

# Study Questions Homework 3 Foundations of Computational Math 1 Fall 2021

## Problem 3.1

Let  $x$  and  $y$  be two vectors in  $\mathbb{R}^n$ .

**3.1.a.** Show that given  $x$  and  $y$  the value of  $\|x - \alpha y\|_2$  is minimized when

$$\alpha_{min} = \frac{x^T y}{y^T y}$$

**3.1.b.** Show that  $x = y\alpha_{min} + z$  where  $y^T z = 0$ , i.e.,  $x$  is easily written as the sum of two orthogonal vectors with specified minimization properties.

## Problem 3.2

Recall that an elementary reflector has the form  $Q = I + \alpha z z^T \in \mathbb{R}^{n \times n}$  with  $\|z\|_2 \neq 0$ .

**3.2.a.** Show that  $Q$  is orthogonal if and only if

$$\alpha = \frac{-2}{z^T z} \text{ or } \alpha = 0$$

**3.2.b.** Given  $v \in \mathbb{R}^n$ , let  $\gamma = \pm\|v\|$  and  $z = v + \gamma e_1$ . Assuming that  $z \neq v$  show that

$$\frac{z^T z}{z^T v} = 2$$

**3.2.c.** Using the definitions and results above show that  $Qv = -\gamma e_1$

## Problem 3.3

Let  $x \in \mathbb{R}^n$  be a known vector with components  $\xi_i = e_i^T x$ ,  $1 \leq i \leq n$  and consider the computation of

$$\nu = \xi_1 - \|x\|_2$$

where  $\|x\|_2^2 = \sum_{i=1}^n \xi_i^2$ . (Recall this is a key computation in the production of a Householder reflector in least squares problems.) When  $\xi_1 > 0$  and  $\xi_1 \approx \|x\|_2$  the cancellation in the subtraction may result in a significant loss of accuracy.

Find an alternate expression for  $\nu$  that does not suffer from cancellation when  $\xi_1 > 0$  and  $\xi_1 \approx \|x\|_2$ . (**Hint:** Consider a difference of squares.)

## Problem 3.4

Consider a Householder reflector,  $H$ , in  $\mathbb{R}^2$ . Show that

$$H = \begin{pmatrix} -\cos(\phi) & -\sin(\phi) \\ -\sin(\phi) & \cos(\phi) \end{pmatrix}$$

where  $\phi$  is some angle.

## Problem 3.5

Suppose you are given the nonsingular tridiagonal matrix  $T \in \mathbb{R}^{n \times n}$ . For example, if  $n = 6$  then

$$\begin{pmatrix} \alpha_1 & \beta_1 & 0 & 0 & 0 & 0 \\ \gamma_2 & \alpha_2 & \beta_2 & 0 & 0 & 0 \\ 0 & \gamma_3 & \alpha_3 & \beta_3 & 0 & 0 \\ 0 & 0 & \gamma_4 & \alpha_4 & \beta_4 & 0 \\ 0 & 0 & 0 & \gamma_5 & \alpha_5 & \beta_5 \\ 0 & 0 & 0 & 0 & \gamma_6 & \alpha_6 \end{pmatrix}$$

**3.5.a** Suppose you use Householder reflectors to transform  $T$  to upper triangular, i.e.,

$$H_{n-1} \dots H_1 T = R.$$

What is the zero/nonzero structure of  $R$ ?

**3.5.b** What is the structure of each of the reflectors  $H_i$ ?

**3.5.c** Let  $T^{(i)} = H_i H_{i-1} \dots H_1 T$ . What is the structure of  $T^{(i)}$ ?

**3.5.d** What is the computational complexity of the factorization, i.e., what is  $k$  in  $O(n^k)$ ? (You do not have to determine the constant in the complexity expression.)

To answer the questions on structure, in addition to characterizing it algebraically, include a  $*$ , 0 matrix pattern for  $n = 6$  to make it clear.

## Problem 3.6

### 3.6.a

This part of the problem concerns the computational complexity question of operation count.

For both  $LU$  factorization and Householder reflector-based orthogonal factorization, we have used elementary transformations,  $T_i$ , that can be characterized as rank-1 updates to the identity matrix, i.e.,

$$T_i = I + x_i y_i^T, \quad x_i \in \mathbb{R}^n \text{ and } y_i \in \mathbb{R}^n$$

Gauss transforms and Householder reflectors differ in the definitions of the vectors  $x_i$  and  $y_i$ . Maintaining computational efficiency in terms of a reasonable operation count usually implies careful application of associativity and distribution when combining matrices and vectors.

Suppose we are to evaluate

$$z = T_3 T_2 T_1 v = (I + x_3 y_3^T)(I + x_2 y_2^T)(I + x_1 y_1^T)v$$

where  $v \in \mathbb{R}^n$  and  $z \in \mathbb{R}^n$ . Show that by using the properties of matrix-matrix multiplication and matrix-vector multiplication, the vector  $z$  can be evaluated in  $O(n)$  computations (a good choice of version for an algorithm) or  $O(n^2)$  computations (a bad choice of version for an algorithm) or  $O(n^3)$  computations (a very bad choice of version for an algorithm).

### 3.6.b

This part of the problem concerns the computational complexity question of storage space.

Recall, that we discussed and programmed an **in-place** implementation of  $LU$  factorization that was very efficient in terms of storage space. An array with  $n^2$  entries initialized with  $array(I, J) = \alpha_{ij}$  could be used to store the  $n^2$  entries needed to specify  $L$  and  $U$ , i.e.,  $\lambda_{ij}$  for  $j < i$ ,  $2 \leq i \leq n$  and  $1 \leq j \leq n - 1$ , and  $\mu_{ij}$  for  $i < j$ ,  $1 \leq i \leq n$  and  $1 \leq j \leq n$ .

Let  $A \in \mathbb{R}^{n \times k}$ ,  $n \geq k$ , and  $rank(A) = k$ . Consider the use of Householder reflectors,  $H_i$ ,  $1 \leq i \leq k$ , to transform  $A$  to upper trapezoidal form, i.e.,

$$H_k H_{k-1} \cdots H_2 H_1 A = \begin{pmatrix} R \\ 0 \end{pmatrix}$$

$R \in \mathbb{R}^{k \times k}$  nonsingular upper triangular

Suppose you are given an array with  $n \times k$  entries initialized with  $array(I, J) = \alpha_{ij}$  and you are to implement your algorithm using minimal storage.

- (i) Are you able to store all of the information needed to specify the  $H_i$ ,  $1 \leq i \leq k$  and  $R$  within the array with  $n \times k$  entries? Justify your answer.
- (ii) If you are not able to store all of the information in the array, how much extra storage do you need and what do you store in it?

## Problem 3.7

### 3.7.a

Consider a set  $\{x_i\}$ ,  $i = 1, \dots, n$ , of vectors  $x_i \in \mathbb{R}^d$  and the vector 2-norm defined by the standard Euclidean inner product, i.e.,  $\langle x, y \rangle = x^T y$  and  $\|x\|_2^2 = \langle x, x \rangle$ . The Karcher mean or intrinsic mean of the vectors with respect to this norm is the unique minimizer of

$$f(x) = \min_{x \in \mathbb{R}^d} \sum_{i=1}^n \|x - x_i\|_2^2$$

and has a well-known closed form given by

$$\bar{x} = \frac{1}{n} \sum_{i=1}^n x_i.$$

Now consider the inner product  $\langle x, y \rangle_M = x^T M y$  and the norm it defines  $\|x\|_M^2 = \langle x, x \rangle_M$ , where  $M$  is a symmetric positive definite matrix. The Karcher mean or intrinsic mean,  $\tilde{x}$ , of the vectors with respect to this  $M$ -norm is the unique minimizer of

$$f_M(x) = \min_{x \in \mathbb{R}^d} \sum_{i=1}^n \|x - x_i\|_M^2$$

Determine a closed form expression for  $\tilde{x}$  and the effect of  $M$ , if any, on the mean. Justify your answer.

### 3.7.b

Suppose  $b \in \mathbb{R}^d$  is a known vector. Consider the unit sphere defined by the  $M$ -norm

$$\mathcal{S} = \{x \in \mathbb{R}^d \mid \|x\|_M^2 = 1\}.$$

Determine the approximating vector  $\hat{x}$  that solves

$$g(x) = \min_{x \in \mathcal{S}} \|x - b\|_M^2.$$

Justify your answer.