# Friend Networks on Convex Polynomials or 1, 2, 3, 4, ? This welcome is brought to you by Social Networks, Food Web Networks, and Convex Polyhedra Steven F. Bellenot Aug 24, 2012

The talk is a slide show. The slides are framed in yellow rectangles. The quotation that follows, is what might have been said while the audience was looking at the slide. The blue comments like this one were added later and not part of the welcome. The title frame above was not the original.



Why did I do social networks and food webs?

There was a talk on big networks the semester before.



Little Rock Lake is in Wisconsin. Note the loop, making the smallmouth bass a cannibal. My food web section in biocalculus, had no cannibals. An example where not all edges are friendly.

#### new stuff



A convex polyhedra is the convex hull of a finite set of points in  $\mathbb{R}^n$ . Network vertices are the extreme points (i.e. (2,2) and (0,4) but not (1,1) nor (3,1)) and Network edges are extreme line segments between vertices (i.e. [(2,2), (0,4)] but not [(0,0), (2,2)]). Two vertices are friends if they are joined by an edge.

A convex polyhedra is the convex hull of its extreme points. We make the extreme points vertices, and connect vertices that share an extrama edge. In two dimensions, this will be a convex polygon. In three dimensions a polyhedra. This picture tries to clarify.

Friends mean adjacency.

Maybe I should delete the first two examples an start with a docedahedran? A cube?

## Clicks

A *click* is a set of vertices, every pair of which are friends.



How are the two graphs for  $K_4$  the same? I think they are two attempts at getting two dimensional graphs with a click of four. Both of course fail. The first failure works as a click in three dimensions, the second failure doesn't work in three dimensions.





Suppose the convex hull of these points in the plane form a click.

Change the plot to larger solid dots. We want to show that the click has at most three vertices.



We start by moving a line from the far left (west) towards the east. We stop when we make contact with a vertex.

I could touch more than one vertex, if so, pick the topmost.



We are rotating the top ray and stop when we come to be in contact with a second point.

There cannot be two new points on the line. The middle one would not be an extreme point.



Repeat with the lower ray. We now have a angle of extreme points. All other extreme points most be inside the angle. So edge from our leftmost point, to any of the other points can be extremal. The click is at most three.





That is what I said.

new stuff



The points between the leftmost and the rightmost are not extreme points. The max click size in one dimension is 2.

new stuff



The tetrahedron the best you can do. This picture follows our two ideas. Move a plane from the far left. rotate it til it touchs a second point. Use the edge between them as a hinge to find the thrid point.

We can continue doing planes until we have a solid angle. This solid angle cannot have four sides because both of the red edges cannot be extrema.

new stuff

Maximal Convex Clie	cks Size	in Dimension	n
	Dimension	Max Convex Click	
	0	1	
	1	2	
	2	3	
	3	4	
	4	?	

Here is a table summarizing the results so far. We can move to four dimensions, what do we expect the answer to be?

new stuff

Maximal Convex Clicks Size in Dimension $n$					
	Dimension	Max Convex Click			
	0	1			
	1	2			

2

3

4

3

4

 $\infty$ 

The answer is infinity!

new stuff

### In 4-dimensions there is no max click size

Pick distinct  $t_i$ ,  $1 \le i \le n$ . The points  $P_i = (t_i, t_i^2, t_i^3, t_i^4)$ ,  $1 \le i \le n$  form a convex click.

$$f(t) = (t - t_3)^2 (t - t_5)^2 = a_0 + a_1 t + a_2 t^2 + a_3 t^3 + a_4 t^4$$

Note  $f(t_3) = f(t_5) = 0 < f(t_i)$  for  $3 \neq i \neq 5$ . Note  $P_i$  dot product  $[a_1, a_2, a_3, a_4] = f(t_i) - a_0$ We have found a linear functional that exposes the edge from  $P_3$  to  $P_5$ .

The points are on a curve called the moment curve. Any two on this curve are extreme points with an extrema edge connecting them. We illustrate this with the  $t_3$  and  $t_5$ , using the function f(t). The polynomial f(f) is always non-negative and strictly positive except for the zero's at  $t_3$  and  $t_5$ .

Looking at the coefficients of f(t) we can see that f is a linear functional, a dot product. The half space that f determines contains just the two points  $P_3$  and  $P_5$  (and the edge between them).

I learned this from a Claremont colleges colloquium. The term the speaker used was neighborly, in this case it was 2-neighborly. The result generalizes, in *d*-dimensions, this construction yields a  $k = \lfloor d/2 \rfloor$  neighborly set. Each k vertices are on an extrema half space.

### **Picture sources**

social I can't find again. Perhaps replace it with http://authenticorganizations.com/harquail/2011/ 01/24/your-authentic-social-network-the-identity-graph/#sthash.aALsBCch.dpbs

foodweb picture is from

https://www.researchgate.net/figure/3D-Visualization-of-the-food-web-of-Little-Rock-Lake-Wisconsin fig1\_221000766 N D Martinez, 1991 Artifacts of Attributes? Effects of Resolution on the Little Rock Lake Food Web, Ecological Monographs 61, 367-392.