

Discrete Morse Theory and Persistent Homology

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Overview

- 1 Discrete Morse Theory
 - Definitions
 - Gradients

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- 3 Persistence vs. DMT
 - Persistence to Discrete Vector Fields
 - Discrete Vector Fields to Persistence

Discrete Morse Theory

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$$f : M \rightarrow \mathbb{R}.$$

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It must satisfy the following two conditions, for every p -simplex $\alpha^{(p)}$ in M :

- 1 $\#\{\beta^{(p+1)} > \alpha^{(p)} \mid f(\beta) \leq f(\alpha)\} \leq 1$;
- 2 $\#\{\tau^{(p-1)} < \alpha^{(p)} \mid f(\tau) \geq f(\alpha)\} \leq 1$.

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Simple example: $f : M \rightarrow \mathbb{R}, f(\alpha) = \dim \alpha$

Critical Points

A simplex $\alpha^{(p)}$ is *critical* if the following two conditions hold:

- 1 $\#\{\beta^{(p+1)} > \alpha^{(p)} \mid f(\beta) \leq f(\alpha)\} = 0$;
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A simplex that is not critical is called *regular*.

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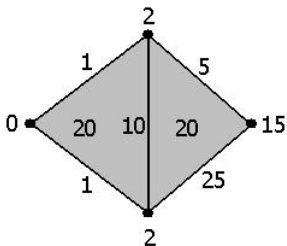
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critical simplex σ
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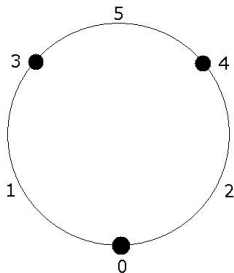
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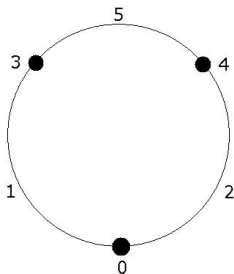
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- 1 Critical vertex = Local minimum
- 2 Critical n -cell = Local maximum
- 3 For $f(\alpha) = \dim \alpha$, every cell is critical.

Here is a discrete Morse function on the circle:

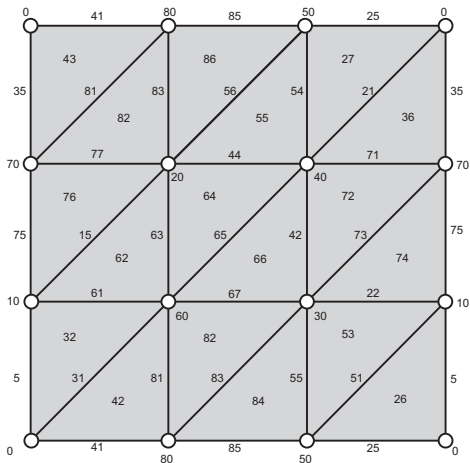


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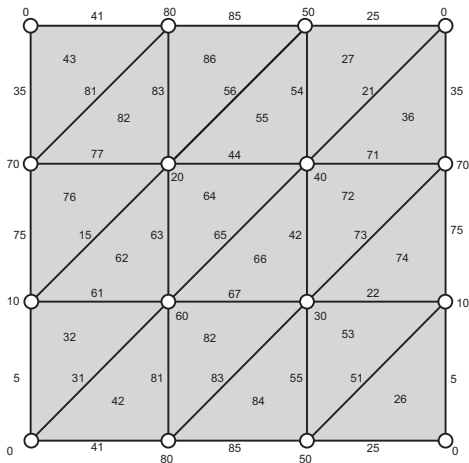


There are two critical cells, $f^{-1}(0)$ and $f^{-1}(5)$.

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The critical cells are $f^{-1}(0)$, $f^{-1}(42)$, $f^{-1}(44)$, and $f^{-1}(86)$.

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So, the torus has the homotopy type of a complex with one vertex, two 1-cells, and one 2-cell.

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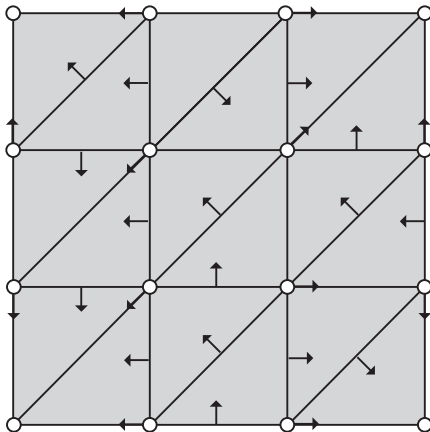
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In the last case, σ is critical.

Here is the gradient field on the torus associated to the above discrete Morse function:



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Definition. Let V be a discrete vector field on M . A V -*path* is a sequence of simplices

$$\alpha_0^{(p)}, \beta_0^{(p+1)}, \alpha_1^{(p)}, \beta_1^{(p+1)}, \dots, \beta_r^{(p+1)}, \alpha_{r+1}^{(p)}$$

with $\{\alpha_i < \beta_i\} \in V$, $i = 0, \dots, r$, and $\beta_i > \alpha_{i+1} \neq \alpha_i$. Such a path is called *closed* if $\alpha_{r+1} = \alpha_0$.

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Another way to look at this, due to Chari, is to consider the Hasse diagram of M , the directed graph whose vertices are the simplices of M with arrows $\beta^{(p+1)} \rightarrow \alpha^{(p)}$ for $\beta > \alpha$. Modify it by reversing an arrow whenever $\{\alpha, \beta\} \in V$. Then V is a gradient if and only if this modified diagram has no directed loops.

Persistent Homology

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If a class α is born in $H_i(M_j)$ and dies in $H_i(M_\ell)$, we say the *persistence* of α is $\ell - j - 1$.

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This creates a pairing of the simplices of M .

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Note that a simplex may get paired with something that is not adjacent to it. Also, each simplex in M is in at most one pair (some simplices do not get paired).

Here's one way to do this. Assume the filtration on M is such that each level M_i is obtained from M_{i-1} by attaching a single simplex σ_i .

Let D be the incidence matrix of the complex M :

$$D_{ij} = \begin{cases} 1 & \text{if } \sigma_i \in \partial\sigma_j \\ 0 & \text{otherwise} \end{cases}$$

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Definition. R is called *reduced* and low_R is a *pairing function* if $\text{low}_R(j) \neq \text{low}_R(j')$ for nonzero columns $j \neq j'$.

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Reduce D by adding columns mod 2 to other columns located to the right:

- $R = D$
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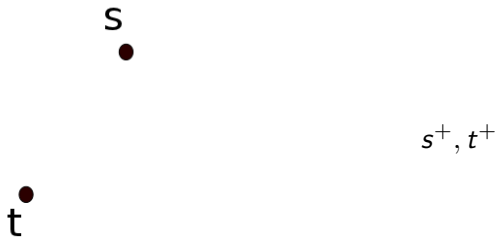
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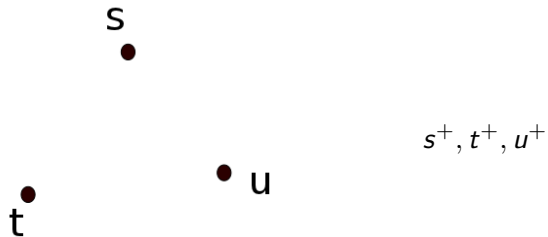
s


s^+

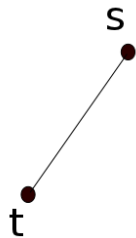
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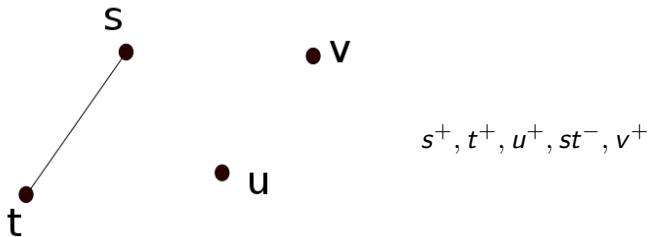


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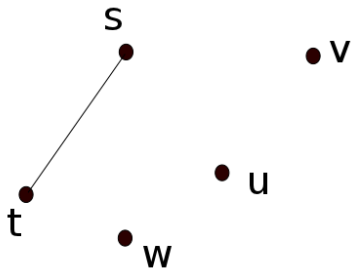


s^+, t^+, u^+, st^-

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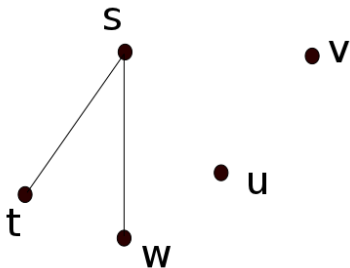


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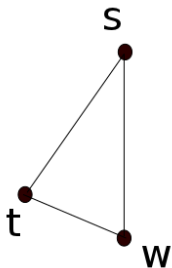
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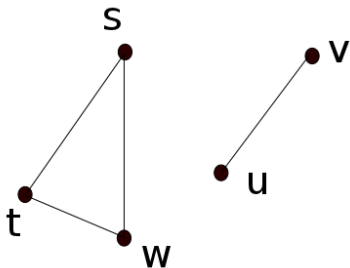
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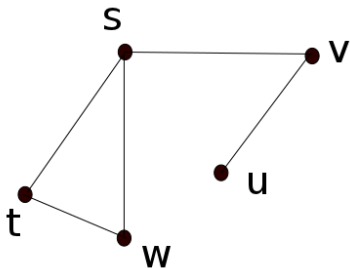
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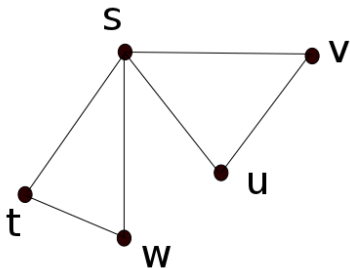
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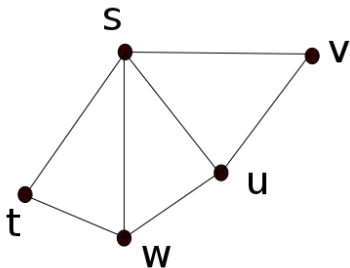
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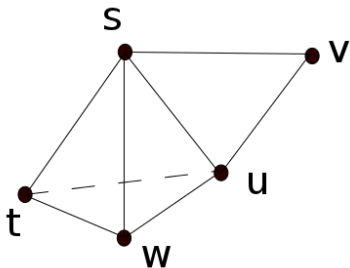
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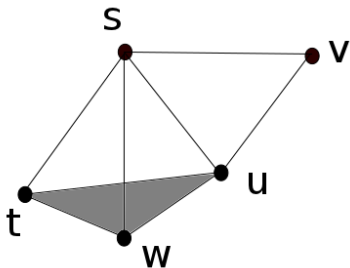
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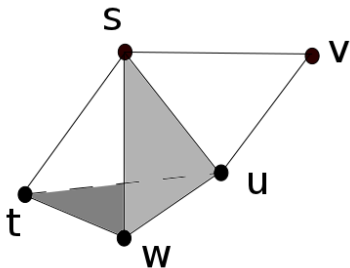
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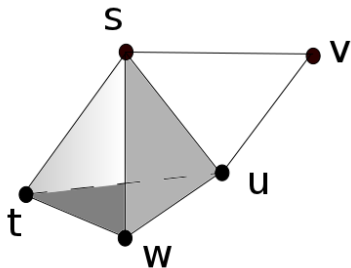
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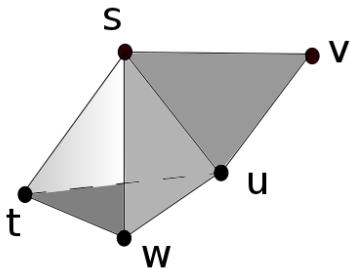
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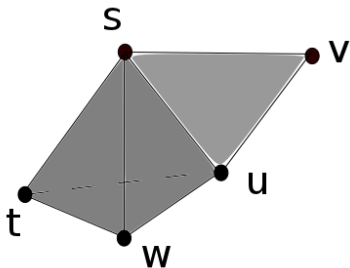
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The resulting pairing arising from this filtration is then

$$\begin{array}{ll} (t, st) & (tu, tuw) \\ (w, sw) & (uw, suw) \\ (v, uv) & (tw, stu) \\ (u, sv) & (su, suv) \end{array}$$

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Note that we wound up with (u, sv) and (tw, stu) . In both cases, the lower dimensional simplex is not a face of the larger one. What that means is that it took a while for the homology classes represented by u and tw to die. This is information you might want to know.

BUT

BUT This looks like a vector field on M if we ignore these anomalies.

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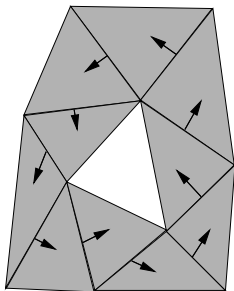
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Note that each pair $\{\alpha, \beta\} \in V_P$ has α positive and β negative.

Theorem. V_P is a gradient.

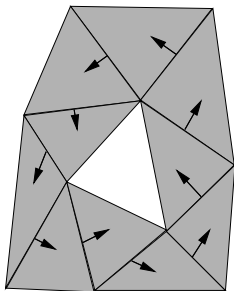
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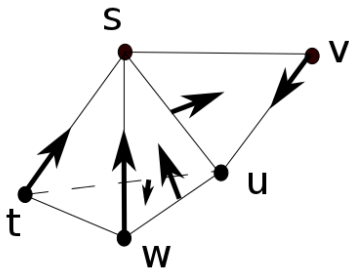
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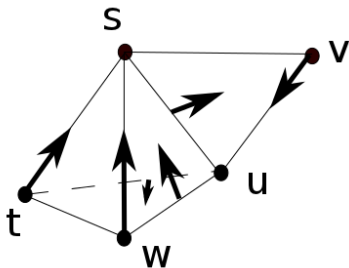


Each of the edges on the interior is positive by assumption. But that can't be—one of those edges had to join the outer loop with the inner triangle and therefore be negative. (N.B. This is a gross oversimplification, but it illustrates the point.)

An Example

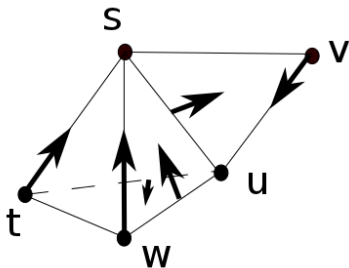


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But we have unique gradient paths

$$sv > v < uv > u$$

and

$$stu > tu < tuw > tw$$

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Conjecture. Those simplices that are paired by P but not in V_P have a unique gradient path joining them.

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The natural thing to try is to filter M by sublevel sets of a discrete Morse function associated to V (there's a fairly canonical choice of such a map). This will lead to a filtration where more than one simplex enters at a time, but that can be dealt with. I'm still working on this.

What You Want:

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What You Can't Have:

What You Want: $P_{V_P} = P$

What You Can't Have: $V_{P_V} = V$

What You Want: $P_{V_P} = P$

What You Can't Have: $V_{P_V} = V$

Example: Take $V = \emptyset$. Then $P_V \neq \emptyset$; in fact some vertex will get paired with one of its edges, and this pair will belong to V_{P_V} .

Future Work

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- 3 Kozlov has characterized discrete Morse functions as poset maps with “small fibers” (fibers of cardinality ≤ 2). Since persistent homology roughly corresponds to discrete Morse theory, what is the theory associated with other poset maps? Is there a connection to multidimensional persistence?

Fin