# Handout WP: Writing Proofs.

## 1. Direct proof of an if-then statement.

Many theorems have the form "if p then q" where p, q are statements. (Note: p, q are often composed of other statements with and's and or's). A direct proof of  $p \Longrightarrow q$  works like this:

Assume: p. To prove: q.

## 2. Proof by contrapositive.

The statement  $p \Longrightarrow q$  is *logically equivalent to* (means: same truth-table) the statement  $\neg q \Longrightarrow \neg p$ . This means we can prove  $p \Longrightarrow q$  like this:

Assume:  $\neg q$ . To prove:  $\neg p$ .

**Remark:** For any statement p you need to be able to compute  $\neg p$  (see the second handout "Quantifiers and Negation" for rules and examples on computing  $\neg p$  for any statement p).

## 3. Proof by contradiction.

One way to prove a statement S is as follows:

Assume:  $\neg S$ .

To prove: a contradiction.

Proving a contradiction means proving something that is obviously wrong (e.g.  $x \neq x$ , or the negation of a given/assumed/proved statement).

## 4. Proving an if-then statement by contradiction.

If S is the statement  $p \Longrightarrow q$  then  $\neg S$  is logically equivalent to  $p \land \neg q$ . So a proof-by-contradiction for  $p \Longrightarrow q$  works like this:

Assume: p and  $\neg q$ . To prove: a contradiction.

## 5. Direct proof for a for-all statement.

If P(x) is a statement involving x, to prove a statement like this  $\forall_{x \in A} P(x)$  you write the following:

Let  $x \in A$ . To prove: P(x).

**Explanation:** When you write "Let  $x \in A$ " then you are telling the reader that x is an element of A, but you are not specifying *which* element of A. That means that x could be any element of A. Once you proved P(x) for such x then P(x) must be true for any element of A.

What not to do: Suppose  $f: S \to T$  and you want to prove  $\forall_{t \in T} P(t)$ . You should start with: "Let  $t \in T$ ". But what if you this do instead:

Some steps ... take t = f(s) ... some steps ... hence P(t).

Then you only proved P(t) for some, but not all, elements of T.

## 6. Direct proof for an exists statement.

If P(x) is a statement involving x, to prove a statement like this  $\exists_{x \in A} P(x)$  you write the following:

Take x :=[write down an expression].

If it is obvious that the expression you wrote down meets the requirements  $x \in A$  and P(x) then the proof is now complete; you have a 1-line proof! (only non-obvious requirements need to be checked). If you wrote a lot of text but not this 1 line, then your proof is still not complete.

#### 7. Proving a for-all statement by contradiction.

If S is the statement  $\forall_{x \in A} P(x)$  then  $\neg S$  is the statement  $\exists_{x \in A} \neg P(x)$ . So if we follow item 3 (i.e. prove S by assuming  $\neg S$  and then proving a contradiction) then the proof would start like this:

Assume: x is an element of A and  $\neg P(x)$ . To prove: a contradiction.

## 8. Proving an exists-statement by contradiction.

If S is the statement  $\exists_{x \in A} P(x)$  then  $\neg S$  is the statement  $\forall_{x \in A} \neg P(x)$ .

Assume:  $\neg P(x)$  for every  $x \in A$ . To prove: a contradiction.

## 9. Proof by cases.

Suppose we need to prove some statement P(x) where x can only have a few possible values. In that case, we can write a separate proof for each possible value.

**Example 1:** Suppose we have to prove P(x) but we know that x can only be u or v or w. Then we write three proofs:

Case 1: x = u. Write a proof for P(u).

Case 2: x = v. Write a proof for P(v).

Case 3: x = w. Write a proof for P(w).

**Example 2:** Suppose we have to prove a statement p but there is some other statement q such that we can easily find a proof for  $q \Longrightarrow p$ . Then we can do the following, we split the proof in two cases:

Case 1: Assume q is true and prove p under that assumption. Case 2: Assume q is false now prove p under that assumption.

These cases combined provide a complete proof for p.

#### 10. Proving an and statement.

To prove  $p \wedge q$  write two separate proofs:

To prove:  $\boldsymbol{p}$ 

To prove: q

11. **Proving an iff statement** (iff = "if and only if").

The statement  $p \iff q$  is logically equivalent to  $(p \implies q) \land (q \implies p)$ . So to prove  $p \iff q$  you have to write two proofs:

To prove:  $p \Longrightarrow q$ 

To prove:  $q \Longrightarrow p$ 

Proofs of "if and only if" statements in math books typically look like:

Assume  $p \ldots$  some math  $\ldots$  hence q. For the converse  $\ldots$ 

Recall that  $q \Longrightarrow p$  is called the **converse** of  $p \Longrightarrow q$ . Math books assume you know that "For the converse" means "To prove:  $q \Longrightarrow p$ ".

#### 12. Proving an or statement.

The statement  $p \lor q$  means that at least one of p or q is true. But which one? That question makes it tricky to give a direct proof of  $p \lor q$ . But you can always replace a statement by a logically equivalent statement. We have several options:

(1)  $p \lor q$  is logically equivalent to  $\neg p \Longrightarrow q$ .

(2)  $p \lor q$  is logically equivalent to  $\neg q \Longrightarrow p$ .

(3)  $\neg(p \lor q)$  is logically equivalent to  $\neg p \land \neg q$ .

That gives us several ways to prove  $p \lor q$ .

Method (1): Assume  $\neg p$ . To prove: q.

Method (2): Assume  $\neg q$ . To prove: p.

Method (3): Assume  $\neg p$  and  $\neg q$ . To prove: a contradiction.

Which method is best? That depends on what you already know. Suppose for instance you see a theorem in the book of the form  $\neg q \Longrightarrow r$ . In that case you want to try method (2) (assume  $\neg q$ ) because then you can use the theorem to conclude r. Hopefully that brings you closer to goal p.

#### 13. Using an or statement.

Suppose you want to prove r, and you are given  $p \lor q$  which means that at least one of p or q is true. But which one? So we write two proofs:

Assume p. To prove r.

Assume q. To prove r.

The two proofs combined show that  $p \lor q$  implies r.

This is the same as "Proof by cases" from item 9; given  $p \lor q$  you distinguish two cases: Case 1: assume p, to prove r. Case 2: assume q, to prove r.

- 14. If you want to use a for-all statement like  $\forall_{x \in A} P(x)$  to prove another statement, often the best strategy is to make a clever choice for one particular element of A, and then use the fact that P is true for that element.
- 15. If you want to use an exists statement like  $\exists_{x \in A} P(x)$  to prove another statement, then you may not choose x. All you know is  $x \in A$  and P(x).