpha, \beta, \gamma, \delta$  :

$$0 \cdot \alpha = 0$$

$$0 < \alpha \ \& \ 1 < \beta \Rightarrow \alpha < \alpha \cdot \beta$$

$$\alpha \leq \beta \ \& \ \gamma \leq \delta \Rightarrow \alpha \cdot \gamma \leq \beta \cdot \delta$$

$$0 < \alpha \leq \beta \ \& \ \gamma < \delta \Rightarrow \alpha \cdot \gamma < \beta \cdot \delta$$

Show also that even when  $\gamma > 0$  in general,

$$\alpha < \beta \text{ does not imply } \alpha \cdot \gamma < \beta \cdot \gamma$$

3. (Cancellation laws) For all ordinals  $\alpha, \beta, \gamma$  :

$$\alpha + \beta < \alpha + \gamma \Rightarrow \beta < \gamma$$

$$\alpha + \beta = \alpha + \gamma \Rightarrow \beta = \gamma$$

$$\alpha \cdot \beta < \alpha \cdot \gamma \Rightarrow \beta < \gamma$$

$$0 < \alpha \ \& \ \alpha \cdot \beta = \alpha \cdot \gamma \Rightarrow \beta = \gamma$$

Show also that, in general

$$0 < \alpha \ \& \ \beta \cdot \alpha = \gamma \cdot \alpha \text{ does not imply } \beta = \gamma$$

4. For all  $\alpha \geq \omega$  and  $n < \omega$  :

$$n + \alpha = \alpha$$

$$(\alpha + 1) \cdot n = \alpha \cdot n + 1 \quad (n > 0)$$

$$(\alpha + 1) \cdot \omega = \alpha \cdot \omega$$

Give an example:

$$(\alpha + \beta) \cdot \gamma \neq \alpha \cdot \gamma + \beta \cdot \gamma$$

5. If  $\alpha \leq \gamma$  then there exists exactly one  $\beta$  such that  $\gamma = \alpha + \beta$ .  
 6. For all ordinals  $\alpha, \beta, \gamma$  with  $\alpha > 1$

$$\beta < \gamma \Rightarrow \alpha^\beta < \alpha^\gamma$$

$$\alpha^{(\beta+\gamma)} = \alpha^\beta \cdot \alpha^\gamma$$

$$(\alpha^\beta)^\gamma = \alpha^{\beta \cdot \gamma}$$

7. If  $\alpha > 0$ , then there is a largest  $\beta$  such that  $\omega^\beta \leq \alpha$  and for that  $\beta$ ,  $\alpha = \omega^\beta + \gamma$  for some  $\gamma < \alpha$ .  
 8. If  $\beta < \gamma$ , then  $\omega^\beta + \omega^\gamma = \omega^\gamma$ . Every ordinal  $\alpha > 0$  can be uniquely written as  $\alpha = \omega^\beta + \gamma$  with  $\gamma < \alpha$   
 9. Every ordinal  $\alpha > 0$  can be uniquely written

$$\alpha = \omega^{\beta_1} \cdot n_1 + \omega^{\beta_2} \cdot n_2 \cdots + \omega^{\beta_k} \cdot n_k$$

$$(\beta_1 > \beta_2 > \cdots > \beta_k, 0 < n_i < \omega)$$

This is called the Cantor normal form.

10. Find the Cantor normal form of

$$\omega \cdot (\omega^\omega + 1) + (\omega^\omega + 1) \cdot \omega$$