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Sharp isoperimetric inequality, discrete PDEs and Semidiscrete optimal transport

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GEOMETRIC MOTIVATIONS

 Σ_t surface, then $\frac{d}{dt}\Big|_{t=0}$ Area $(\Sigma_t) = 0$: Minimal Surfaces CM = 0.



Otto Frei, 1964

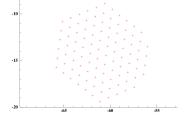
 $\frac{d}{dt}\Big|_{\substack{t=0\\|\Omega_t|=1}} \operatorname{Area}(\partial\Omega_t) = 0: \text{Constant mean curvature } CM = \lambda.$



Classical problem: $\min \{ \text{Area}(\partial \Omega) : |\Omega| = 1 \}$. Solution: $\Omega = B$. How to translate these classical principles to the discrete world?

PHYSICAL MOTIVATIONS

Minimize $\sum_{i \neq j} V(|x_i - x_j|)$ with $V(r) = \frac{1}{r^{12}} - \frac{1}{r^6}$ or similar:



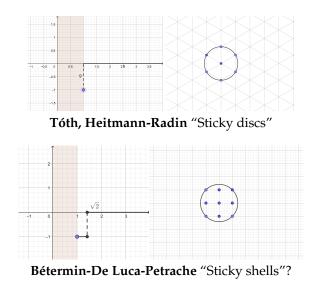
Survey by Blanc-Lewin 2015, N = 100

Tóth 1956, Heitmann-Radin 1980: sticky disk rigidity. **Schmidt 2013, De Luca-Friesecke 2017**: *N*^{3/4}-fluctuations around hexagon (sticky disc).

Theil 2006: minimizer resistent to perturbations of *V* (in the limit $N \rightarrow \infty$)

Bétermin-Petrache 2018 observation: \mathbb{Z}^2 sometimes outperforms A₂. **Bétermin-De Luca-Petrache 2018**: crystallization to \mathbb{Z}^2 for "sticky shell" model (robust under perturbation).

SIMPLIFIED MODELS



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What's the crystal shape?

(Discrete isoperimetric problems)

CONTINUUM ISOPERIMETRIC INEQUALITIES

Find $H \subset \mathbb{R}^d$ such that for all $\Omega \subset \mathbb{R}^d$:

$$\frac{\operatorname{Area}(\partial H)^{d}}{\operatorname{Vol}(H)^{d-1}} \leq \frac{\operatorname{Area}(\partial \Omega)^{d}}{\operatorname{Vol}(\Omega)^{d-1}}$$

- Equivalently: *H* minimizes perimeter at fixed volume.
- Can be extended to general notions of "Vol(Ω)" and "Area($\partial \Omega$)":

 $\begin{array}{l} P_g(\Omega) := \int_{\partial\Omega} g(\nu(x)) dS(x) \\ V_w(\Omega) := \int_{\Omega} w(x) dV(x) \end{array} \right\} \quad \text{con} \quad H : \mathbb{R}^d \to \mathbb{R} \left\{ \begin{array}{l} w, g \ge 0, \\ H(\lambda x) = \lambda H(x) \\ H \text{ convex} \end{array} \right.$

- In this talk we only consider w = 1 for simplicity.
- Optimizer $H = \{x \in \mathbb{R}^d : \forall \nu, x \cdot \nu < g(\nu)\}$ Wulff 1901.
- Del Nin Petrache 2021: discrete-continuum limit for crystals and quasicrystals.

CONTINUUM CASE PROOFS – 1/4

Strategy 1: PDE + convexity

(Cabré-Ros Oton-Serra 2013, Trudinger 1994)

$$\begin{cases} \Delta u(x) = \frac{P_g(\Omega)}{|\Omega|} & x \in \Omega, \\ \frac{\partial u}{\partial \nu}(x) = g(\nu(x)) & x \in \partial \Omega. \end{cases}$$

Solution exists, is regular and unique up to constant summand.

$$\partial_{\Omega} u := \{ x \in \Omega : u(y) \ge u(x) + \nabla u(x) \cdot (y - x) \mid y \in \overline{\Omega} \}.$$

(*x* such that tg. plane to graph of *u* at *x* supports graph($u|_{\Omega}$).)

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CONTINUUM CASE PROOFS -2/4

Claim: $H \subset \nabla u(\partial_{\Omega} u)$.

• $p \in H$ means: $p \cdot \nu < g(\nu)$ for all $\nu \in \mathbb{S}^{d-1}$.

• Let
$$x \in \overline{\Omega}$$
 min of $u(y) - p \cdot y$.

- If $p \in \partial \Omega$ then $\frac{\partial (u(y) p \cdot y)}{\partial \nu} \leq 0 \Leftrightarrow \frac{\partial u}{\partial \nu}(x) \leq p \cdot \nu < g(\nu)$, contradiction.
- ► So *x* is interior. It follows:

• $p = \nabla u(x)$ (*u* smooth),

• $u(y) \ge u(x) + p \cdot (y - x)$ (for all $y \in \overline{\Omega}$).

Therefore $p \in \nabla u(\partial_{\Omega} u)$, proving the claim.

We get
$$|H| \le |\nabla u(\partial_{\Omega} u)| = \int_{\nabla u(\partial_{\Omega} u)} dp \le \int_{\partial_{\Omega} u} \det[D^2 u(x)] dx$$

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CONTINUUM CASE PROOFS -3/4

Linear algebra: $\lambda_1 \leq \lambda_2 \leq \cdots \leq \lambda_d$ eigenvalues of $D^2u(x)$, then

$$\det[D^2 u(x)] = \prod_{j=1}^d \lambda_j \le \left(\frac{1}{d} \sum_{j=1}^d \lambda_j\right)^d = \left(\frac{\operatorname{tr}[D^2 u(x)]}{d}\right)^d = \left(\frac{\Delta u(x)}{d}\right)^d.$$

We get:

$$\begin{split} |H| &\leq |\partial_{\Omega} u| \left(\frac{\Delta u(x)}{d}\right)^{d} = |\partial_{\Omega} u| \left(\frac{P_{g}(\Omega)}{d \cdot |\Omega|}\right)^{d} \leq |\Omega| \left(\frac{P_{g}(\Omega)}{d \cdot |\Omega|}\right)^{d}.\\ &\frac{|H|}{d^{d}} \leq \frac{P_{g}(\Omega)^{d}}{|\Omega|^{d-1}}. \end{split}$$

Finally, due to the duality $W \leftrightarrow H$:

$$P_g(H) = \int_{\partial H} g(\nu(x)) dS \stackrel{*}{=} \int_{\partial H} x \cdot \nu(x) dS = \int_H \operatorname{div}(x) dx = d|H|.$$

CONTINUUM CASE PROOFS -4/4

Strategy 2: Mass transportation proof (sketch)

(Cordero Erausquin - Nazaret - Villani 2004, Figalli-Maggi-Pratelli 2010, Gromov 1983)

- Let $T : \Omega \to H$ transport density $\frac{|H|}{|\Omega|} 1_{\Omega}(x)$ to $1_H(x)$.
- $g^*(y) := \min\{\lambda > 0 : \lambda y \in H\}$ dual norm to *g*.
- Then $g^*(T(x)) \leq 1$ for a.e. $x \in \Omega$.

$$d|\Omega| \left(\frac{|H|}{|\Omega|}\right)^{\frac{1}{d}} = d \int_{\Omega} (\det \nabla T)^{\frac{1}{d}} dx$$

$$\leq \int_{\Omega} \operatorname{div} T \, dx = \int_{\partial^*\Omega} T \cdot \nu_{\Omega} \, dS$$

$$\leq \int_{\partial^*\Omega} g^*(T(x)) g(\nu_{\Omega}(x)) dS \leq P_g(\Omega).$$

DISCRETE SHARP ISOPERIMETRIC INEQUALITY

Setup:

- $V \subset \mathbb{R}^d$ possible positions of points (atoms).
- G = (V, E) undirected graph of possible edges (bonds).
- $g: E \to [0, +\infty)$ weight (energy) of edges.

Looking for edge-isoperimetric inequalities of the form

$$\forall \Omega \subset V \text{ finite}, \quad (\sharp \Omega)^{d-1} \leq C(\sharp_g \partial \Omega)^d := C \left(\sum_{(x,y) \in \overrightarrow{\partial \Omega}} g(x,y) \right)^d.$$

Sharp inequality: Equality actually achieved for some Ω ⊂ V.
 Interesting case: Equality achieved for ∞-many values of #Ω.

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SAMPLE RESULTS

- Hamamuki 2014: \mathbb{Z}^d product graph with nearest-neighbor edges, constant *g* (cubes optimize).
- ► Gomez-Petrache 2020: The triangular lattice (with g = 1) does not have a sharp inequality as above, but rather

$$\frac{(\sharp \overrightarrow{\partial \Omega} - 6)^2}{4 \sharp \Omega - \sharp \overrightarrow{\partial \Omega} + 2} \geq 12,$$

optimized only by "perfect hexagons" (follows via result for honeycomb graph).

- Gomez-Petrache 2020: In honeycomb graph hexagons optimize.
- **Gomez-Petrache 2020**: In the 1-skeleton of the Voronoi diagram of the BCC, rhombic dodecahedron configurations optimize.
- ► **Gomez-Petrache 2020**: Sharp inequalities occur for special geometries of *V*, such as reciprocals of Coxeter triangulations.

THE DISCRETE RESULT

- Auxiliary symmetric $A : V \times V \to \mathbb{R}$ with $E = \{A = 0\}$, defines Laplacian $\Delta_A u(x) := \sum_{y \in V} A^2(y, x)(u(x) - u(y))$.
- Solve discrete PDI (discrete PDE not solvable in general)

$$\begin{cases} \Delta_A u(x) \leq \frac{1}{\sharp\Omega} \sum_{(x,y) \in \overrightarrow{\partial\Omega}} g(x,y) & \text{for } x \in \Omega \\ u(y) - u(x) = \frac{g(x,y)}{A(x,y)} & \text{for } (x,y) \in \overrightarrow{\partial\Omega}. \end{cases}$$

► Some notations:

$$\begin{aligned} \partial u(x) &:= \{ p \in \mathbb{R}^d : (\forall z \in \overline{\Omega}), \ u(x) \leq u(z) + p \cdot (x - z) \}, \\ \partial^{\text{prox}} u(x) &:= \{ p \in \mathbb{R}^d : (\forall z : z \sim x), \ u(x) \leq u(z) + p \cdot (x - z) \}, \\ H_g &:= \left\{ p \in \mathbb{R}^d : (\forall (x, y) \in \overrightarrow{\partial \Omega}), \ p \cdot (y - x) \leq \frac{g(x, y)}{A(x, y)} \right\}. \end{aligned}$$

THE DISCRETE RESULT (GOMEZ-PETRACHE 2020)

We find an analogue to strategy 1:

$$\begin{aligned} |H_{g}| & \stackrel{(a)}{\leq} \quad \left| \bigcup_{x \in \Omega} \partial u(x) \right| \stackrel{(b)}{=} \sum_{x \in \Omega} |\partial u(x)| \\ &\stackrel{(c)}{\leq} \quad \sum_{x \in \Omega} |\partial^{\text{prox}} u(x)| \\ &\stackrel{(d)}{\leq} \quad \sum_{x \in \Omega} c_{x} \left(\Delta_{A} u(x) \right)^{d} \\ &\stackrel{(e)}{\leq} \quad \left(\max_{x \in \Omega} c_{x} \right) \sharp \Omega \left(\frac{1}{\sharp \Omega} \sum_{(x,y) \in \overrightarrow{\partial \Omega}} g(x,y) \right)^{d} = \max_{x \in \Omega} c_{x} \frac{\left(\sharp_{g} \overrightarrow{\partial \Omega} \right)^{d}}{(\sharp \Omega)^{d-1}}. \end{aligned}$$

Generalization of arithmetic-geometric inequality:

 $|\partial^{\operatorname{prox}} u(x)| \leq c_x |\Delta_A u(x)|^d.$

Main Result – 1/2

Necessary conditions for equality:

- ► $c_x, x \in \Omega$ are all equal.
- G_Ω dual graph of a face-to-face decomposition of ∂u(Ω) into convex polyhedra of equal volume
- ► The PDI becomes a PDE (equality achieved).

Sufficient conditions for equality. (assume $\Omega \subset V$ connected in *G*).

- The complex made of vertices and edges of G|_Ω ∪ ∂Ω is reciprocal to the collection of *d* and (*d* − 1)-cells of an equal-volume Voronoi tessellation of a convex polyhedron *H*.
- ► If F_{x,y} denotes the (d − 1)-dimensional facet of the Voronoi tessellation which is dual to edge (x, y) of G, then A²(x,y)|y-x| H^{d-1}(F_{x,y}) takes the same value for all edges {x, y} of G.

Main Result – 2/2

Relation to optimal transport

Functions *u* achieving equality are precisely of the form $u = \lambda u_{Alek} + \ell$ where $\lambda > 0$, ℓ is an affine function and u_{Alek} is the Aleksandrov solution to a semidiscrete optimal transport problem between $\frac{|H|}{\sharp\Omega} \sum_{x \in \Omega} \delta_x$ and $1_H(x) dx$ with transport cost $|x - y|^2$.

The subdifferential optimization problem

Set $H_v(c) := \{p : v \cdot p \le c\}$ and for fixed $\mathcal{V} \subset \mathbb{R}^d$ we define

$$C_{\mathcal{V}} := \max \left\{ \left| \bigcap_{v \in \mathcal{V}} H_v(c_v) \right| : \vec{c} = (c_v)_{v \in \mathcal{V}} \in \mathbb{R}^{\mathcal{V}}, \ \sum_{v \in \mathcal{V}} c_v = 1 \right\}.$$

• Optimizer exists iff Span $\mathcal{V} = \mathbb{R}^d$, $\sum_{v \in \mathcal{V}} v = 0$.

• The face F_v with normal v of the optimizer has area $d C_{\mathcal{V}}|v|$.

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THE $\partial^{\text{prox}} u(x)$ -INEQUALITY

- ► We need $|\partial^{\text{prox}} u(x)| \le c_x (\Delta_A u(x))^d$ and discussion of equality case if $c_x, x \in \Omega$ all equal.
- Apply the above with $\mathcal{V}_x := \{(y x)A^2(x, y) : y \sim x\}$, and then $c_x = C_{\mathcal{V}_x}$.
- We have

$$\frac{1}{\Delta_A u(x)} \partial^{\text{prox}} u(x) = \bigcap_{v \in \mathcal{V}_x} H_v(c_v) \quad \text{with} \quad c_v = \frac{A^2(y, x)(u(y) - u(x))}{\Delta_A u(x)}.$$

• Equality case (assuming Ω connected): $|F_{x,y}|/(|y-x|A^2(x,y))$ constant over all edges.

FURTHER CONNECTIONS/DIRECTIONS

- Aurenhammer 1987, Rybnikov 1999: translation between liftings, weighted Voronoi tessellations, reciprocal graphs.
- Mérigot 2013, Benamou-Froese 2017: link of the above to semidiscrete optimal transport.
- Trudinger 1994: further continuum isoperimetric inequalities (possible extension to operators of higher degree).
- Optimal discrete PDI in general graphs: first bounds in Gomez-Petrache 2020, general Cheeger type bounds to be explored.
- More complicated optimal inequalites in periodic graphs do exist (cf. triangular graph case above), classification missing.