## Frame Design Using Projective Riesz Energy

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Point Distributions Seminar March 17, 2021

This is joint work with Doug Hardin and Ed Saff.

## Frames



#### Frames

A set of vectors  $X = \{x_i\}_{i \in I}$  is a frame for a separable Hilbert space H if there exist A, B > 0 such that for every  $x \in H$ ,

$$A\|x\|^2 \leq \sum_{i\in I} |\langle x, x_i\rangle|^2 \leq B\|x\|^2.$$

The constants A, B are called the frame bounds.

When A = B, X is called a tight frame, which generalizes the concept of an orthonormal basis in the sense that we have the following recovery formula

$$x = \frac{1}{A} \sum_{i \in I} \langle x, x_i \rangle x_i, \quad \forall x \in H.$$



robustness to erasures stability in reconstruction

 $\begin{array}{l} \mbox{signal processing} \\ \rightarrow & \mbox{coding theory} \\ \mbox{quantum inf. theory} \end{array}$ 

R. B. Holmes and V. I. Paulsen. Optimal frames for erasures. Linear Algebra Appl. 377 (2004): 31-51
 J. M. Renes, R. Blume-Kohout, A. J. Scott, and C. M. Caves. Symmetric informationally complete quantum measurements. J. Math. Phys., 45(6):2171–2180, 2004.

[3] T. Strohmer and R. W. Heath, Jr. Grassmannian frames with applications to coding and communication. Appl. Comput. Harmon. Anal., 14(3):257–275, 2003.

S(d, N) = the set of all unit norm frames (with N frame vectors) for  $\mathbb{H}^d$ = the collection of all N-point configurations on  $\mathbb{S}^{d-1}$ 

We will focus on finding unit norm finite frames  $X = [x_1, \cdots, x_N]$  that are

• nearly tight: 
$$XX^* \approx \frac{N}{d}I_d$$

► well-separated: coherence  $\xi(X) = \max_{i \neq i} |\langle x_i, x_j \rangle|$  is small.

 $\min_{X\in\mathcal{S}(d,N)}\xi(X) \text{ is the best line-packing problem [4]}.$ 

[4] J. H. Conway, R. H. Hardin, and N. J. A. Sloane. Packing lines, planes, etc.: Packings in Grassmannian spaces. Exp. Math. 5.2 (1996)

## Desirable properties of frames

A frame  $X = \{x_i\}_{i=1}^N$  is an equiangular tight frame (ETF) if

► X is tight, and ►  $\frac{|\langle x_i, x_j \rangle|}{||x_i||_2 ||x_j||_2}$  is constant for any pair  $i \neq j$ .

ETF is exactly tight and best-separated, but it does not always exist.

[5] M. Fickus and D. G. Mixon. Tables of the existence of equiangular tight frames.[6] S. Waldron. An introduction to finite tight frames. Boston: Birkhäuser, 2018.

## Previous work

Minimizing frame potential [7]

$$\min_{X \in \mathcal{S}(d,N)} \sum_{i \neq j} |\langle x_i, x_j \rangle|^2$$
(1)

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Minimizing frame potential [7]

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The minimizing configurations of (1) are exactly unit norm tight frames.



[7] J. Benedetto and M. Fickus. Finite normalized tight frames. Adv. Comput. Math. 2003

$$\min_{X \in \mathcal{S}(d,N)} \sum_{i \neq j} |\langle x_i, x_j \rangle|^{\rho}$$
(2)

- Bigger *p* value promotes separation, since (2) approaches coherence minimizing as *p* → ∞.
- If p > 2, ETF (when exists) is the minimizer.

[8] M. Ehler and K. Okoudjou. Minimization of the Probabilistic *p*-frame Potential, J. Statist. Plann. Inference, 142 (2012)
[9] D. Bilyk, D., A. Glazyrin, R. Matzke, J. Park, and O. Vlasiuk. Optimal measures for p-frame energies on spheres. (2019)
[10] X. Chen, V. Gonzalez, E. Goodman, S. Kang, and K. Okoudjou. Universal optimal configurations for the p-frame potentials. Adv. Comput. Math (2020)

Solving 
$$\min_{X \in S(d,N)} \sum_{i \neq j} \frac{1}{\left(\sqrt{2 - 2|\langle x_i, x_j \rangle|^2}\right)^s}$$
 produces nearly tight and well-separated (low coherence) frames.

$$\sqrt{2-2|\langle x,y\rangle|^2} = \|xx^*-yy^*\|_F$$

Set up

$$\min_{X \in \mathcal{S}(d,N)} \sum_{i \neq j} f\left(\sqrt{2 - 2|\langle x_i, x_j \rangle|^2}\right)$$

where *f* is decreasing and convex on  $(0, \sqrt{2}]$ .

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$$\begin{split} f(t) &= \frac{1}{t^s}, s > 0 \\ \text{decreasing, strictly convex} \\ \hline f(t) &= -\log(t) \\ \text{decreasing, strictly convex} \\ \hline f(t) &= -\log(t) \\ \text{decreasing, strictly convex} \\ \hline f(t) &= \left(\frac{2-t^2}{2}\right)^{p/2}, p > 0 \\ \text{decreasing, } f \text{ is not convex} \\ \hline f(t) &= \left(\frac{x_i, x_j}{2}\right)^{p/2} \\ \hline f(t) &= \left(\frac{2-t^2}{2}\right)^{p/2}, p > 0 \\ \text{decreasing, } f \text{ is not convex} \\ \hline f(t) &= \left(\frac{x_i, x_j}{2}\right)^{p/2} \\ \hline f(t) &= \left(\frac{x_i, x_j}{2}\right)^{p/2}, p > 0 \\ \hline f(t) &= \left(\frac{x_i,$$

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 $c \rightarrow N$ 

S(d, N) = the collection of all *N*-point configurations on  $\mathbb{S}^{d-1}$  $\mathcal{P}(d, N)$  = the collection of all *N*-point configurations on  $\mathbb{P}^{d-1}$ 

$$X = \{x_i\}_{i=1}^{i}$$

$$P = \{p_i = x_i x_i^*\}_{i=1}^{N}$$

$$\min_{X \in \mathcal{S}(d,N)} \sum_{i \neq j} f\left(\sqrt{2 - 2|\langle x_i, x_j \rangle|^2}\right) = \min_{P \in \mathcal{P}(d,N)} \sum_{i \neq j} f(||p_i - p_j||)$$

# Results on $\mathbb{R}^2$

#### Theorem 1 (C. Hardin, Saff, 2020)

If *f* is a non-increasing convex function, then equally distributed points on half circle is an optimal configuration of

$$\min_{X \in \mathcal{S}(d,N)} \sum_{i \neq j} f\left(\sqrt{2 - 2|\langle x_i, x_j \rangle|^2}\right)$$

If, in addition, *f* is strictly convex, then no other *N*-point configuration is optimal.



## Results related to ETF

### Theorem 2 (C., Hardin, Saff, 2020)

Let *f* be decreasing and strictly convex, then ETF (when exists) is the unique minimizer of

$$\min_{X\in\mathcal{S}(d,N)}\sum_{i\neq j}f\left(\sqrt{2-2|\langle x_i,x_j\rangle|^2}\right).$$

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#### Theorem 2 (C., Hardin, Saff, 2020)

#### This is a stronger theorem:

Let g be decreasing and strictly convex, then ETF (when exists) is the unique minimizer of

$$\min_{X\in\mathcal{S}(d,N)}\sum_{i\neq j}g\left(2-2|\langle x_i,x_j\rangle|^2\right).$$

 $|\langle x, y \rangle|^p = g(2-2|\langle x, y \rangle|^2)$ , where  $g(t) = \left(\frac{2-t}{2}\right)^{p/2}$  is decreasing and strictly convex when p > 2.

## Results related to ETF

Proof of Theorem 2: Let  $p_i = x_i x_i^*$ .

$$J := \sum_{i \neq j} \|p_i - p_j\|^2 = \sum_{i=1}^{N} \sum_{j \neq i} (2 - 2\langle p_i, p_j \rangle)$$
  
=  $\sum_{i=1}^{N} \left( 2(N-1) - 2 \sum_{j=1}^{N} \langle p_i, p_j \rangle + 2 \right)$   
=  $2N^2 - 2 \sum_{i,j=1}^{N} |\langle x_i, x_j \rangle|^2 \le 2N^2 - 2N^2/d,$ 

$$\begin{split} E(X) &= \frac{N(N-1)}{1} \cdot \frac{1}{N(N-1)} \sum_{i \neq j} g(\|p_i - p_j\|^2) \\ &\geq \frac{N(N-1)}{1} g\left( \sum_{i \neq j} \frac{1}{N(N-1)} \|p_i - p_j\|^2 \right) = N(N-1) g\left( \frac{1}{N(N-1)} J \right) \\ &\geq N(N-1) g\left( \frac{2N^2 - 2N^2/d}{N(N-1)} \right) = N(N-1) g\left( \frac{2N(1-1/d)}{(N-1)} \right). \end{split}$$

## Projective Riesz kernel

$$\frac{1}{\left(\sqrt{2-2|\langle x_i, x_j\rangle|^2}\right)^s} = \frac{1}{\|p_i - p_j\|^s}$$

#### Theorem 3 (Separation)

Let  $s > \alpha = \dim(\mathbb{P}^{d-1})$ . If  $X_N$  is an *N*-point minimizing configuration of

$$\min_{X \in \mathcal{S}(d,N)} \sum_{i \neq j} \frac{1}{\left(\sqrt{2 - 2|\langle x_i, x_j \rangle|^2}\right)^s},\tag{3}$$

then

$$\xi(X_N) \leq \sqrt{1 - \frac{C^2}{2}N^{-2/\alpha}},$$

where the constant C depends on s and d.

Proof uses a result from [11]

Suppose  $A \subset \mathbb{R}^m$  is compact and supports an upper  $\alpha$ -regular measure  $\mu$ . Let  $s > \alpha$ ,  $N \ge 2$  be fixed. If  $P_N$  is an N-point minimizing configuration on A for the s-Riesz energy minimizing problem, then

$$\delta(P_N) \geq C_1 N^{-\frac{1}{\alpha}},$$

where 
$$C_1 = \left(rac{\mu(\mathcal{A})}{C_{\mathcal{A}}}(1-rac{lpha}{s})
ight)^{1/lpha} \left(rac{lpha}{s}
ight)^{rac{1}{s}}.$$

$$\delta^2(P_N) = \min_{i \neq j} \|p_i - p_j\|^2 = \min_{i \neq j} (2 - 2|\langle x_i, x_j \rangle|^2) = 2 - 2\xi^2(X_N)$$

[11] D. P. Hardin, E. B. Saff, and J. T. Whitehouse. Quasi-uniformity of minimal weighted energy points on compact metric spaces. J. Complexity 28.2 (2012)

#### Theorem 4 (Nearly tight)

Let  $s \ge 0$ . If  $X_N$  is an *N*-point minimizing configuration of

$$\min_{X\in\mathcal{S}(d,N)}\sum_{i\neq j}\frac{1}{\left(\sqrt{2-2|\langle x_i,x_j\rangle|^2}\right)^s},$$

then

$$\lim_{N\to\infty}\frac{1}{N}X_NX_N^*=\frac{1}{d}I_d.$$
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#### Theorem 4 (Nearly tight)

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$$\lim_{N \to \infty} \frac{1}{N} X_N X_N^* = \frac{1}{d} I_d.$$
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The counting measure converges weak\* to the uniform measure.

Theorem 4 says that the optimal configurations are nearly tight asymptotically. A more desirable result would be

$$\left\|X_N X_N^* - \frac{N}{d}I_d\right\|_F = \mathcal{O}(N^{-q}), \quad q > 0.$$
<sup>(5)</sup>

Numerical experiments suggest this holds for small values of *s*.

## Numerical experiment: tightness



## Numerical experiment: separation



A method to generate well separated tight frame:

- 1. Generate a random frame  $X \in S(d, N)$ .
- 2. Using *X* as an initial configuration, use an optimization algorithm (such as gradient descent) to find a local minimizer *Y* for

$$\min_{X \in \mathcal{S}(d,N)} \sum_{i \neq j} \frac{1}{\left(\sqrt{2 - 2|\langle x_i, x_j \rangle|^2}\right)^s}$$
(3)

for some s > d - 1. *Y* is expected to be well-separated and nearly tight according to our theorems.

3. Solve  $\min_{X \in S(d,N)} \sum_{i \neq j} |\langle x_i, x_j \rangle|^2$  using *Y* as the initial configuration.

## Numerical experiment: a proposal



## Thank you

This work is sponsored by NSF