## Sign Uncertainty Principle $\approx$ Tug-of-war


$\triangle$ Sign uncertainty principle
$\square$ LP bounds for sphere packing
Why these problems are so challenging and surprising

We say $f: \mathbb{R}^{d} \rightarrow \mathbb{R}$ is eventually nonnegative (E.NN.) if $f(x) \geq 0, \quad$ for all sufficiently large $|x|$. and we define

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$$
\mathcal{A}_{+}(d)=\left\{\begin{array}{l}
f, \widehat{f} \in L^{1}\left(\mathbb{R}^{d}\right) \text { both even and real-valued } \\
\widehat{f}(0) \leq 0, f(0) \leq 0 \\
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Thm (Bourgain, Clozel, Kahane, 2010)

$$
\frac{1}{\sqrt{2 \pi e}} \leq \frac{\mathbb{A}_{+}(d)}{\sqrt{d}} \leq \frac{1+o(1)}{\sqrt{2 \pi}} .
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## Thm (Cohn, Gonçalves, 2018)

$$
\begin{gathered}
\frac{1}{\sqrt{2 \pi e}} \leq \frac{\mathbb{A}_{-}(d)}{\sqrt{d}} \leq(1+o(1))\left(\frac{1}{\sqrt{2 \pi}}-0.079 \ldots\right) . \\
\mathbb{A}_{-}(d)=\inf \{r(f): f \text { radial, } \widehat{f}=-f, f(0)=0, f \text { E.NN. }\}
\end{gathered}
$$

$\mathbb{A}_{ \pm}(d) \gtrsim \sqrt{d} \rightsquigarrow \pm 1$ Uncertainty Principles

Thm (Gonçalves, Oliveira e Silva, Steinerberger, 2016; Cohn, Gonçalves, 2018)

Existence of Optimal: $\exists f \in \mathcal{A}_{ \pm}(d)$ such that

$$
r(f)=\mathbb{A}_{ \pm}(d)
$$

we can assume $f$ radial, $\widehat{f}= \pm f$ and $f(0)=0$.
Multiple Roots: $f(|x|)$ has infinitely many double roots for $|x|>r(f)$.

## Sphere Packing Problem

What is the most dense arrangement of non-overlapping equal spheres in $\mathbb{R}^{d} ? \rightsquigarrow \Delta(d)=$ largest density

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(Thue, 1910)

honeycomb ~ 91\%
(Hales, 1998)


Cannonball Packing
~ 70\%


## Thm (Cohn, Elkies, 2003) [Linear Programming Bounds]

Let

$$
\mathcal{A}_{L P}(d)=\left\{\begin{array}{l}
f, \widehat{f} \in L^{1}\left(\mathbb{R}^{d}\right) \text { radial and real-valued } \\
f(0)=\widehat{f}(0)=1 \\
f \text { E.NN. and } \widehat{f} \geq 0 .
\end{array}\right.
$$

$$
\text { and } \mathbb{A}_{L P}(d)=\inf _{f \in \mathcal{A}_{L P}(d)} r(f) \text {. Then }
$$

$$
\Delta(d) \leq \operatorname{vol}\left(\frac{1}{2} B^{d}\right) \mathbb{A}_{L P}(d)^{d}
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and $\mathbb{A}_{L P}(d)=\inf _{f \in \mathcal{A}_{L P}(d)} r(f)$. Then

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Viazovska showed $\mathbb{A}_{L P}(8)=\sqrt{2}$ and $\mathbb{A}_{L P}(24)=\sqrt{4}$ with a construction using modular forms.

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- Cohn \& G. (2018) showed that

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- Recent numerical evidence by Afkhami-Jeddi, Cohn et al in connection with modular bootstraps for CFTs indicates

$$
c=\frac{1}{\pi}=.31 \ldots
$$

| $d$ | Best Packing | $\mathbb{A}_{L P}(d)$ | $\mathbb{A}_{-}(d)$ | $\mathbb{A}_{+}(d)$ |
| :---: | :---: | :---: | :---: | :---: |
| 1 | $\mathbb{Z}$ | 1 | 1 | $? ?$ surprise |
| 2 | Honeycomb | $?=(4 / 3)^{\frac{1}{4}}$ | $?=(4 / 3)^{\frac{1}{4}}$ | $?$ |
| 8 | $E 8$ | $\sqrt{2}$ | $\sqrt{2}$ | $?$ |
| 12 | $?$ | $?$ | $?$ | $\sqrt{2}$ |
| 24 | Leech | $\sqrt{4}$ | $\sqrt{4}$ | $?$ |

## New Sign Uncertainty Principles <br> (G., Oliveira e Silva, Ramos - arXiv March 2020)

- Spherical Harmonics and Jacobi Polynomials ( $\rightsquigarrow$ Spherical Designs and Quadrature).


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- Spherical Harmonics and Jacobi Polynomials ( $\rightsquigarrow$ Spherical Designs and Quadrature).
- Fourier Series, Bessel-Dini Series.
- Discrete Fourier and Hankel Transf..
- Functions of the Hamming cube (Complexity of Boolean Functions).
- Hankel Transf., Hilbert Transf. and other Smooth Conv. Kernels.

For $f: \mathbb{Z}_{2 q+1} \rightarrow \mathbb{C}$ we define the DFT

$$
\widehat{f}(n)=\left.\ell \sum_{m=-q}^{q} f(n) e^{-2 \pi i x \ell m}\right|_{x=\ell n}\left(\ell=\frac{1}{\sqrt{2 q+1}}\right)
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$$

## Thm

Let
$\mathcal{A}_{ \pm}^{\text {disc }}[q]=\left\{\begin{array}{l}f, \widehat{f}: \mathbb{Z}_{2 q+1} \rightarrow \mathbb{R} \text { both even and real-valued } \\ \widehat{f}(0) \leq 0, \pm f(0) \leq 0 .\end{array}\right.$
Then

$$
\mathbb{A}_{ \pm}^{\text {disc }}[q]:=\min \{\sqrt{k(f) k( \pm \widehat{f})}\} \gtrsim \sqrt{2 q+1},
$$

where $k(f)=\min \{k>0: f(n) \geq 0$ if $n \geq k\}$.

## Numerical evidence that $\frac{\mathbb{A}_{l}^{\text {disc }}[q]}{\sqrt{2 q+1}} \rightarrow \mathbb{A}_{-}(1)$



## Numerical evidence for




The function $q \mapsto k=\mathbb{A}_{ \pm}^{\text {disc }}[q]$ is a stairway: $k \mapsto q_{ \pm}^{\text {jump }}(k)$


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It turns out that (numerically):

$$
q_{-}^{j u m p}(k)=\left\lfloor\frac{k^{2}-2 k+2}{2}\right\rfloor_{k \geq 4}=5,8,13,18,25,32, \ldots
$$

$$
q_{+}^{\text {jump }}(k) \approx\left\lfloor(k-1)^{2} \times \text { golden. ratio }\right\rfloor_{k \geq 3}=6,14,25,40,58,79, \ldots
$$

In higher dimensions we use a Disc. Hankel Transf.

$$
H_{d}^{\text {disc }}(f)(m)=\frac{2}{j_{q+1}} \sum_{n=1}^{q} f(n) \frac{J_{d / 2-1}\left(\frac{j j_{j} n}{j_{q+1}}\right)}{J_{d / 2}\left(j_{n}\right)^{2}} \quad\left[j_{n}=n^{\text {th }} \text {-zero of } J_{d / 2-1}\right]
$$

if $\mathrm{d}=1$ it is a translated DFT.

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if $\mathrm{d}=1$ it is a translated DFT.
Let

$$
\mathcal{A}_{ \pm}^{\text {disc }}[d, q]=\left\{\begin{array}{l}
f, H_{d} f:\{1, . ., q\} \rightarrow \mathbb{R} \\
H_{d} f(0) \leq 0, \pm f(0) \leq 0 .
\end{array}\right.
$$

and

$$
\mathbb{A}_{ \pm}^{\text {disc }}[d, q]:=\min \left\{\sqrt{k(f) k\left( \pm H_{d} f\right)}\right\} .
$$

Then $j_{\mathrm{A}}^{\text {disc }[d, q]}$ $\gtrsim \sqrt{2 \pi j_{q+1}}$

$$
\begin{aligned}
& \frac{j_{d_{ \pm}^{\text {disc }[d, q]}}}{\sqrt{2 \pi j_{q+1}}} \rightarrow \mathbb{A}_{ \pm}(d) \text { numerically }(q \rightarrow \infty) \\
& q \mapsto k=\mathbb{A}_{ \pm}^{\text {disc }}[d, q] \text { is a stairway }
\end{aligned}
$$

$$
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& \frac{j_{\mathbb{A}}^{d_{1} \text { isc }[d, q]}}{} \sqrt{2 \pi j_{q+1}} \rightarrow \mathbb{A}_{ \pm}(d) \text { numerically }(q \rightarrow \infty) \\
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$q_{-}^{j u m p}(2, k) \approx\left\lfloor\frac{\sqrt{3}\left(k^{2}-2 k+2\right)}{4}\right\rfloor_{k \geq 4}=4,7,11,16,21,28,35,43,52,62,$. $q_{-}^{j u m p}(\mathbf{8}, k) \approx\left\lfloor\frac{k^{2}}{4}\right\rfloor_{k \geq 4}=4,6,9,12,16,20,25,30,36,42, \ldots$, $q_{-}^{j u m p}(24, k) \approx\left\lfloor\frac{k^{2}+6 k-8}{8}\right\rfloor_{k \geq 4}=4,5,8,10,13,15,19,22,26,29, \ldots$, $q_{+}^{j u m p}(12, k) \approx\left\lfloor\frac{k^{2}+2 k-1}{4}\right\rfloor_{k \geq 3}=3,5,8,11,15,19,24,29,35,41, \ldots$,

## numerically

$$
\begin{aligned}
& \frac{j_{\mathbb{A}_{-}^{\text {disc }[d, q]}}}{\sqrt{2 \pi j_{q+1}}} \rightarrow 1,\left(\frac{4}{3}\right)^{\frac{1}{4}}, \sqrt{2}, \sqrt{4} \quad(d=1,2,8,24) \\
& \frac{j_{\mathbb{A}_{-}^{d i s c}[d, q]}}{\sqrt{2 \pi j_{q+1}}} \rightarrow \frac{1}{\sqrt{2 \text { golden.ratio }}}, \sqrt{2} \quad(d=1,12)
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\end{aligned}
$$

d Best Packing $\mathbb{A}_{L P}(d)$
1
$\mathbb{Z}$
1

| $\mathbb{A}_{-}(d)$ | $\mathbb{A}_{+}(d)$ |
| :---: | :---: |
| 1 | $? ? ? ?$ |

2 Honeycomb $?=(4 / 3)^{\frac{1}{4}} \quad ?=(4 / 3)^{\frac{1}{4}} \quad$ ?
8
E8
$\sqrt{2}$
$\sqrt{2}$
?
12
?
?
?
$\sqrt{2}$
24
Leech
$\sqrt{4}$
$\sqrt{4}$
?

$$
\begin{aligned}
& \frac{j_{\mathbb{A}_{-}^{d i s c}[d, q]}}{\sqrt{2 \pi j_{q+1}}} \rightarrow 1,\left(\frac{4}{3}\right)^{\frac{1}{4}}, \sqrt{2}, \sqrt{4} \quad(d=1,2,8,24) \\
& \frac{j_{\mathbb{A}_{+}^{d i s c}[d, q]}}{\sqrt{2 \pi j_{q+1}}} \rightarrow \frac{1}{\sqrt{2 \text { golden.ratio }}}, \sqrt{2} \quad(d=1,12)
\end{aligned}
$$

| $d$ | Best Packing | $\mathbb{A}_{L P}(d)$ | $\mathbb{A}_{-}(d)$ | $\mathbb{A}_{+}(d)$ |
| :---: | :---: | :---: | :---: | :---: |
| 1 | $\mathbb{Z}$ | 1 | 1 | $? ? ? ?$ |

2 Honeycomb $?=(4 / 3)^{\frac{1}{4}} \quad ?=(4 / 3)^{\frac{1}{4}} \quad$ ?

| 8 | E8 | $\sqrt{2}$ | $\sqrt{2}$ | $?$ |
| :---: | :---: | :---: | :---: | :---: |
| 12 | $?$ | $?$ | $?$ | $\sqrt{2}$ |
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There are Thm's to be proven here!

