

Closed Form Solutions of Linear Difference Equations

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Object of the Thesis:

Algorithm *sol/ver* that solves difference operators.

- 1 Transformations
- 2 Invariant Data
- 3 Table of base equations

Outline

- 1 Difference Operator
- 2 Example
- 3 Transformations
- 4 Main Idea
- 5 Invariant Local Data
 - Finite Singularity
 - Generalized Exponent
- 6 Liouvillian
- 7 Special Functions

Linear Difference Equation

- Difference Equation:

Let $DE : \mathbb{C}^{n+2} \rightarrow \mathbb{C}$. Then a difference equation is an equation of the form

$$DE(f(x), f(x+1), \dots, f(x+n), x) = 0 \quad (n \geq 1)$$

- A recurrence relation

Let $R : \mathbb{C}^{n+1} \rightarrow \mathbb{C}$. Then a recurrence relation is an equation of the form

$$f(x+n) = R(f(x), f(x+1), \dots, f(x+n-1), x) \quad (n \geq 1)$$

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Linear Difference Equation

A difference equation is called linear, if it is in the form of

$$a_n(x)f(x+n) + a_{n-1}(x)f(x+n-1) + \cdots + a_0(x)f(x) + a(x) = 0$$

where $a, a_i : \mathbb{C} \rightarrow \mathbb{C}$ for $i = 0, \dots, n$.

Then it naturally defines a recurrence relation by

$$f(x+n) = -\frac{a_{n-1}(x)}{a_n(x)}f(x+n-1) - \cdots - \frac{a_0(x)}{a_n(x)}f(x) - \frac{a(x)}{a_n(x)}$$

A difference equation is called homogeneous if $a(x) = 0$.

In this talk we will only consider homogeneous linear difference equations with coefficients in $\mathbb{C}(x)$.

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Linear Difference Operator

Let τ be the shift operator: $\tau(u(x)) = u(x + 1)$

Then a Linear Difference Operator L is

$$L = a_n \tau^n + a_{n-1} \tau^{n-1} + \cdots + a_0 \tau^0 \text{ where } a_i \in \mathbb{C}(x).$$

L corresponds to a difference equation

$$a_n(x)f(x+n) + a_{n-1}(x)f(x+n-1) + \cdots + a_0(x)f(x) = 0.$$

Example:

- If $L = \tau - x$ then the equation $L(f(x)) = 0$ is $f(x+1) - xf(x) = 0$ and $\Gamma(x)$ is a solution of L .

We will see some examples of what *solver* can do.
(with Maple worksheet)

GT-Transformation

Notation:

- $V(L)$ = solution space of L .
- ① Term Product: L_2 is a term product of L_1 when $V(L_2)$ can be written as $V(L_1)$ multiplied by a hypergeometric term.
- ② Gauge Equivalence: L_2 is gauge equivalent to L_1 if there exists $G \in \mathbb{C}(x)[\tau]$ that bijectively maps $V(L_1)$ to $V(L_2)$.
- ③ GT-Equivalence: $L_2 \sim_{gt} L_1$ if a combination of (1) and (2) can map $V(L_1)$ to $V(L_2)$. Such map is called GT-Transformation.

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Main Idea

Observation:

If two operators are gt-equivalent and if one of them has closed form solutions, then so does the other.

Idea:

- Find base equations: Find parameterized families of equations with known solutions.
- Solve every equation \sim_{gt} to a base equation.

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Questions

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Yes

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- \bullet $LbIK = z\tau^2 + (2 + 2v + 2x)\tau - z$
 Solutions: Modified Bessel functions of the first and second kind, $I_{v+x}(z)$ and $K_{v+x}(-z)$
- \bullet $LbJY = z\tau^2 - (2 + 2v + 2x)\tau + z$
 Solutions: Bessel functions of the first and second kind, $J_{v+x}(z)$ and $Y_{v+x}(z)$
- \bullet $LWW = \tau^2 + (z - 2v - 2x - 2)\tau - v - x - \frac{1}{4} - v^2 - 2vx - x^2 + n^2$
 Solution: Whittaker function $W_{x,n}(z)$
- \bullet $LWM = \tau^2(2n + 2v + 3 + 2x) + (2z - 4v - 4x - 4)\tau - 2n + 1 + 2v + 2x$
 Solution: Whittaker function $M_{x,n}(z)$
- \bullet $L2F1 = (z - 1)(a + x + 1)\tau^2 + (-z + 2 - za - zx + 2a + 2x + zb - c)\tau - a + c - 1 - x$
 Solution: Hypergeometric function ${}_2F_1(a + x, b; c; z)$
- \bullet $Ljc = \tau^2 - \frac{1}{2} \frac{(2x+3+a+b)(a^2-b^2+(2x+a+b+2)(2x+4+a+b)z)}{(x+2)(x+2+a+b)(2x+a+b+2)} \tau + \frac{(x+1+a)(x+1+b)(2x+4+a+b)}{(x+2)(x+2+a+b)(2x+a+b+2)}$
 Solution: Jacobian polynomial $P_x^{a,b}(z)$
- \bullet $Lgd = \tau^2 - \frac{(2x+3)z}{x+2} \tau + \frac{x+1}{x+2}$
 Solution: Legendre functions $P_x(z)$ and $Q_x(z)$
- \bullet $Lgr = \tau^2 - \frac{2x+3+\alpha-z}{x+2} \tau + \frac{x+1+\alpha}{x+2}$
 Solution: Laguerre polynomial $L_x^{(\alpha)}(z)$
- \bullet $Lgb = \tau^2 - \frac{2z(m+x+1)}{x+2} \tau - \frac{2m+x}{x+2}$
 Solution: Gegenbauer polynomial $C_x^m(z)$
- \bullet $Lgr1 = (x + 2)\tau^2 + (x + z - b + 1)\tau + z$
 Solution: Laguerre polynomial $L_x^{(b-x)}(z)$
- \bullet $Lkm = (a + x + 1)\tau^2 + (-2a - 2x - 2 + b - c)\tau + a + x + 1 - b$
 Solution: Kummer's function $M(a + x, b, c)$
- \bullet $L2F0 = \tau^2 + (-zb + zx + z + za - 1)\tau + z(b - x - 1)$
 Solution: Hypergeometric function ${}_2F_0(a, b - x; ; z)$
- \bullet $Lge = (x + 2)\tau^2 + (-ab - d + (a + 1)(1 + x))\tau + ax - a(b + d)$
 Solution: Sequences whose ordinary generating function is $(1 + ax)^b(1 + bx)^d$

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Local data

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Local data

Main Algorithm

- 1 Compute local data of L .
- 2 Compare the data with those in the table and find a base equation that matches the data. If there is no such base equation then return \emptyset .
 - 1 Compute candidate values for each parameters.
 - 2 Construct a set cdd by plugging values found in step 1 to corresponding parameters.
- 3 For each $L_c \in cdd$ check if $L \sim_{gt} L_c$ and if so
 - 1 Generate a basis of solutions or a solution of L_c by plugging in corresponding parameters.
 - 2 Apply the term transformation and the gauge transformation to the result from 1.
 - 3 Return the result of step 2 as output and stop the algorithm.

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Algorithm

TryBessel:

Input: $L \in \mathbb{C}(x)[\tau]$

- 1 Compute the local data of $L_{v,z} = z\tau^2 + (2 + 2v + 2x)\tau - z$ (Bessel recurrence).
- 2 Compute local data of L that is invariant under \sim_{gt} .
- 3 Compare the local data of $L_{v,z}$ with that of L .
- 4 If compatible, compute v, z from this comparison.
- 5 Check if $L \sim_{gt} L_{v,z}$, and if so, return solution(s).

Note: Step 1 is done only once, and then stored in a table.

Remark: Checking $L \sim_{gt} L_{v,z}$ and computing the gt-transformation can only be done after we have found the values of the parameter v, z .

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Invariant Local Data

Question: If $L \sim_{gt} L_{v,z}$, how to find v, z from L ?

Need data that is invariant under \sim_{gt}

Two sources

- 1 Finite Singularities (valuation growths)
- 2 Singularity at ∞ (generalized exponents)

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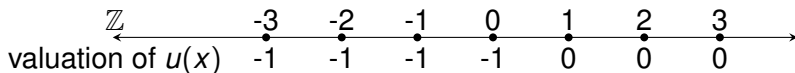
Finite Singularity: Valuation Growth

Suppose $L_1 \sim_g L_2$ and $G = r_k(x)\tau^k + \cdots + r_0(x)$, $r_i(x) \in \mathbb{C}(x)$
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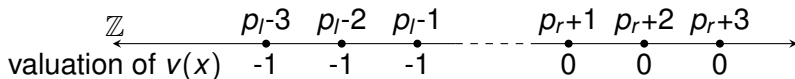
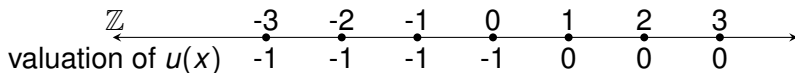
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To calculate $f(s+n)$ with values of $f(s), \dots, f(s+n-1)$, $s \in \mathbb{C}$,

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Definition

Let $L = a_n\tau^n + \cdots + a_0\tau^0$ with $a_i \in \mathbb{C}[x]$. $q \in \mathbb{C}$ is called a *problem point* of L if q is a root of the polynomial $a_0(x)a_n(x-n)$. $p \in \mathbb{C}/\mathbb{Z}$ is called a *finite singularity* of L if it contains a problem point.

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Finite Singularity: Valuation Growth

Definition

Let $u(x) \in \mathbb{C}(x)$ be a non-zero meromorphic function. The *valuation growth* of $u(x)$ at $p = q + \mathbb{Z}$ is

$$\liminf_{n \rightarrow \infty} (\text{order of } u(x) \text{ at } x = n + q) \\ - \liminf_{n \rightarrow \infty} (\text{order of } u(x) \text{ at } x = -n + q)$$

Definition

Let $p \in \mathbb{C}/\mathbb{Z}$ and L be a difference operator. Then $\text{Min}_p(L)$ resp. $\text{Max}_p(L)$ is the minimum resp. maximum valuation growth at p , taken over all meromorphic solutions of L .

Theorem

If $L_1 \sim_g L_2$ then they have the same $\text{Min}_p, \text{Max}_p$ for all $p \in \mathbb{C}/\mathbb{Z}$.

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Finite Singularity: Valuation Growth

Theorem

$\text{Max}_p - \text{Min}_p$ is \sim_{gt} invariant for all $p \in \mathbb{C}/\mathbb{Z}$.

Invariant data: Compute all $p \in \mathbb{C}/\mathbb{Z}$ for which $\text{Max}_p \neq \text{Min}_p$
store $[p, \text{Max}_p - \text{Min}_p]$ for all such p .

Note: Since $p \in \mathbb{C}/\mathbb{Z}$ and not in \mathbb{C} , the parameters computed from such data are determined mod $r\mathbb{Z}$ for some $r \in \mathbb{Q}$.

Suppose we need parameter $\nu \bmod \mathbb{Z}$ but find it mod $\frac{1}{2}\mathbb{Z}$, then we need to check two cases.

Finite Singularity: Valuation Growth

Theorem

$\text{Max}_p - \text{Min}_p$ is \sim_{gt} invariant for all $p \in \mathbb{C}/\mathbb{Z}$.

Invariant data: Compute all $p \in \mathbb{C}/\mathbb{Z}$ for which $\text{Max}_p \neq \text{Min}_p$
store $[p, \text{Max}_p - \text{Min}_p]$ for all such p .

Note: Since $p \in \mathbb{C}/\mathbb{Z}$ and not in \mathbb{C} , the parameters computed from such data are determined mod $r\mathbb{Z}$ for some $r \in \mathbb{Q}$.
Suppose we need parameter $\nu \bmod \mathbb{Z}$ but find it mod $\frac{1}{2}\mathbb{Z}$, then we need to check two cases.

Singularity at ∞ : Generalized Exponent

Definition

If $\tau - ct^v(1 + \sum_{i=1}^{\infty} a_i t^{\frac{i}{r}})$, with $t = 1/x$, is right hand factor of L for some $v \in \frac{1}{r}\mathbb{Z}$, $c \in \mathbb{C}^*$, $a_i \in \mathbb{C}$, $r \in \mathbb{N}$, then the dominant term $ct^v(1 + a_1 t^{\frac{1}{r}} + \cdots + a_r t^1)$ is called a *generalized exponent* of L .

We say two generalized exponents

$g_1 = c_1 t^{v_1}(1 + a_1 t^{\frac{1}{r}} + \cdots + a_r t^1)$ and

$g_2 = c_2 t^{v_2}(1 + b_1 t^{\frac{1}{r}} + \cdots + b_r t^1)$ are equivalent if

$c_1 = c_2$, $v_1 = v_2$, $a_i = b_i$ for $i = 1 \dots r-1$ and $a_r \equiv b_r \pmod{\frac{1}{r}\mathbb{Z}}$ and denote $g_1 \sim_r g_2$

Theorem

Generalized exponents are invariant up to \sim_r under Gauge equivalence.

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Singularity at ∞ : Generalized Exponent

Generalized exponents are not invariant under term-product.

Definition

Suppose $\text{ord}(L) = 2$ and let $\text{genexp}(L) = \{a_1, a_2\}$ such that $v(a_1) \geq v(a_2)$. Then we define the *set of quotient of the two generalized exponents* as

if $v(a_1) > v(a_2)$

$$\text{Gquo}(L) = \left\{ \frac{a_1}{a_2} \right\} \text{ and}$$

if $v(a_1) = v(a_2)$ then we define

$$\text{Gquo}(L) = \left\{ \frac{a_1}{a_2}, \frac{a_2}{a_1} \right\}.$$

Theorem

If $L_1 \sim_{gt} L_2$ then $\text{Gquo}(L_1) = \text{Gquo}(L_2) \bmod \sim_r$

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Outline

- 1 Difference Operator
- 2 Example
- 3 Transformations
- 4 Main Idea
- 5 Invariant Local Data
 - Finite Singularity
 - Generalized Exponent
- 6 Liouvillian**
- 7 Special Functions

Liouvillian Solutions of Linear Difference Equations: Property

Theorem (Hendriks Singer 1999)

If $L = a_n\tau^n + \cdots + a_0\tau^0$ is irreducible then

\exists Liouvillian Solutions $\iff \exists b_0 \in \mathbb{C}(x)$ such that

$$a_n\tau^n + \cdots + a_0\tau^0 \sim_g \tau^n + b_0\tau^0$$

Remark

Operators of the form $\tau^n + b_0\tau^0$ are easy to solve, so if we know b_0 then we can solve L .

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for some unknown $b_0 \in \mathbb{C}(x)$.

If we can find b_0 then we can solve $\tau^n + b_0 \tau^0$ and hence solve L .

Notation

write $b_0 = c\phi$ where $\phi = \frac{\text{monic poly}}{\text{monic poly}}$ and $c \in \mathbb{C}^*$.

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c is easy to compute, the main task is to compute ϕ .

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Let $L = a_n \tau^n + \cdots + a_0 \tau^0 \in \mathbb{C}[x][\tau]$ then the finite singularities of L are $Sing = \{q + \mathbb{Z} \in \mathbb{C}/\mathbb{Z} \mid q \text{ is root of } a_0 a_n\}$

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If $q_1 + \mathbb{Z}, \dots, q_k + \mathbb{Z}$ are the finite singularities then we may

assume $\phi = \prod_{i=1}^k \prod_{j=0}^{n-1} (x - q_i - j)^{k_{i,j}}$ with $k_{i,j} \in \mathbb{Z}$.

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Example of Operator of order 2 with one finite singularity at $p = \mathbb{Z}$

Suppose $L = a_2\tau^2 + a_1\tau + a_0$ and that

$$L \sim_g \tau^2 + c \cdot x^{k_0}(x-1)^{k_1}$$

- 1 c can be computed from a_0/a_2
- 2 $k_0 + k_1$ can be computed from a_0/a_2
- 3 $\max\{k_0, k_1\} = \text{Max}_{\mathbb{Z}}(L)$
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Items 2, 3, 4 determine k_0, k_1 up to a permutation.

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Example from

The On-Line Encyclopedia of Integer Sequences™ (OEIS™)

A000246 = (1, 1, 1, 3, 9, 45, 225, 1575, 11025, 99225, ...)

Number of permutations in the symmetric group S_n that have odd order.

- $\tau^2 - \tau - x(x+1)$
- $Sing = \{\mathbb{Z}\}$ and $c = 1$.
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$$min = 0, \quad max = 2, \quad sum = 2$$

- So the exponents of $x \cdots (x-1) \cdots$ must be a permutation of 0, 2
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Suppose $L = a_3\tau^3 + a_2\tau^2 + a_1\tau + a_0$ is gauge equivalent to

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This determines k_0, k_1, k_2 up to a permutation, and also l_0, l_1, l_2 up to a permutation.

Worst case is $3! \cdot 3!$ combinations (actually: 1/3 of that).

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Liouvillian Solutions of Linear Difference Equations:

Example $L = x\tau^3 + \tau^2 - (x+1)\tau - x(x+1)^2(2x-1)$

- $Sing = \{\mathbb{Z}, \frac{1}{2} + \mathbb{Z}\}$ and $c = -2$.

- At \mathbb{Z} ,

$$min = 0, \quad max = 1, \quad sum = 2$$

So the exponents of $x^{\cdots}(x-1)^{\cdots}(x-2)^{\cdots}$ must be a permutation of 0, 1, 1

- At $\frac{1}{2} + \mathbb{Z}$,

$$min = 0, \quad max = 1, \quad sum = 1$$

So the exponents of $(x - \frac{1}{2})^{\cdots}(x - \frac{3}{2})^{\cdots}(x - \frac{5}{2})^{\cdots}$ must be a permutation of 0, 0, 1

Liouvillian Solutions of Linear Difference Equations:

Example $L = x\tau^3 + \tau^2 - (x+1)\tau - x(x+1)^2(2x-1)$

- $Sing = \{\mathbb{Z}, \frac{1}{2} + \mathbb{Z}\}$ and $c = -2$.

- At \mathbb{Z} ,

$$min = 0, \quad max = 1, \quad sum = 2$$

So the exponents of $x^{\cdots}(x-1)^{\cdots}(x-2)^{\cdots}$ must be a permutation of 0, 1, 1

- At $\frac{1}{2} + \mathbb{Z}$,

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Candidates of $c\phi$ are

- ① $-2x^1(x-1)^1(x-2)^0(x-1/2)^1(x-3/2)^0(x-5/2)^0$
- ② $-2x^1(x-1)^1(x-2)^0(x-1/2)^0(x-3/2)^1(x-5/2)^0$
- ③ $-2x^1(x-1)^1(x-2)^0(x-1/2)^0(x-3/2)^0(x-5/2)^1$
- ④ $-2x^0(x-1)^1(x-2)^1(x-1/2)^0(x-3/2)^0(x-5/2)^1$
- ⑤ $-2x^0(x-1)^1(x-2)^1(x-1/2)^0(x-3/2)^1(x-5/2)^0$
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- ⑦ $-2x^1(x-1)^0(x-2)^1(x-1/2)^1(x-3/2)^0(x-5/2)^0$
- ⑧ $-2x^1(x-1)^0(x-2)^1(x-1/2)^0(x-3/2)^0(x-5/2)^1$
- ⑨ $-2x^1(x-1)^0(x-2)^1(x-1/2)^0(x-3/2)^1(x-5/2)^0$

Liouvillian Solutions of Linear Difference Equations:

Example $L = x\tau^3 + \tau^2 - (x+1)\tau - x(x+1)^2(2x-1)$

Remark

$$\tau^n - c\phi \sim_g \tau^n - c\tau^k(\phi) \text{ for } k = 1 \dots n-1$$

$$\textcircled{1} -2x^1(x-1)^1(x-2)^0(x-1/2)^1(x-3/2)^0(x-5/2)^0$$

$$\textcircled{2} -2x^1(x-1)^1(x-2)^0(x-1/2)^0(x-3/2)^1(x-5/2)^0$$

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Only need to try 1, 2, 3, the others are redundant.

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Only need to try 1, 2, 3, the others are redundant.

Liouvillian Solutions of Linear Difference Equations:

Example $L = x\tau^3 + \tau^2 - (x+1)\tau - x(x+1)^2(2x-1)$

- $\tau^3 - 2x(x-1)(x-1/2)$ is gauge equivalent to L
- Gauge transformation is $\tau + x - 1$.
- Basis of solutions of $\tau^3 - 2x(x-1)(x-1/2)$ is

$$\{(\xi^k)^x v(x)\} \quad \text{for } k = 0 \dots 2$$

where $v(x) = 3^x 2^{x/3} \Gamma(\frac{x}{3}) \Gamma(\frac{x-1}{3}) \Gamma(\frac{x-2}{3})$ and $\xi^3 = 1$.

- Thus, Basis of solutions of L is

$$\{(\xi^k)^{x+1} v(x+1) + (x-1)(\xi^k)^x v(x)\} \quad \text{for } k = 0 \dots 2$$

Liouvillian Solutions of Linear Difference Equations:

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Outline

- 1 Difference Operator
- 2 Example
- 3 Transformations
- 4 Main Idea
- 5 Invariant Local Data
 - Finite Singularity
 - Generalized Exponent
- 6 Liouvillian
- 7 Special Functions**

- $LbIK = z\tau^2 + (2 + 2v + 2x)\tau - z$
 Solutions: Modified Bessel functions of the first and second kind, $I_{v+x}(z)$ and $K_{v+x}(-z)$
- $LbJY = z\tau^2 - (2 + 2v + 2x)\tau + z$
 Solutions: Bessel functions of the first and second kind, $J_{v+x}(z)$ and $Y_{v+x}(z)$
- $LWW = \tau^2 + (z - 2v - 2x - 2)\tau - v - x - \frac{1}{4} - v^2 - 2vx - x^2 + n^2$
 Solution: Whittaker function $W_{x,n}(z)$
- $LWM = \tau^2(2n + 2v + 3 + 2x) + (2z - 4v - 4x - 4)\tau - 2n + 1 + 2v + 2x$
 Solution: Whittaker function $M_{x,n}(z)$
- $L2F1 = (z - 1)(a + x + 1)\tau^2 + (-z + 2 - za - zx + 2a + 2x + zb - c)\tau - a + c - 1 - x$
 Solution: Hypergeometric function ${}_2F_1(a + x, b; c; z)$
- $Ljc = \tau^2 - \frac{1}{2} \frac{(2x+3+a+b)(a^2-b^2+(2x+a+b+2)(2x+4+a+b)z)}{(x+2)(x+2+a+b)(2x+a+b+2)} \tau + \frac{(x+1+a)(x+1+b)(2x+4+a+b)}{(x+2)(x+2+a+b)(2x+a+b+2)}$
 Solution: Jacobian polynomial $P_x^{a,b}(z)$
- $Lgd = \tau^2 - \frac{(2x+3)z}{x+2} \tau + \frac{x+1}{x+2}$
 Solution: Legendre functions $P_x(z)$ and $Q_x(z)$
- $Lgr = \tau^2 - \frac{2x+3+\alpha-z}{x+2} \tau + \frac{x+1+\alpha}{x+2}$
 Solution: Laguerre polynomial $L_x^{(\alpha)}(z)$
- $Lgb = \tau^2 - \frac{2z(m+x+1)}{x+2} \tau - \frac{2m+x}{x+2}$
 Solution: Gegenbauer polynomial $C_x^m(z)$
- $Lgr1 = (x + 2)\tau^2 + (x + z - b + 1)\tau + z$
 Solution: Laguerre polynomial $L_x^{(b-x)}(z)$
- $Lkm = (a + x + 1)\tau^2 + (-2a - 2x - 2 + b - c)\tau + a + x + 1 - b$
 Solution: Kummer's function $M(a + x, b, c)$
- $L2F0 = \tau^2 + (-zb + zx + z + za - 1)\tau + z(b - x - 1)$
 Solution: Hypergeometric function ${}_2F_0(a, b - x; ; z)$
- $Lge = (x + 2)\tau^2 + (-ab - d + (a + 1)(1 + x))\tau + ax - a(b + d)$
 Solution: Sequences whose ordinary generating function is $(1 + ax)^b(1 + x)^d$

Special Functions: Functions and their Local Data

Operator	Val	Gquo
<i>LbIK</i>	$\{\}$	$\{-\frac{1}{4}T^2z^2(1 - (1 + 2v)T)\}$
<i>LbJY</i>	$\{\}$	$\{\frac{1}{4}T^2z^2(1 - (1 + 2v)T)\}$
<i>LWW</i>	$\{[-n + \frac{1}{2} - v, 1], [n + \frac{1}{2} - v, 1]\}$	$\{-3 - 2\sqrt{2}(1 - \frac{1}{2}\sqrt{2}z)T, -3 + 2\sqrt{2}(1 + \frac{1}{2}\sqrt{2}z)T\}$
<i>LWM</i>	$\{[-n + \frac{1}{2} - v, 1], [n + \frac{1}{2} - v, 1]\}$	$\{1 - 2\sqrt{-z}T - 2zT^2, 1 + 2\sqrt{-z}T - 2zT^2\}$
<i>L2F1</i>	$\{[-a + c, 1], [-a, 1]\}$	$\{-\frac{1}{z-1}(1 + (2b - c)T), (-z + 1)(1 + (-2b + c)T)\}$
<i>Ljc</i>	$\{[0, 1], [-a, 1], [-b, 1], [-a - b, 1]\}$	$\{2z^2 - 2z\sqrt{z^2 - 1} - 1, 2z^2 + 2z\sqrt{z^2 - 1} - 1\}$
<i>Lgd</i>	$\{[0, 2]\}$	$\{2z^2 - 2z\sqrt{z^2 - 1} - 1, 2z^2 + 2z\sqrt{z^2 - 1} - 1\}$
<i>Lgr</i>	$\{[0, 1], [-\alpha, 1]\}$	$\{1 + 2\sqrt{-z}T - 2zT^2, 1 - 2\sqrt{-z}T - 2zT^2\}$
<i>Lgr1</i>	$\{[0, 1]\}$	$\{zT(1 + 2bT)\}$
<i>Lgb</i>	$\{[0, 1], [-2m, 1]\}$	$\{-2z\sqrt{z^2 + 1} - 2z^2 - 1, 2z\sqrt{z^2 + 1} - 2z^2 - 1\}$
<i>Lkm</i>	$\{[-a, 1], [-a + b, 1]\}$	$\{1 - 2\sqrt{c}T + 2cT^2, 1 + 2\sqrt{c}T + 2cT^2\}$
<i>L2F0</i>	$\{[b, 1]\}$	$\{\frac{T}{z}(1 + (b - 2a)T)\}$
<i>Lge</i>	$\{[0, 1], [b + d, 1]\}$	$\{a(1 + (d - b)T), \frac{1}{a}(1 + (-b - d)T)\}$

Effectiveness of *solver*

Found 10,659 sequences in OEIS™ that satisfy a second order recurrence but not a first order recurrence.

- 9,455 were reducible
- 161 irreducible Liouvillian
- 86 Bessel
- 330 Legendre
- 374 Hermite
- 21 Jacobi
- 8 Kummer
- 44 Laguerre
- $7 {}_2F_1$
- $14 {}_2F_0$
- 77 Generating function $(1+x)^a(1+bx)^c$
- 82 Not yet solved

Example from

The On-Line Encyclopedia of Integer Sequences™ (OEIS™)

A096121 = (2, 8, 60, 816, 17520, 550080, 23839200,...)

Number of full spectrum rook's walks on a $(2 \times n)$ board.

- Difference Operator: $\tau^2 - (1 + x)(x + 2)\tau - (1 + x)(x + 2)$
- Val: $\{\}$
- Gquo: $\{-T^2(1 - 3T)\}$

Modified Bessel functions of the first and second kind,
 $I_{\nu+x}(z)$ and $K_{\nu+x}(-z)$.

- Difference Operator: $z\tau^2 + (2 + 2\nu + 2x)\tau - z$
- Val: $\{\}$
- Gquo: $\{-\frac{1}{4}T^2z^2(1 - (1 + 2\nu)T)\}$

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Example from

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- Comparing G_{quo} ,
 $\{-T^2(1-3T)\}$ and $\{-\frac{1}{4}z^2T^2(1-(1+2v)T)\}$,
 we get candidates of $z = \{2, -2\}$
 and candidates of $v = \{\frac{1}{2}, 1\}$
- We get four candidates to check \sim_{gt} ,

$$2\tau^2 - (2x+4)\tau - 2, \quad 2\tau^2 - (2x+3)\tau - 2$$

$$2\tau^2 + (2x+4)\tau - 2, \quad 2\tau^2 + (2x+3)\tau - 2.$$

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- Comparing Gquo,
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$$2\tau^2 + (2x+4)\tau - 2, 2\tau^2 + (2x+3)\tau - 2.$$

Example from

The On-Line Encyclopedia of Integer Sequences™ (OEIS™)

- $2\tau^2 - (2x + 4)\tau - 2 \sim_{gt} L$ (When $v = 1, z = -2$)
- Term-transformation is $x + 2$ and gauge-transformation is 1.
- Applying gt-transformation to $l_{1+x}(2)$ and $K_{1+x}(-2)$ we get basis of a basis of solutions of L ,

$$\{l_{1+x}(2)\Gamma(x+2), K_{1+x}(-2)\Gamma(x)\}$$

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Example from

The On-Line Encyclopedia of Integer Sequences™ (OEIS™)

A000262 = (1, 1, 3, 13, 73, 501, 4051, 37633, 394353,...)

Number of “sets of lists”:

number of partitions of $\{1, \dots, n\}$ into any number of lists.

- Difference Operator: $\tau^2 - (3 + 2x)\tau + x(x + 1)$
- Val: $\{[0, 2]\}$
- Gquo: $\{1 - 2T + 2T^2, 1 + 2T + 2T^2\}$

Laguerre polynomial $L_x^{(\alpha)}(z)$.

- Difference Operator: $Lgr = \tau^2 - \frac{2x+3+\alpha-z}{x+2}\tau + \frac{x+1+\alpha}{x+2}$
- Val: $\{[0, 1], [-\alpha, 1]\}$
- Gquo: $\{1 - 2\sqrt{-z}T - 2zT^2, 1 + 2\sqrt{-z}T - 2zT^2\}$

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- Gquo: $\{1 - 2\sqrt{-z}T - 2zT^2, 1 + 2\sqrt{-z}T - 2zT^2\}$

Example from

The On-Line Encyclopedia of Integer Sequences™ (OEIS™)

- Comparing G_{quo} ,
 $\{1 - 2T + 2T^2, 1 + 2T + 2T^2\}$ and
 $\{1 - 2\sqrt{-z}T - 2zT^2, 1 + 2\sqrt{-z}T - 2zT^2\}$,
 we get $z = -1$.
- $\text{Val} = \{[0, 2]\}$ is a special case of Lgr when $\alpha = 0$.
- We get one candidate to check \sim_{gt} ,

$$\tau^2 - \frac{2x+4}{x+2}\tau + \frac{x+1}{x+2}$$

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Example from

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- Term-transformation is x and gauge-transformation is $\frac{x+1}{x}\tau - \frac{x^2+2x}{x}$.
- Applying gt-transformation to $L_x^{(0)}(-1)$,

$$\{(x+1)L_{x+1}^{(0)}(-1) - (x+2)L_x^{(0)}(-1)\}\Gamma(x)$$

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Example from

The On-Line Encyclopedia of Integer Sequences™ (OEIS™)

A068770 = (1, 1, 16, 264, 4480, 77952, 1386496, 25135616,...)
Generalized Catalan numbers.

- Difference Operator: $(3 + x)\tau^2 + (-48 - 32x)\tau + 224x$
- Val: $\{[0, 2]\}$
- Gquo: $\{\frac{9}{7} - \frac{4}{7}\sqrt{2}, \frac{9}{7} + \frac{4}{7}\sqrt{2}\}$

Jacobian polynomial $P_x^{a,b}(z)$

- Difference Operator: $Lgd = \tau^2 - \frac{(2x+3)z}{x+2}\tau + \frac{x+1}{x+2}$
- Val: $\{[0, 2]\}$
- Gquo: $\{2z^2 - 2z\sqrt{z^2 - 1} - 1, 2z^2 + 2z\sqrt{z^2 - 1} - 1\}$

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Example from

The On-Line Encyclopedia of Integer Sequences™ (OEIS™)

- Comparing G_{quo} ,
 $\{\frac{9}{7} - \frac{4}{7}\sqrt{2}, \frac{9}{7} + \frac{4}{7}\sqrt{2}\}$ and
 $\{2z^2 - 2z\sqrt{z^2 - 1} - 1, 2z^2 + 2z\sqrt{z^2 - 1} - 1\}$,
 we get candidates of $z = \{\frac{2}{7}\sqrt{14}, -\frac{2}{7}\sqrt{14}\}$.
- $\text{Val} = \{[0, 2]\}$ is used to find the right base equation.

We get 2 candidate to check \sim_{gt} ,

$$\tau^2 - \frac{2(2x+3)\sqrt{14}}{7(x+2)}\tau + \frac{1+x}{x+2}$$

$$\tau^2 + \frac{2(2x+3)\sqrt{14}}{7(x+2)}\tau - \frac{1+x}{x+2}$$

Example from

The On-Line Encyclopedia of Integer Sequences™ (OEIS™)

- Comparing G_{quo} ,
 $\{\frac{9}{7} - \frac{4}{7}\sqrt{2}, \frac{9}{7} + \frac{4}{7}\sqrt{2}\}$ and
 $\{2z^2 - 2z\sqrt{z^2 - 1} - 1, 2z^2 + 2z\sqrt{z^2 - 1} - 1\}$,
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- $\text{Val} = \{[0, 2]\}$ is used to find the right base equation.

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$$\tau^2 + \frac{2(2x+3)\sqrt{14}}{7(x+2)}\tau - \frac{1+x}{x+2}$$

Example from

The On-Line Encyclopedia of Integer Sequences™ (OEIS™)

- $\tau^2 - \frac{2}{7} \frac{(2x+3)\sqrt{14}}{x+2} \tau + \frac{1+x}{x+2}$
(When $z = \frac{2}{7}\sqrt{14}$)
- Term-transformation is $4\sqrt{14}$ and
gauge-transformation is $\frac{1}{x}(\tau - 16)$.
- Applying gt-transformation to $\{P_x(\frac{2}{7}\sqrt{14}), Q_x(\frac{2}{7}\sqrt{14})\}$,
we get

$$\left\{ -\frac{1}{x}(4^{x+2}14^{\frac{1}{2}x}P_x(\frac{2}{7}\sqrt{14}) + 4^{x+1}14^{\frac{1}{2}x+\frac{1}{2}}P_{x+1}(\frac{2}{7}\sqrt{14}), \right.$$

$$\left. -\frac{1}{x}(4^{x+2}14^{\frac{1}{2}x}Q_x(\frac{2}{7}\sqrt{14}) + 4^{x+1}14^{\frac{1}{2}x+\frac{1}{2}}Q_{x+1}(\frac{2}{7}\sqrt{14})) \right\}$$

Example from

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- $\tau^2 - \frac{2}{7} \frac{(2x+3)\sqrt{14}}{x+2} \tau + \frac{1+x}{x+2}$
(When $z = \frac{2}{7}\sqrt{14}$)
- Term-transformation is $4\sqrt{14}$ and
gauge-transformation is $\frac{1}{x}(\tau - 16)$.
- Applying gt-transformation to $\{P_x(\frac{2}{7}\sqrt{14}), Q_x(\frac{2}{7}\sqrt{14})\}$,
we get

$$\left\{ -\frac{1}{x}(4^{x+2}14^{\frac{1}{2}x}P_x(\frac{2}{7}\sqrt{14}) + 4^{x+1}14^{\frac{1}{2}x+\frac{1}{2}}P_{x+1}(\frac{2}{7}\sqrt{14}), \right.$$

$$\left. -\frac{1}{x}(4^{x+2}14^{\frac{1}{2}x}Q_x(\frac{2}{7}\sqrt{14}) + 4^{x+1}14^{\frac{1}{2}x+\frac{1}{2}}Q_{x+1}(\frac{2}{7}\sqrt{14})) \right\}$$

How to add a new base equation

One advantage of *solver* is we can add base equation to it.
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