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1. Determine the number of digits of the following numbers.

- (a) 1000!
- (b)  $2^{10000}$

2. Determine the floating point representation of  $e = \exp(1)$  up to 10, 50 and 500 digits.

3. Determine the following numbers exact and up to 10 digits in their floating point representation.

- (a)  $\cos(\pi/6)$
- (b)  $\ln(2e)$
- (c)  $\arctan(\sqrt{3}+2)$

4. Determine the factorization of the following numbers,

- (a) 98765432101234567890
- (b)  $(2^{10})^{(2^{10})}$

5. Let

$$h(x) = \frac{x^5 - 3x^4 - 4x^3 - 11x^2 + 6x - 11}{x^5 - 5x^4 + 4x^3 - x^2 + 5x - 4}$$

- (a) Let f(x) be the numerator of h(x).
- (b) Let q(x) be the denomerator of h(x).
- (c) Let d(x) be the gcd(f(x), g(x)).
- (d) Determine f(x)/d(x) and g(x)/d(x).
- (e) Type normal(h). What happens?
- (f) Type convert(h, 'parfrac', x). What happens?

6. (a) Let  $f(x) = e^{-x^2}$ . Determine f'(x).

- (b) Let  $f(x) = \ln(x)$ . Determine F(x) such that F'(x) = f(x).
- (c) Let h(x) be as in exercise 5. Determine h'(x).
- (d) Determine H(x) such that H'(x) = h(x).

7. (a) Determine the taylor series in x = 0 of sin(x) up to order 10.

- (b) Determine the taylor series in x = 0 of  $e^x$  up to order 20
- (c) Let

$$f(x) = \frac{1}{1 - x - x^2}.$$

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Determine the taylor series in x = 0 of f(x) up to order 20. What can you say about the coefficients of this taylor series?

- 8. Solve the following (system of) equation(s).
  - (a)  $ax^2 + bx + c = 0$ .
  - (b)  $x^8 + x^3 x^2 1$ .

(c) 
$$\sin(x) = \cos(x)$$
.  
(d) 
$$\begin{cases} x + y = a \\ cx - 3y = b \end{cases}$$
(e) 
$$\begin{cases} x^2 + 1 = y \\ y - 3 = x \end{cases}$$
(f) 
$$\begin{cases} x^2 + y^2 - 5 = 0 \\ xy - y^2 + 2 = 0 \end{cases}$$
(g)  $x^2 + x + 1 = y^2$ ,  $x, y \in \mathbb{Z}$   
(h)  $x^2 + x + 1 = y^2$ ,  $x, y \in \mathbb{Z}/2\mathbb{Z}$  and  $x, y \in \mathbb{Z}/11\mathbb{Z}$   
(i) Solve  $x^8 + x^3 - x^2 - 1$  numerical.

- 9. Let

$$A = \begin{pmatrix} 1 & 0 & 1 & 0 \\ 1 & 1 & 1 & 1 \\ 0 & 1 & 1 & 1 \\ 0 & 0 & 0 & 1 \end{pmatrix}, B = \begin{pmatrix} 1 & 0 & 1 & 0 \\ 1 & 1 & 1 & 1 \\ 0 & 1 & 1 & 1 \\ 1 & 0 & 0 & 1 \end{pmatrix}.$$

- (a) Determine det(A) and det(B).
- (b) Determine  $A^{-1}$  and  $B^{-1}$ .
- (c) Determine C = AB.
- 10. Let

$$A = \left(\begin{array}{cccc} x & 2 & 3 & 4 \\ 5 & 6 & 7 & 8 \\ 9 & 10 & 11 & 12 \\ 13 & 14 & 15 & y \end{array}\right)$$

For which  $x, y \in \mathbb{R}$  holds:  $\det(A) = 0$ ?

## 11. Let

$$A = \left(\begin{array}{cccc} -12 & 12 & 4 & 0 \\ 4 & 4 & 4 & 4 \\ -40 & 4 & 4 & 4 \\ 21 & 0 & 0 & 4 \end{array}\right)$$

- (a) Determine  $A^4$ .
- (b) Determine the characteristic polynomial of A. Can you explain the relation between this result and the result in (a)?

12. Let

$$A = \left(\begin{array}{cccc} 1 & x_1 & x_1^2 & x_1^3 \\ 1 & x_2 & x_2^2 & x_2^3 \\ 1 & x_3 & x_3^2 & x_3^3 \\ 1 & x_4 & x_4^2 & x_4^3 \end{array}\right)$$

- (a) Determine det(a).
- (b) Factorize det(a).
- (c) Generalize this exercise to a  $6 \times 6$  matrix

13. Plot the following surfaces

Plot the following surfaces 
$$\begin{cases} x = \sin(s)\cos(t) & s \in [0, 2\pi] \\ y = \cos(s)\cos(t) & t \in [0, 2\pi] \\ z = \sin(t) & s \in [0, 2\pi] \end{cases}$$
(a) 
$$\begin{cases} x = \sin(s)\cosh(t) & s \in [0, 2\pi] \\ y = \cos(s)\cosh(t) & t \in [-2, 2] \\ z = \sinh(t) & t \in [-2, 2] \end{cases}$$
(b) 
$$\begin{cases} x = \sin(s)t & s \in [0, 2\pi] \\ y = \cos(s)t & t \in [-4, 4] \\ z = t & s \in [0, 2\pi] \end{cases}$$
(c) 
$$\begin{cases} x = \sin(s)(t) & s \in [0, 2\pi] \\ y = \cos(s)(t) & t \in [0, 2\pi] \\ z = t & t \in [0, 2\pi] \end{cases}$$
(d) 
$$\begin{cases} x = \sin(s)(2 + \cos(t)) & s \in [0, 2\pi] \\ y = \cos(s)(2 + \cos(t)) & t \in [0, 2\pi] \\ z = \sin(t) & t \in [-1..1] \end{cases}$$
(e) 
$$\begin{cases} x = \sin(s)(4 + t\cos(\frac{1}{2}s)) & s \in [0, 2\pi] \\ y = \cos(s)(4 + t\cos(\frac{1}{2}s)) & t \in [-1..1] \end{cases}$$

14. Solve, using Maple, for all  $a \in \mathbb{R}$  the following system of linear equations:

$$\begin{cases} 2ax + 6y = 5 \\ 5x + (a-2)y = 7 \end{cases}$$

Is the solution always correct?

15. Compute the extremal values of the function

$$f: x \mapsto x^5/5 - x^4/2 + x^2 - x$$
.

16. Let

$$A = \begin{pmatrix} 1 & 0 & 2 \\ 2 & 5 & 3 \\ 4 & 2 & 6 \end{pmatrix} \text{ and } B = \begin{pmatrix} 3 & 5 \\ 0 & 1 \\ 7 & 4 \end{pmatrix}.$$

Compute  $A^{-1}$ , AB and  $(A^2 + A)B$  using Maple.

17. Use Maple to verify the formula

$$4\arctan(1/5) - \arctan(1/239) = \pi/4$$
.

This formula can be used to find approximations of  $\pi/4$ . How?

- 4
- 18. Experiment with Maple and its on-line help system.
- 19. Let  $p \in \mathbb{N}$  be a prime number.
  - (a) Write a Maple-routine *multitable* which returns the multiplication table in  $\mathbb{Z}_p$  as in the following example for p = 5.

$$\left(\begin{array}{cccc} 1 & 2 & 3 & 4 \\ 2 & 4 & 1 & 3 \\ 3 & 1 & 4 & 2 \\ 4 & 3 & 2 & 1 \end{array}\right),$$

where the  $ij^{th}$  element is equal to  $ij \mod p$ .

(b) Write a Maple-routine divtable which returns the table of division in  $\mathbb{Z}_p$  as in the folling example for p = 5.

$$\left(\begin{array}{cccc} 1 & 2 & 3 & 4 \\ 3 & 1 & 4 & 2 \\ 2 & 4 & 1 & 3 \\ 4 & 3 & 2 & 1 \end{array}\right),$$

where the  $ij^{th}$  element is equal to  $ij^{-1} \mod p$ .

20. For the Fibonacci numbers  $F_n$  we have the following formula.

$$F_n = rac{1}{\sqrt{5}} \left( \left( rac{1+\sqrt{5}}{2} 
ight)^{n+1} - \left( rac{1-\sqrt{5}}{2} 
ight)^{n+1} 
ight).$$

- (a) Write a Maple-routine to compute the Fibonacci numbers using their recursive definition (don't forget the remember option).
- (b) Write a Maple-routine to compute the Fibonacci numbers using above-mentioned formula.
- (c) Compare the speed of these routines.
- 21. Write a routine in Maple which computes the Fibonacci-polynomials  $F_n(x)$ . The Fibonacci-polynomials are defined as follows:

$$\begin{cases} F_0(x) = 1 \\ F_1(x) = x \\ F_n(x) = xF_{n-1}(x) + F_{n-2}(x) & \text{voor } n \ge 2 \end{cases}$$

Compute  $F_{50}(y)$ .

- 22. Write a Maple-routine *elsym* such that  $\operatorname{elsym}(i, [x_1, x_2, \dots, x_n])$  returns the *i*th elementary symmetric polynomial in  $x_1, x_2, \dots, x_n$ .
- 23. Write a package to do integer arithmetic. Do this as follows:
  - Choose a base B (B a global variable).
  - Represent a non-negative integer by a list  $[a_{n-1}, \ldots, a_1, a_0]$  where  $0 \le a_i < B$  and  $a_{n-1} \ne 0$  or by [0]. A list  $[a_{n-1}, \ldots, a_1, a_0]$  will represent the number  $a_0 + a_1B + \cdots + a_{n-1}B^{n-1}$  and [0] will represent 0.

- Represent an integer by a list [s,p] where  $s \in \{-1,0,1\}$  and p the representation of a non-negative integer. [0,[0]] will represent  $0, [1,[a_{n-1},\ldots,a_0]]$  will represent  $a_0+a_1B+\cdots+a_{n-1}B^{n-1}$  and  $[-1,[a_{n-1},\ldots,a_0]]$  will represent  $-(a_0+a_1B+\cdots+a_{n-1}B^{n-1})$ .
- Write Maple-routines for:
  - comparing integers
  - adding integers
  - subtracting integers
  - multiplying integers

Be sure that all operations you perform are elementary. Test your code on examples using B=1000.

- 24. Extend the package from the previous exercise with a routine to compute the quotient and remainder when dividing two base B numbers.
- 25. (a) Write a Maple-routine representation such that for positive integers a and B the call representation (a, B) returns the base B representation of a, i.e. if  $a = a_0 + a_1 B + \cdots + a_n B^n$  it returns

$$[a_n, a_{n-1}, \ldots, a_1 a_0, B].$$

For example representation (1000000, 121) returns [68, 36, 56, 121].

- (b) Write a maple-routine *len* which returns the number of digits in the above representation, i.e. len(1000000, 121) = 9.
- (c) Write a Maple-routine  $minimal\_base$  which, for a positive integer a, computes a base B for which len(a, B) is minimal.
- (d) Can len(a, B) be smaller than length(a)? If yes, give an example.
- 26. The balanced ternary notation.

Prove that each non-zero integer can be written in a unique way as

$$d_0 + d_1 3 + d_2 3^2 + \dots + d_n 3^n$$
,

where  $n \in \mathbb{N}$ ,  $d_i \in \{-1,0,1\}$  and  $d_n \neq 0$ . This is called the *balanced ternary* notation. How can you see whether an integer in balanced ternary notation is positive or negative?

27. The mixed radix notation.

Let  $m_0, m_1, \ldots, m_n$  be integers > 1. Prove that each  $a \in \{0, 1, \ldots, m_0 m_1 \cdots m_n - 1\}$  can be written in a unique way as

$$d_0 + d_1 m_0 + d_2 m_0 m_1 + \cdots + d_n m_0 m_1 \cdots m_{n-1}$$

where  $d_i \in \{0, 1, \dots, m_i - 1\}$ . This is called the *mixed radix* notation.

28. Let B be a positive integer > 1 and m an integer satisfying 0 < m < B. Prove that  $|B/(m+1)|m \ge |B/2|$ .

- 6
  - 29. (a) Write a Maple-routine to compute  $a^n$  for positive integers a and n. Do this in two ways: by successive multiplication by a and by a divide-and-conquer technique. Compare the speed of both routines.
    - (b) Do the same as in (a) but now to compute  $a^n$  for a positive integer n and  $a \in \mathbb{Z}_m$  (m a positive integer).
    - (c) For a positive integer n let  $l_n$  be the minimal number of multiplications needed to compute  $a^n$ . For example the minimal number of multiplications needed to compute  $a^15$  is 5 (see lecture notes). Write a Maple-routine which computes for a positive integer m the set  $\{n \in \mathbb{N} \mid l_n = m\}$ .
  - 30. Notation as in Lemma 1 of *Elementary arithmetic of the integers*. Prove that:

$$1 = s_2 \le -s_3 < s_4 < -s_5 < \cdots$$

$$1 = t_1 \le -t_2 < t_3 < -t_4 < \cdots$$

$$s_{i-1}u_i - s_iu_{i-1} = (-1)^{i+1}b \qquad i = 1, \dots, n+1$$

$$t_{i-1}u_i - t_iu_{i-1} = (-1)^i a \qquad i = 1, \dots, n+1.$$

- 31. Prove that for  $a, b \in \mathbb{Z}$  such that  $a \geq b > 0$  and not both 1, there exist integers s and t such that gcd(a, b) = sa + tb and |s| < b, |t| < b. Do s and t computed in the extended Euclid's algorithm satisfy these inequalities?
- 32. Prove Corollary 4 of Elementary arithmetic of the integers. Use the fact that  $F_n = \left( (\frac{\sqrt{5}+1}{2})^n (\frac{\sqrt{5}-1}{2})^n \right) / \sqrt{5}$ .
- 33. Show that for positive integers  $a \ge b$  the time for computing q and r such that a = qb + r (long division) can be bounded by  $O(\ln(a)(\ln(q) + 1))$ . Use this to show that the computation time of the extended Euclid's algorithm can be bounded by  $O(\ln^2(a))$ .
- 34. Do exercices 10, 11 and 12 from Chapter 2 of [1].
- 35. Write a Maple-routine Karatsuba which multiplies two n-digit base B numbers using Karatsuba's method. Use the package of exercise 23 to do your base B arithmetic. Compare the speed of your routine with multiplication in Maple (take B=10). How does the speed of your routine depend on the size of the input?
- 36. Write a Maple-routine  $my\_gcd$  which for  $a, b \in \mathbb{Z}$  computes the greatest common divisor of a and b. You may use the Maple-routines irem and iquo for this (do you need both?). Write your routine such that it will print output as in the following example.

$$30 = 1 * 18 + 12$$
  
 $18 = 1 * 12 + 6$   
 $12 = 2 * 6 + 0$ 

37. Write a Maple-routine  $my\_gcdex$  which for  $a, b \in \mathbb{Z}$  retruns a list [d, s, t] such that d is the greatest common divisor of a and b and d = sa + tb. You may use the Maple-routines

*irem* and *iquo* for this (do you need both?). Write your routine such that it will print output as in the following example.

$$30 = 1 * 18 + 12$$
  $12 = 1 * 30 + -1 * 18$   
 $18 = 1 * 12 + 6$   $6 = -1 * 30 + 2 * 18$   
 $12 = 2 * 6 + 0$   $0 = 3 * 30 + -5 * 18$ 

38. The Maple-routines *irem* and *iquo* work do the following:

For  $a, b \in \mathbb{Z}$  we have a = iquo(a, b) \* b + irem(a, b) where -b < irem(a, b) < b and sgn(irem(a, b)) = sgn(a) (when  $\text{irem}(a, b) \neq 0$ ).

Write Maple-routines  $my\_irem$  and  $my\_iquo$  such that for  $a, b \in \mathbb{Z}$  the following holds:  $a = my\_iquo(a, b) * b + my\_irem(a, b)$  where  $-b/2 < my\_irem(a, b) < b/2$ .

Rewrite your routine from the previous exercise but now using my\_irem and my\_iquo instead of irem and iquo.

Compare both routines on some large inputs. Which is faster?

- 39. For randomly chosen integers a and b the probability that they are relatively prime is  $6/\pi^2$  (E. Cesaro, 1881). Check this theorem using Maple. For random number generation you can use teh Maple-routine rand.
- 40. Greatest common divisor in the ring of Gaussian integers.

Let  $G = \{n + mi \mid n, m \in \mathbb{Z}\}$   $(i = \sqrt{-1})$ . Show that G is a ring under usual addition and multiplication. G is called the ring of Gaussian integers. Define the norm N on G by  $N(n + mi) = n^2 + m^2$ . Claim: for  $a, b \in G$ ,  $b \neq 0$ , there are  $q, r \in G$  such that a = qb + r and  $0 \leq N(r) < N(b)$ . This can be seen as follows. Let c = a/b,  $c_1 = \Re(c), c_2 = \Im(c), q_1 = \operatorname{round}(c_1), q_2 = \operatorname{round}(c_2)$  and  $q = q_1 + q_2i$ . Here for  $a \in \mathbb{R}$ , round(a) = n is the integer satisfying  $n \leq a < n + 1/2$  or  $n - 1/2 \leq a \leq n$ . Now prove that  $N(q - c) \leq 1/2$  and  $N(a - bq) \leq N(b)/2$ .

Write a Maple-routine  $Gauss\_div$  which for  $a, b \in G$   $(b \neq 0)$  returns a list [q, r] such that  $q, r \in G$ , a = qb + r and  $0 \leq N(r) < N(b)$  (round is a builtin Maple-routine).

Write a Maple-routine  $Gauss\_gcd$  which for  $a, b \in G$  (not both 0) returns a list [d, s, t] such that d is the greatest common divisor of a and b satisfying  $\Re(d) > 0$ ,  $\Im(d) \ge 0$  and d = sa + tb.

Test your routines for several (large) a and b.

- 41. Write a Maple-routine gcd2 to compute a greatest common divisor of two numbers  $a, b \in \{n + m\sqrt{2} \mid n, m \in \mathbb{Z}\}.$
- 42. For randomly chosen integers a and b the probability that they are relatively prime is  $6/\pi^2$  (E. Cesaro, 1881). Check this theorem using Maple. For random number generation you can use teh Maple-routine rand.

43. For integers a > b > 0 we can write  $\frac{a}{b}$  as

where  $a_0, a_1, \ldots, a_n$  are positive integers. This expression is denoted by  $[a_0, a_1, \ldots, a_n]$ and is called a continued fraction.

- (a) Show how the numbers  $a_0, a_1, \ldots, a_n$  can be extracted from the results during the execution of Euclid's algorithm.
- (b) Write a Maple-routine calc\_frac which computes the rational number corresponding to a continued fraction.
- (c) Write a Maple-routine determine\_frac which computes the continued fraction corresponding to a rational number.
- (d) Test whether your routines are each others inverse.
- 44. Write a Maple-routine for multiplying polynomials using Karatsuba's method.
- 45. Experiment using Maple's normal function and study Maple's on-line help on this function. What normal form is used by Maple?
- 46. Simplify the following expressions using Maple (think of expand, normal, simplify, factor, combine, radsimp, etc.).
  - (a)  $(\exp(x) + x)/(\exp(2x) + 2x \exp(x) + x^2)$
  - (b)  $(x^5 + 40x^4 + 595x^3 + 3905x^2 + 9680x + 1331)^{1/3}$
  - (c)  $(x-2)^{3/2}/(x^2-4x+4)^{1/4}$
  - (d)  $(\sqrt{(x)} y)/(x y^2)$
  - (e)  $1/(2+5^{1/3})$
  - (f)  $\cos(x+y) + \sin(x)\sin(y) + 2^{x+y}$ (g)  $2\cos(x)^2 \cos(2x)$
- 47. Write a Maple-routine recursive which returns the recursive representation of a multivariate polynomial. For example

recursive
$$(x^3y^2z^2 + x^3y^2 + 3x^3y + zyx^2 - xy^4 + xzy^2 + yz - y + 3z, [x, y, z])$$

should return

$$((z^2+1)y^2+3y)x^3+zyx^2-(y^4-zy^2)x+(z-1)y+3z.$$

48. Write a Maple-routine horner which returns the representation of a polynomial according to Horner's evaluation rule. For example

$$horner(79x^5 + 56x^4 + 63x^2 + 57x - 59)$$

should return

$$-59 + (57 + (63 + (56 + 79x)x^2)x)x.$$

Do not use Maple's convert routine for this.

49.

- 50. Define the nth Legendre polynomial  $P_n(x)$  as the coefficient of  $y^n$  in the Taylor expansion of  $1/\sqrt{(1-2xy+y^2)}$ .
  - (a) Write a Maple-routine to compute the nth Legendre polynomial.
  - (b) Compute  $\int_{-1}^{\hat{1}} P_n(x) P_m(x)$  for several n and m and conjecture a general formula for this integral.
  - (c) Let  $Q_n(x)$  be defined by  $Q_0(x) = 1$ ,  $Q_1(x) = x$  and  $Q_n(x) = (n-1)/n(xQ_{n-1}(x) Q_{n-2}(x)) + xQ_{n-1}(x)$ . Write a Maple-routine to compute the polynomials  $Q_n$ .
  - (d) COmpare the polynomials  $P_n$  and  $Q_n$ .
- 51. An (ordinary, linear) differential operator of order n is an object of the following form:

$$P = \sum_{i=0}^{n} P_i(x) \left(\frac{d}{dx}\right)^i,$$

where the  $P_i(x)$  are functions of x (usually assumed to be differentiable infinitely many times). One can think of P as an operator acting on the set of functions, which are infinitely many times differentiable, as follows:

$$Pf = \sum_{i=0}^{n} P_i(x) \left( \frac{d^i f}{dx^i} \right).$$

Now it is clear how to add these operators and how to multiply them with functions. For operators P and Q we can also define their composition (or product) PQ as the unique operator R such that for all functions Rf = P(Qf). Notice that in general we have  $PQ \neq QP$ .

Implement differential operators in Maple, i.e. find a good data structure to represent them and write routines for the following operations:

- Multiplication with function
- Addition
- Composition
- Action on functions
- 52. A partition of a positive integer n is a list  $[a_1, a_2, \ldots, a_r]$  where  $a_1, a_2, \ldots, a_r$  are positive integers such that  $a_1 \ge a_2 \ge \cdots \ge a_r$  and  $n = a_1 + a_2 + \cdots + a_r$ .
  - (a) Write a Maple-routine partitions such that partitions (n) returns a list of all partitions of n.
    - Hint: First write a routine P such that P(n,k) returns all partitions  $[a_1,a_2,\ldots,a_r]$ .
  - (b) Let  $N_n$  denote the number of partitions of n. Write a Maple-routine N which computes  $N_n$  (without actually computing all permutations).

A second way to compute  $N_n$  can be extracted from the following identity, due to Euler.

$$\frac{1}{(1-x)(1-x^2)(1-x^3)\cdots} = 1 + N_1x + N_2x^2 + N_3x^3 + \cdots$$

Try to understand this identity and write a Maple-routine which computes  $N_n$  using this identity (use the Maple-routines taylor and coeff). Compare the speed of both routines.

53. In this exercise we will do some 'experimental mathematics'. Let

$$\prod_{n=1}^{\infty} (1 - x^n) = \sum_{n=0}^{\infty} a_n x^n.$$

- (a) Write a Maple-routine approx that computes  $\prod_{n=1}^{\infty} (1-x^n) \mod x^m$ .
- (b) Look at the differences of the exponents of succeeding terms appearing. Can you now guess what terms will appear.
- (c) Look at the coefficients of the terms appearing. Can you now guess what the coefficients will be in general.
- (d) Write now a Maple-routine that computes  $\prod_{n=1}^{\infty} (1-x^n) \mod x^m$  without computing a product of polynomials.
- 54. Check that a is a primitive nth root of unity in  $\mathbb{F}_p$ .
  - (a) p = 41, n = 8, a = 14.
  - (b) p = 97, n = 32, a = 28.
  - (c) p = 12289, n = 1024, a = -1987.
  - (d) p = 104857601, n = 1048576, a = -25430071.
- 55. (a) Write a Maple-routine fft to perform the fast Fourier transformation, i.e. when p is a prime number, n a power of 2,  $\omega$  a primitive nth root of unity in  $\mathbb{F}_p$  and  $a = [a_0, a_1, \ldots, a_{n-1}]$  a list of elements of  $\mathbb{F}_p$ , then  $fft(a, \omega, p)$  returns a list  $[A_0, A_1, \ldots, A_{n-1}]$ , the Fourier transform of a.
  - (b) Write a Maple-routine *ifft* for the inverse Fourier transformation.
  - (c) Check that your routines are each others inverse by the following example:  $a = [1,4,3,10,5,16,5,0], \omega = 2, p = 17.$
  - (d) Get an idea of the speed of your routines by studying the following big example:  $p=12289, n=2^{10}, \omega=-1987$  and a list of random numbers of length  $2^{10}$ .
- 56. Study the Maple-routines *FFT* and *iFFT*, the fast Fourier transforms and its inverse over the complex numbers (remind the on-line help). Compare the speed of these routines and your routines *fft* and *ifft*.
- 57. Write a Maple-routine to evaluate a polynomial in the *n*th roots of unity (of a finite prime field) using Horner's rule. Do the same again but now using your *fft*-routine from exercise 55. Compare the speed of both routines.
- 58. Write a Maple-routine to interpolate a polynomial in the *n*th roots of unity (of a finite prime field) using Lagrange's interpolation formula. Do the same again but now using your *ifft*-routine from exercise 55. Compare the speed of both routines.
- 59. Write a Maple-routine to multiply two polynomials using the fast Fourier transformation. Use your routines fft and ifft of exercise 55. Use your routine to multiply  $(x+1)^{15}$  and

- $(x+2)^{15}$  in  $\mathbb{F}_p[x]$  where p=104857601 (use the primitive 32th root of unity  $\omega=43262874$ ).
- 60. Compute all prime numbers  $\leq 2^{31}$  of the form  $1 + 2^{20}k$  (use the Maple-routine *isprime*). How many are there? Compare this with the result in the lecture notes.
- 61. (a) Write a Maple-routine to compute a primitive element of  $\mathbb{F}_p$  (use the Maple-routine numtheory/factorset/).
  - (b) Write a Maple-routine that returns  $[n, \omega]$ , where n is the largest power of 2 that divides p-1 and  $\omega$  is a primitive nth root of unity in  $\mathbb{F}_n$ .
  - (c) Use the routine of (b) to compute a primitive  $2^{20}$ th root of unity in  $F_p$ , where  $p = 2^{20} \cdot 625 + 1$ .
- 62. Do some experiments using your fft routine of exercise 55. You can use your routine of exercise 61 to find primitive roots of unity. Use random lists as input.
- 63. Write a Maple-routine to compute the resultant of two polynomials with coefficients in Q. Do this in two ways: using the definition of the resultant and using the recursive formula which comes from pseudo-division of the polynomials (see lecture notes). Compare the speed of these routines and the Maple-routine resultant.
- 64. Write a Maple-routine resultant\_ex which computes the resultant r of two polynomials f, g, with coefficients in  $\mathbb{Q}$ , and also polynomials a and b such that r = af + bg,  $\deg(a) < \deg(g)$  and  $\deg(b) < \deg(f)$ .

## REFERENCES

[1] K.O. Geddes, S.R. Czapor, G. Labahn Algorithms for computer algebra, 1992, Kluwer