

Handout WP: Writing Proofs.

1. Direct proof of an if-then statement.

Many theorems have the form "if p then q " where p, q are statements.
(Note: p, q are often composed of other statements with and's and or's).
A direct proof of $p \implies q$ works like this:

Assume: p .
To prove: q .

2. Proof by contrapositive.

The statement $p \implies q$ is *logically equivalent to* (means: same truth-table) the statement $\neg q \implies \neg p$. This means we can prove $p \implies q$ like this:

Assume: $\neg q$.
To prove: $\neg p$.

Remark: For any statement p you need to be able to compute $\neg p$.

3. Proof by contradiction.

One way to prove a statement S is as follows:

Assume: $\neg S$.
To prove: a contradiction.

Proving a contradiction means proving something that is obviously wrong (e.g. $x \neq x$, or the negation of a given/assumed/proved statement).

4. Proving an if-then statement by contradiction.

If S is the statement $p \implies q$ then $\neg S$ is logically equivalent to $p \wedge \neg q$.
So a proof-by-contradiction for $p \implies q$ works like this:

Assume: p and $\neg q$.
To prove: a contradiction.

5. Direct proof for a for-all statement.

If $P(x)$ is a statement involving x , to prove a statement like this $\forall_{x \in A} P(x)$ you write the following:

Let $x \in A$.
To prove: $P(x)$.

Explanation: When you write "Let $x \in A$ " then you are telling the reader that x is an element of A , but you are not specifying *which* element of A . That means that x could be any element of A . Once you proved $P(x)$ for such x then $P(x)$ must be true for any element of A .

What not to do: Suppose $f : S \rightarrow T$ and you want to prove $\forall_{t \in T} P(t)$. You should start with: "Let $t \in T$ ". But what if you do instead:

Some steps ... take $t = f(s)$... some steps ... hence $P(t)$.
Then you only proved $P(t)$ for some, but not all, elements of T .

6. Direct proof for an exists statement.

If $P(x)$ is a statement involving x , to prove a statement like this $\exists_{x \in A} P(x)$ you write the following:

Take $x :=$ [write down an expression].

If it is obvious that the expression you wrote down meets the requirements $x \in A$ and $P(x)$ then the proof is now complete; you have a 1-line proof! (only non-obvious requirements need to be checked). If you wrote a lot of text but not this 1 line, then your proof is still not complete.

7. Proving a for-all statement by contradiction.

If S is the statement $\forall_{x \in A} P(x)$ then $\neg S$ is the statement $\exists_{x \in A} \neg P(x)$. So if we follow item 3 (i.e. prove S by assuming $\neg S$ and then proving a contradiction) then the proof would start like this:

Assume: x is an element of A and $\neg P(x)$.

To prove: a contradiction.

8. Proving an exists-statement by contradiction.

If S is the statement $\exists_{x \in A} P(x)$ then $\neg S$ is the statement $\forall_{x \in A} \neg P(x)$.

Assume: $\neg P(x)$ for every $x \in A$.

To prove: a contradiction.

9. Proof by cases.

Suppose we need to prove some statement $P(x)$ where x can only have a few possible values. In that case, we can write a separate proof for each possible value.

Example 1: Suppose we have to prove $P(x)$ but we know that x can only be u or v or w . Then we write three proofs:

Case 1: $x = u$. Write a proof for $P(u)$.

Case 2: $x = v$. Write a proof for $P(v)$.

Case 3: $x = w$. Write a proof for $P(w)$.

Example 2: Suppose we have to prove a statement p but there is some other statement q such that we can easily find a proof for $q \implies p$. Then we can do the following, we split the proof in two cases:

Case 1: Assume q is true and prove p under that assumption.

Case 2: Assume q is false now prove p under that assumption.

These cases combined provide a complete proof for p .

10. Proving an and statement.

To prove $p \wedge q$ write two separate proofs:

To prove: p

To prove: q

11. **Proving an iff statement** (iff = “if and only if”).

The statement $p \iff q$ is logically equivalent to $(p \implies q) \wedge (q \implies p)$.

So to prove $p \iff q$ you have to write two proofs:

To prove: $p \implies q$

To prove: $q \implies p$

Proofs of “if and only if” statements in math books typically look like:

Assume $p \dots$ some math \dots hence q . For the converse \dots

Recall that $q \implies p$ is called the **converse** of $p \implies q$. Math books assume you know that “For the converse” means “To prove: $q \implies p$ ”.

12. **Proving an or statement.**

The statement $p \vee q$ means that at least one of p or q is true. But which one? That question makes it tricky to give a direct proof of $p \vee q$. But you can always replace a statement by a logically equivalent statement.

We have several options:

(1) $p \vee q$ is logically equivalent to $\neg p \implies q$.

(2) $p \vee q$ is logically equivalent to $\neg q \implies p$.

(3) $\neg(p \vee q)$ is logically equivalent to $\neg p \wedge \neg q$.

That gives us several ways to prove $p \vee q$.

Method (1): Assume $\neg p$. To prove: q .

Method (2): Assume $\neg q$. To prove: p .

Method (3): Assume $\neg p$ and $\neg q$. To prove: a contradiction.

Which method is best? That depends on what you already know. Suppose for instance you see a theorem in the book of the form $\neg q \implies r$. In that case you want to try method (2) (assume $\neg q$) because then you can use the theorem to conclude r . Hopefully that brings you closer to goal p .

13. **Using an or statement.**

Suppose you want to prove r , and you are given $p \vee q$ which means that at least one of p or q is true. But which one? So we write two proofs:

Assume p . To prove r .

Assume q . To prove r .

The two proofs combined show that $p \vee q$ implies r .

This is the same as “Proof by cases” from item 9; given $p \vee q$ you distinguish two cases: Case 1: assume p , to prove r . Case 2: assume q , to prove r .

14. If you want to **use a for-all statement** like $\forall_{x \in A} P(x)$ to prove another statement, often the best strategy is to make a clever choice for one particular element of A , and then use the fact that P is true for that element.

15. If you want to **use an exists statement** like $\exists_{x \in A} P(x)$ to prove another statement, then you *may not choose* x . All you know is $x \in A$ and $P(x)$.