

# Semidefinite Programming Bounds for the Average Kissing Number

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joint work with

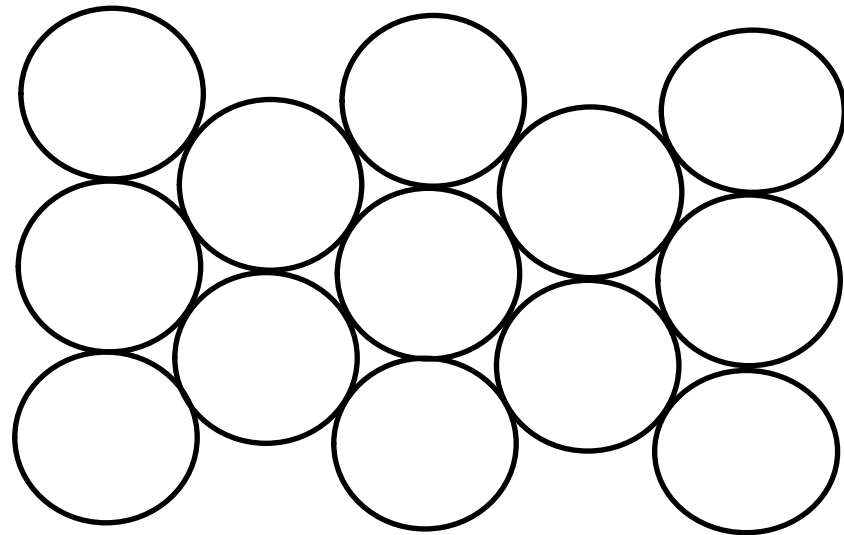


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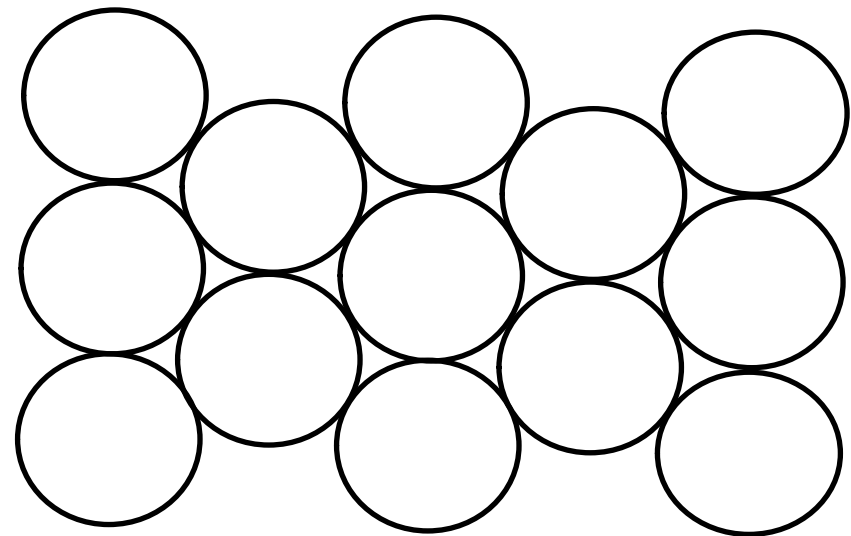


Packing of unit disks by hexagonal lattice

Each disk touches 6 disks

⇒ Average kissing number of the hexagonal lattice is 6

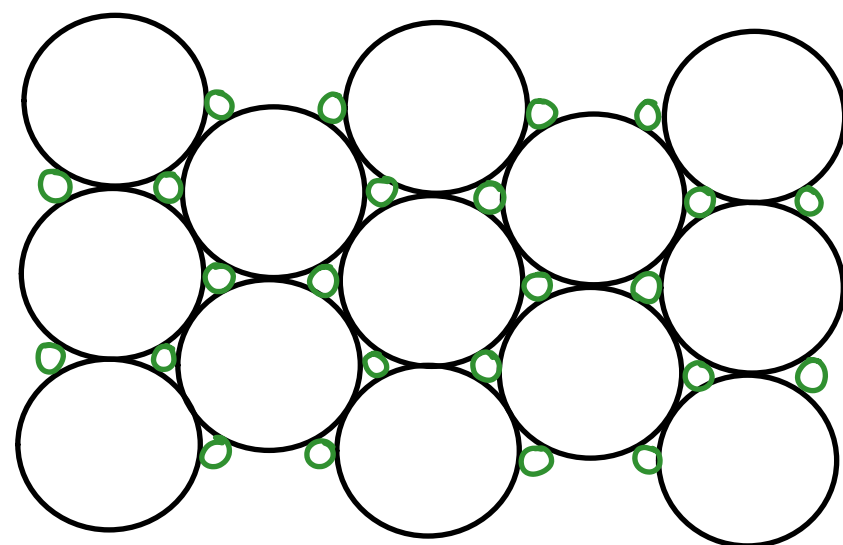
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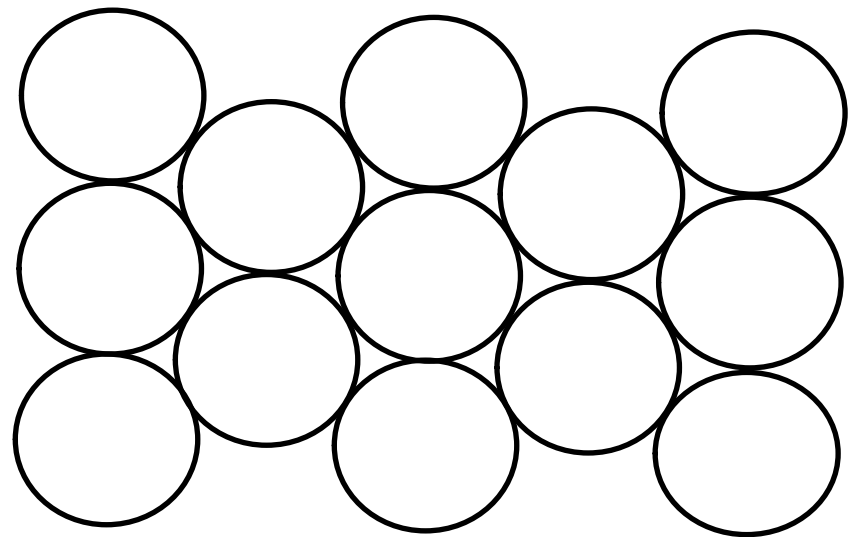
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- any black disk touches 12 disks
  - any green disk touches 3 disks
  - the average is still 6

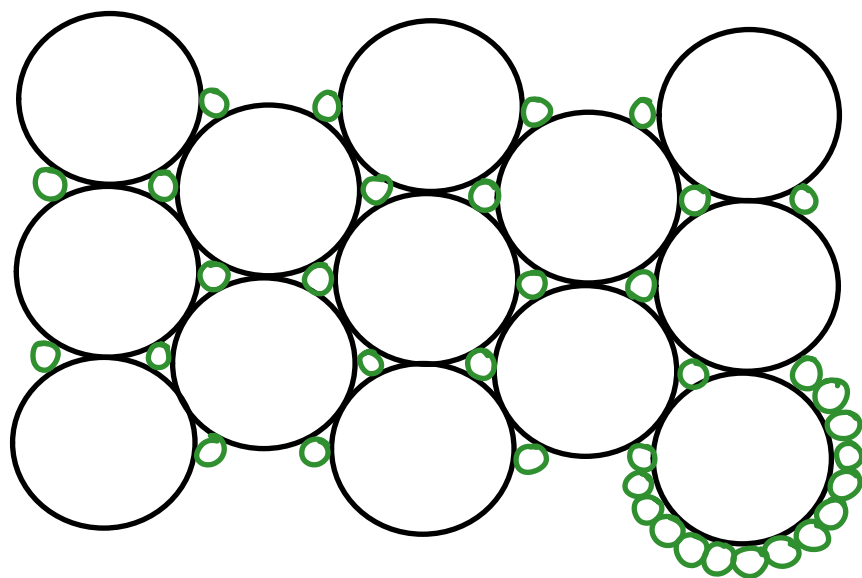
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Can we increase the average kissing number by adding more smaller disks?



# Average Kissing Number

Packing of balls in  $\mathbb{R}^n$ : finite set of interior-disjoint closed balls

Contact graphs of a packing  $\mathcal{P}$ : graph with vertex set  $\mathcal{P}$  in which two balls  $X$  and  $Y$  are adjacent if  $X \cap Y \neq \emptyset$

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Average kissing number in  $\mathbb{R}^n$ :

$$k_n = \sup \{ \bar{\delta}(G) : G \text{ is the contact graph of a packing of balls in } \mathbb{R}^n \}$$

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Koebe-Andrew-Thurston: Contact graphs for packings of disks on the plane are simple planar graphs  $\Rightarrow k_2 = 6$

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 $G(\mathcal{P}, E)$  contact graph of  $\mathcal{P}$ ,  $\tau_n$  kissing number

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$$\bar{d}(G) = \frac{2|E|}{|\mathcal{P}|} \leq 2\tau_n \quad \Rightarrow \quad k_n \leq 2\tau_n$$

# Upper Bound

First nontrivial upper bound by Kuperberg & Schramm  
 $\Rightarrow k_3 \leq 8 + 4\sqrt{3} = 14.928\dots$

Glazyrin refines this approach :  $k_3 \leq 13.955$   
and extends it to higher dimension  
beats  $2\tau_n$  upper bound for  $n=4,5$

Our goal: Refine Glazyrin's approach by using  
Semidefinite programming.

# Notations

Euclidean inner product :  $x \cdot y = \sum_{i=1}^n x_i y_i$ ,  $x, y \in \mathbb{R}^n$

( $n-1$ )dim unit sphere  $S^{n-1} = \{x \in \mathbb{R}^n : \|x\|=1\}$

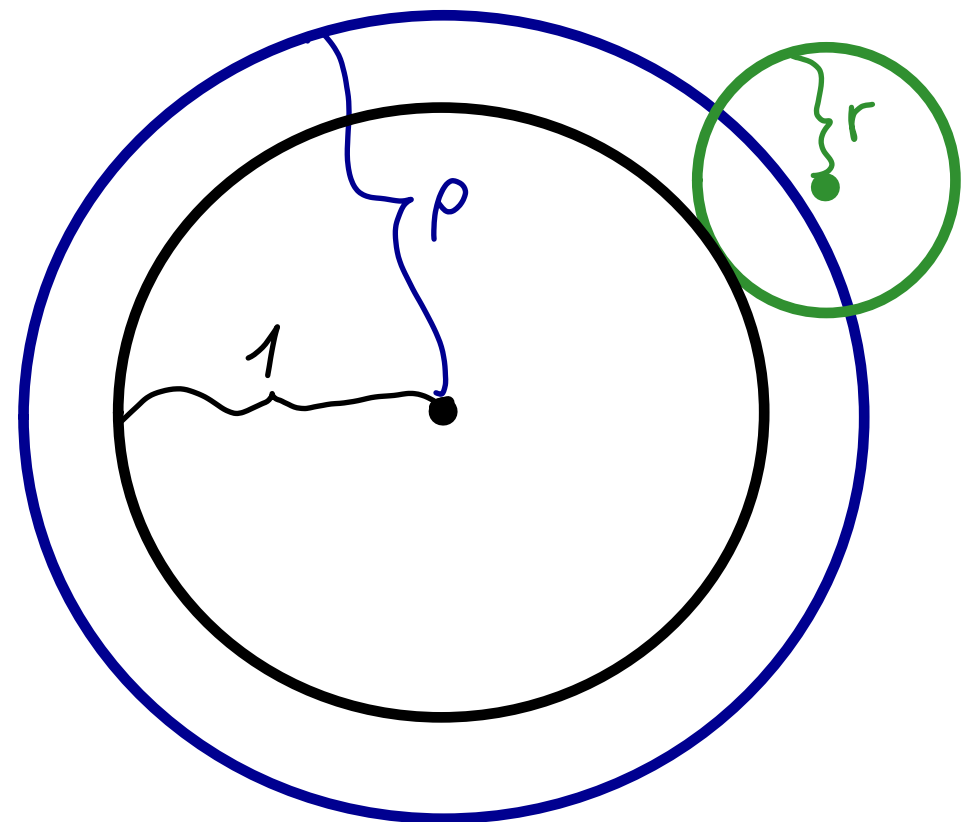
distance between  $x, y \in S^{n-1}$  :  $\arccos x \cdot y$

surface measure of ( $n-1$ )dim sphere of radius  $\rho$  :  $\omega_\rho$  ( $\omega = \omega_1$ )

spherical cap in  $S^{n-1}$  of center  $x \in S^{n-1}$  and radius  $\alpha$  :  $\{y \in S^{n-1} : x \cdot y \geq \cos \alpha\}$

normalized area of this cap :  $\frac{\omega(S^{n-2})}{\omega(S^{n-1})} \int_{\cos \alpha}^1 (1-u^2)^{\frac{n-3}{2}} du$

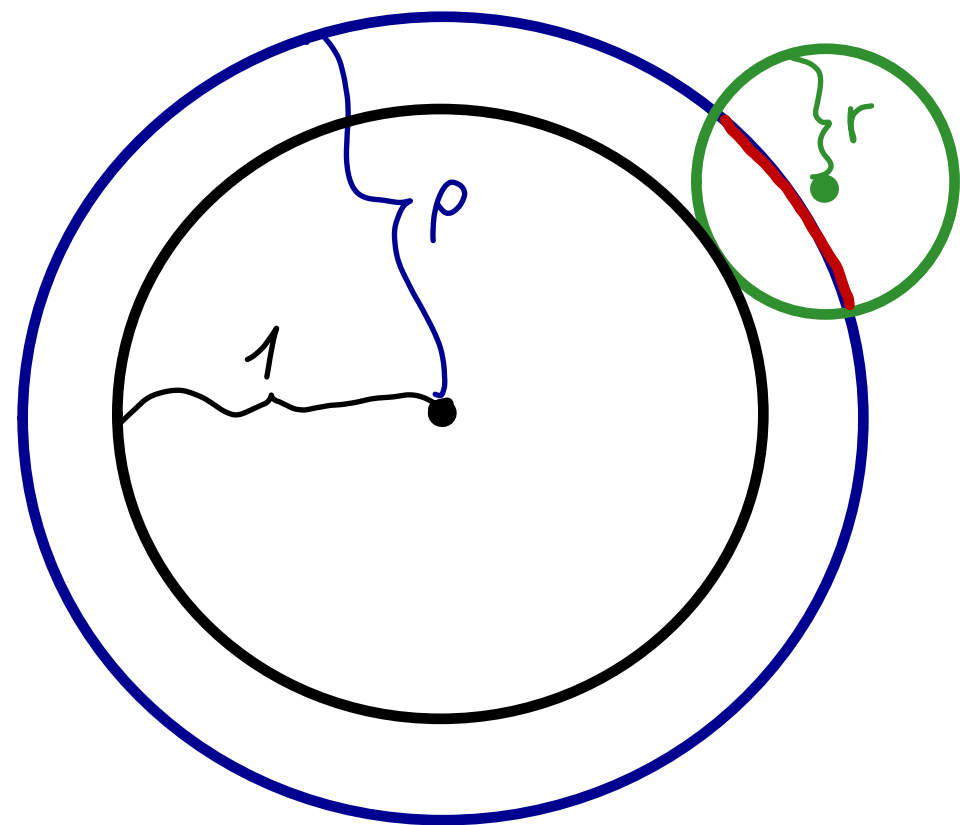
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$B_r \cap \rho S^{n-1}$  is either empty or a spherical cap.

Normalized area of this spherical cap:

$$A_{n,\rho}(r) = \frac{\omega_\rho(B_r \cap \rho S^{n-1})}{\omega_\rho(\rho S^{n-1})}$$

as a function of  $r$  is monotonically increasing.

# Glazyrin's Upper Bound

Lemma: If  $n \geq 3, p > 1, r > 0$  then  $A_{n,p}(r) + A_{n,p}(\frac{1}{r}) \geq 2A_{n,p}(1)$



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Fix  $\rho > 1$  and consider unit ball at origin.

Any configuration of pairwise interior-disjoint balls tangent to central unit ball covers a fraction of  $\rho S^{n-1}$  centered at origin.

$\text{dens}_n$ : sup of covered fraction over all configurations

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$$\sum_{\{X,Y\} \in E} A_{n,\rho}\left(\frac{r(X)}{r(Y)}\right) + A_{n,\rho}\left(\frac{r(Y)}{r(X)}\right) \geq 2 A_{n,\rho}(1) |E|$$

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For  $\rho < 3: A_{n,\rho}(1) > 0$

$$\leq |\mathcal{P}| \text{dens}_n(\rho)$$

$$\Rightarrow \frac{2|E|}{|\mathcal{P}|} \leq \frac{\text{dens}_n(\rho)}{A_{n,\rho}(1)}$$

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$$K_n \leq \frac{\text{dens}_n(\rho)}{A_{n,\rho}(1)}$$

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Theorem: If  $n \geq 3$  and  $1 < p < 3$ , then  $k_n \leq \frac{\text{dens}_n(p)}{A_{n,p}(1)}$



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For  $\text{dens}_n(\rho) \leq 1$ ,  $\rho = \sqrt{3}$  :

$k_3 \leq 14.928\dots$	(Kuperberg & Schramm)
$k_4 \leq 34.680\dots$	} (Glazyrin)
$k_5 \leq 77.756\dots$	

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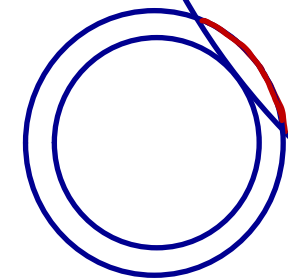
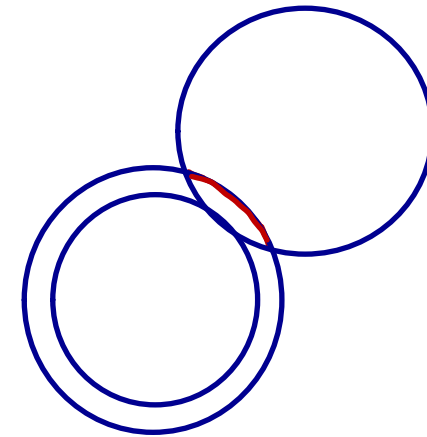
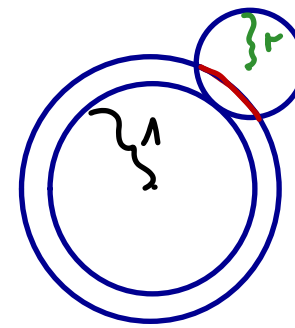
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Glazyrin improved upper bound for  $\text{dens}_3(\sqrt{3}) \Rightarrow k_3 \leq 13.955$ .

# Refining Glazyrin's approach using Semidefinite Programming

•  $A_{n,p}$  increasing in  $r$  and has a limit  $A_{n,p}(\infty)$

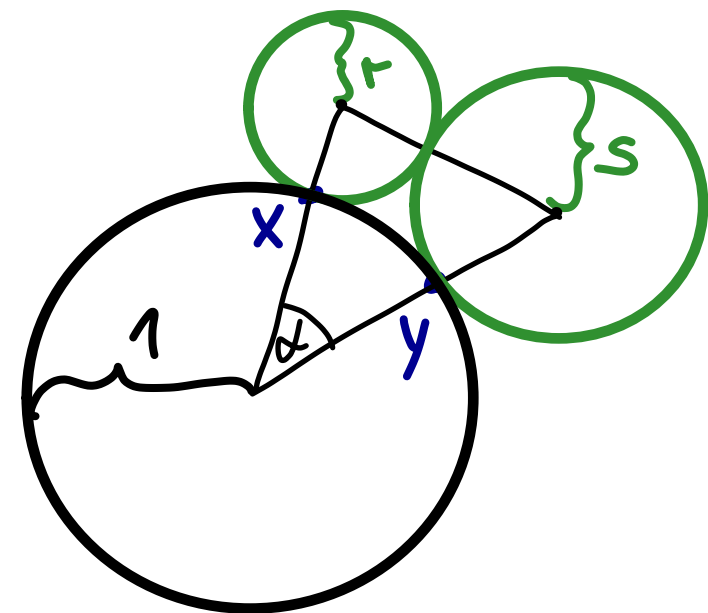
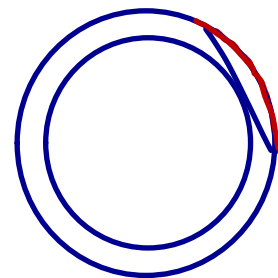


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$$\cdot (r+s)^2 = (1+s)^2 + (1+r)^2 - 2(1+r)(1+s) \cos \alpha$$

$$\Rightarrow x \cdot y \leq \frac{1+r+s-rs}{1+r+s+rs} =: ip(r,s)$$



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- $(r+s)^2 = (1+s)^2 + (1+r)^2 - 2(1+r)(1+s)\cos\alpha \Rightarrow x \cdot y \leq \frac{1+r+s-rs}{1+r+s+rs} = \text{ip}(r,s)$
- **kernel**: real-valued square integrable function on  $V \times V$  ( $V$  measure space)
- If  $f: V \rightarrow \mathbb{R}$  is square integrable  $\Rightarrow f \otimes f^*$  is the kernel of  $(x,y) \mapsto f(x)f(y)$
- If  $F: [0,1] \rightarrow \mathbb{R}$  is a kernel,  $U \subseteq [0,1]$  finite, then  $(F(u,v))_{u,v \in U}$  is a **principal submatrix** of  $F$
- $P_k^n$ : Jacobi polynomial degree  $k$ ,  $\alpha = \beta = \frac{(n-3)}{2}$ ,  $P_k^n(1) = 1$

# Semidefinite Programming Bound

Theorem: Let  $n \geq 3$ ,  $1 < p < 3$ ,  $R$  s.t.  $R > \frac{p-1}{2}$ ,  $r: [0,1] \rightarrow [\frac{p-1}{2}, R]$  increasing bijection

$a: [0,1] \rightarrow \mathbb{R}$  with  $a(u) \geq A_{n,p}(r(u))^{1/2}$  for all  $u \in [0,1]$   
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Fix  $d \geq 0$ , for  $k=0, \dots, d$  let  $F_k: [0,1]^2 \rightarrow \mathbb{R}$  be a kernel

$$f(t, u, v) = \sum_{k=0}^d F_k(u, v) P_k^n(t) \quad \text{for } t \in [-1, 1], u, v \in [0, 1]$$

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- 1) every principal submatrix of  $F_0 - a \otimes a^* \succ 0$ ,
  - 2) every principal submatrix of  $F_k \succ 0$  for all  $k=0, \dots, d$ , and
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- If
- 1) every principal submatrix of  $F_0 - a \otimes a^* \not\geq 0$ ,
  - 2) every principal submatrix of  $F_k \geq 0$  for all  $k=0, \dots, d$ , and
  - 3)  $f(t, u, v) \leq 0$  whenever  $-1 \leq t \leq \text{ip}(r(u), r(v))$

then

$$\text{dens}_n(p) \leq \max \{ f(1, u, u) : u \in [0, 1] \}$$

# Semidefinite Programming Bound

Proof:  $\Delta$  normalized area of  $\rho S^{n-1}$  covered by the configuration of  $\mathcal{P}$

Assume  $\Delta > 0$ , assume each ball has radius  $\geq \frac{\rho-1}{2}$

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$B$  represented by  $(x, u)$ ,  $x = B \cap \text{central } B_1$

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By 1) & 2) we have

$$0 \leq \sum_{(x,u),(y,v) \in I} f(x \cdot y, u, v) a(u) a(v) - a(u)^2 a(v)^2$$

# Semidefinite Programming Bound

Proof:  $\Delta$  normalized area of  $\rho S^{n-1}$  covered by the configuration of  $\mathcal{P}$

Assume  $\Delta > 0$ , assume each ball has radius  $\geq \frac{\rho-1}{2}$

Let  $B \in \mathcal{P}$ , if its radius is in  $[\frac{\rho-1}{2}, R]$ , let  $u \in [0, 1]$  s.th.  $r(u) = \text{radius}$ , otherwise  $u=1$ .

$B$  represented by  $(x, u)$ ,  $x = B \cap \text{central } B_1$

$I \subseteq S^{n-1} \times [0, 1]$  set of pairs representing each  $B \in \mathcal{P}$

By 1) & 2) we have

here we use 3)  $f(t, u, v) \leq 0$  whenever  $-1 \leq t \leq 1$

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$$\Rightarrow \sum_{(x, u) \in I} a(u)^2 \leq \max \{ f(1, u, u) : u \in [0, 1] \}$$

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• One way: Fix  $N > 0$  and functions  $p_0, \dots, p_N: [0, 1] \rightarrow \mathbb{R}$ .

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$\Rightarrow$  We can rephrase SDP. Different choice of  $p_0, \dots, p_N$  gives different SDP

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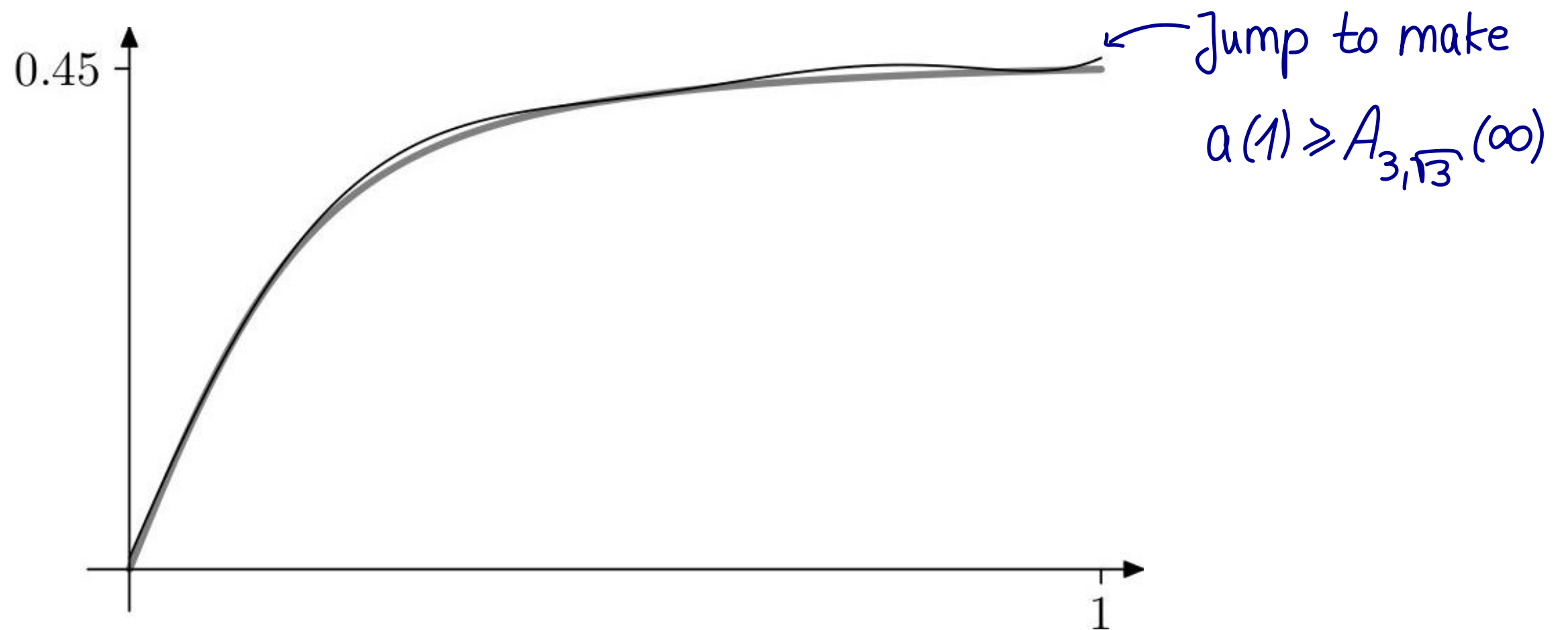
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Interplay between choice of  $p_i$  and quality of approximation  $a$  of  $u \mapsto A_{n,p}(r(u))^{1/2}$

We used two approaches: 1) Set  $p_i$  as step functions

2) Take functions  $p_i$  to be polynomials

# Approximation of $u \mapsto A_{n,p}(r(u))^{1/2}$



A polynomial of degree 6 (in black) that approximates  $u \mapsto A_{3, \sqrt{3}}(r(u))^{1/2}$  (in gray) from above; here  $R=30$

# Results

DIMENSION	LOWER BOUND	PREVIOUS UPPER BOUND	NEW UPPER BOUND
3	12.612	13.955	13.606
4	24	34.681	27.439
5	40	77.757	64.022
6	72	156	121.105
7	126	268	223.144
8	240	480	408.386
9	272	726	722.629

} SDP with polynomials  
 $d=10, N=6, 8, R=30$

} Step function approach  
 $p=2, N=30, R \approx 184.25$

Rigorous verification by Julia library of D., de Laat, Moustrou



Thank you!

# Step function approach

Fix  $R > \frac{p-1}{2}$ ,  $r: [0,1] \rightarrow [\frac{p-1}{2}, R]$ ,  $r(u) = (R - \frac{p-1}{2})u + \frac{p-1}{2}$

Fix  $N > 0$  and points  $0 = s_0 < s_1 < \dots < s_N < s_{N+1} = 1$ .

Let  $S_i = [s_i, s_{i+1}]$  for  $i = 0, \dots, N-1$ ,  $S_N = [s_N, s_{N+1}]$

$p_i$  is 1 on  $S_i$ , 0 otherwise.

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$$a(u) = \begin{cases} A_{n,p}(r(s_{i+1}))^{1/2} & \text{if } u \in S_i \text{ for some } i < N \\ A_{n,p}(\infty)^{1/2} & \text{if } u \in S_N \end{cases}$$

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$\min \max \{f_{ij}(1) : i=0, \dots, N\}$

$$f_{ij}(t) = \sum_{k=0}^d A_{k,ij} P_k^n(t),$$

$$f_{ij}(t) \leq 0 \quad \text{whenever } -1 \leq t \leq \text{ip}(r(s_i), r(s_j))$$

$(A_{0,ij} - \alpha_i \alpha_j)_{i,j=0}^N$  is positive semidefinite,

$A_k \in \mathbb{R}^{(N+1) \times (N+1)}$  is positive semidefinite for  $k=0, \dots, d$ .

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← By using Sampling

# Polynomial approach

Fix  $N > 0$  and let  $p_i(u) = u^i$  for  $i = 0, \dots, N$

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Solve LP to get to obtain  $\alpha_0, \dots, \alpha_N$  (coefficients of  $p_0, \dots, p_N$ )

Fix  $\varepsilon > 0$ , and finite sample  $S$ :  $\alpha_0 p_0(u) + \dots + \alpha_N p_N(u) \geq A_{n,p} (r(u))^{1/2} + \varepsilon$  for sample points  $u$

$$\alpha_0 p_0(1) + \dots + \alpha_N p_N(1) \geq A_{n,p} (\infty)^{1/2}$$

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We get approximation  $a$  and

use sum of squares to check  $f(t, u, v) \leq 0$  for  $\{(u, v, t) : s_i(u, v, t) \geq 0 \ i = 1, \dots, 4\}$

for some polynomials  $s_1, \dots, s_4$



# Polynomial approach

min  $\lambda$

$$f(t, u, v) = \sum_{k=0}^d \sum_{i,j=0}^N A_{k,ij} p_i(u) p_j(v) P_k^n(t)$$

$$f = -s_1 q_1 - s_2 q_2 - s_3 q_3 - s_4 q_4 - q_5$$

$$\lambda - f(1, u, u) = l_1(u) + (1-u)u l_2(u)$$

$q_1, \dots, q_5$  sum of squares polynomials in  $u, v, t$

$l_1, l_2$  sum of squares polynomials in  $u$

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