

A limited-memory Riemannian symmetric rank-one trust-region method with a restart strategy

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Joint work with Kyle Gallivan

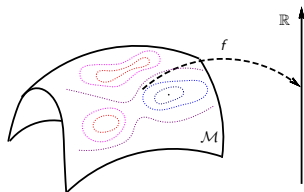
Riemannian Optimization

Problem Statement

Problem: Given $f(x) : \mathcal{M} \rightarrow \mathbb{R}$,
 solve

$$\min_{x \in \mathcal{M}} f(x)$$

where $\mathcal{M}$ is a Riemannian manifold.



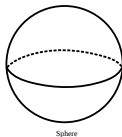
Unconstrained optimization problem on a constrained space.

Riemannian manifold = manifold + Riemannian metric

Riemannian Manifold

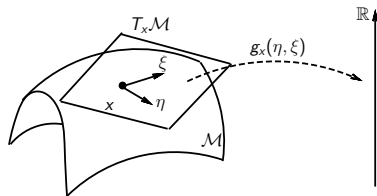
Problem Statement

Manifolds:



- Stiefel manifold: $\text{St}(p, n) = \{X \in \mathbb{R}^{n \times p} | X^T X = I_p\}$;
- Grassmann manifold $\text{Gr}(p, n)$: all p -dimensional subspaces of $\mathbb{R}^n$;
- And many more.

Riemannian metric:



A Riemannian metric, denoted by g , is a smoothly-varying inner product on the tangent spaces;

Riemannian Manifold

A Nonexhaustive Review

- Steepest descent: Smith 1994; Helmke-Moore 1994; Iannazzo-Porcelli 2019;
- Conjugate gradient: Smith 1994; Gallivan-Absil 2010; Ring-Wirth 2012; Sato-Iwai 2015;
- Trust region Newton: Absil-Baker-Gallivan 2007;
- Quasi-Newton:
 - BFGS Ring-Wirth 2012;
 - Trust-region SR1 Huang-Absil-Gallivan 2015;
 - Broyden family (including BFGS) Huang-Absil-Gallivan 2015;
 - BFGS, LBFGS Huang-Absil-Gallivan 2018;
 - Limited-memory trust-region SR1 Huang-Gallivan 2022;

Riemannian Manifold

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 - Limited-memory trust-region SR1 Huang-Gallivan 2022;

A Limited-memory Trust-region Symmetric Rank-One Update Method with a Restart Strategy (LRTRSR1-R)

The idea is based on

- Intrinsic representation of tangent vectors¹;
- An efficient solver for the Euclidean LSR1 subproblem²;
- Slight modifications to LRTRSR1;

¹W. Huang, P.-A. Absil, and K. A. Gallivan. Intrinsic representation of tangent vectors and vector transport on matrix manifolds. *Numerische Mathematik*, 136(2):523-543, 2017.

²Johannes Brust, Jennifer B. Erway, and Roummel F. Marcia. On solving L-SR1 trust- region subproblems. *Computational Optimization and Applications*, 66(2):245-266, Mar 2017

Euclidean TRSR1 to Riemannian TRSR1

Euclidean trust region SR1

Require: Initial iterate x_0 ;

1, Set $k \leftarrow 0$;

while not accurate enough **do**

2, Compute s_k by approximately solving

$$\operatorname{argmin}_{\|s\| \leq \Delta_k} \nabla f(x_k)^T s + \frac{1}{2} s^T B_k s;$$

$$3, \rho_k \leftarrow \frac{f(x_k) - f(x_k + s_k)}{m_k(0) - m_k(s_k)};$$

$$4, x_{k+1} \leftarrow \begin{cases} x_k + s_k & \text{if } \rho_k > c \\ x_k & \text{otherwise.} \end{cases}$$

5, Update radius to get Δ_{k+1} ;

6, Update B_k to get B_{k+1} ;

7, $k \leftarrow k + 1$;

end while

Riemannian trust region SR1

Require: Initial iterate x_0 ;

1, Set $k \leftarrow 0$;

while not accurate enough **do**

2, Compute s_k by approximately solving

$$\operatorname{argmin}_{\|s\| \leq \Delta_k, s \in T_{x_k} \mathcal{M}} g(\operatorname{grad} f(x_k), s) + \frac{1}{2} g(s, B_k s);$$

$$3, \rho_k \leftarrow \frac{f(x_k) - f(R_{x_k}(s_k))}{m_k(0) - m_k(s_k)};$$

$$4, x_{k+1} \leftarrow \begin{cases} R_{x_k}(s_k) & \text{if } \rho_k > c \\ x_k & \text{otherwise.} \end{cases}$$

5, Update radius to get Δ_{k+1} ;

6, Update B_k to get B_{k+1} ;

7, $k \leftarrow k + 1$;

end while

Trust-region LSR Subproblem

Euclidean Version

Local problem in Euclidean trust region:

$$\min_{\|s\| \leq \Delta} m(s) = s^T g + \frac{1}{2} s^T B s;$$

- at k -th iteration: $\Delta = \Delta_k$, $g = \nabla f(x_k)$ and $B = B_k$;
- SR1 update: $(y_k = \nabla f(x_{k+1}) - \nabla f(x_k))$

$$B_{k+1} = B_k + \frac{(y_k - B_k s_k)(y_k - B_k s_k)^T}{(y_k - B_k s_k)^T s_k};$$

- Compact representation:

$$\begin{aligned} B_{k+1} &= B_{k-m+1} + (Y_{k+1} - B_{k-m+1} S_{k+1}) \\ &\quad (D_{k+1} + L_{k+1} + L_{k+1}^T - S_{k+1}^T B_{k-m+1} S_{k+1})^{-1} (Y_{k+1} - B_{k-m+1} S_{k+1})^T; \\ &= \gamma I + \Phi M^{-1} \Phi^T; \end{aligned}$$

where $Y_{k+1} = [y_{k-m+1}, \dots, y_k]$, $S_{k+1} = [s_{k-m+1}, \dots, s_k]$, and $\gamma I = B_{k-m+1}$.

Trust-region LSR Subproblem

Euclidean Version

Local problem in Euclidean trust region:

$$\min_{\|s\| \leq \Delta} m(s) = s^T g + \frac{1}{2} s^T B s;$$

Theorem ([NW06])

The vector s^ is a global solution of the trust region subproblem*

$$\min_{\|s\| \leq \Delta} \nabla f(x)^T s + \frac{1}{2} s^T B s$$

if and only if s^ is feasible and there is a scalar $\lambda \geq 0$ such that the following conditions hold:*

$$(B + \lambda I)s^* = -\nabla f(x), \quad \lambda(\Delta - \|s^*\|) = 0, \quad (B + \lambda I) \succeq 0.$$

Trust-region LSR Subproblem

Euclidean Version

Local problem in Euclidean trust region:

$$\min_{\|s\| \leq \Delta} m(s) = s^T g + \frac{1}{2} s^T B s;$$

Eigenvalue decomposition

$$B = \gamma I + \Phi M^{-1} \Phi^T = [P \ P_{\perp}] \left(\gamma I + \begin{bmatrix} S & \\ & 0_{(n-m) \times (n-m)} \end{bmatrix} \right) [P \ P_{\perp}]^T,$$

where $\Phi = QR$ is qr decomposition, $RM^{-1}R^T = USU^T$ is eigenvalue decomposition, $P = QU$ and $[P, P_{\perp}]$ is orthonormal.

- Eigenvalues of B : $\Lambda = \gamma + \text{diag}(s_1, \dots, s_m, 0, \dots, 0)$;
- $s(\lambda) = -(B + \lambda I)^{-1}g$ for finding the solution s^* ;
- $h(\lambda) = \frac{1}{\|s(\lambda)\|} - \frac{1}{\Delta}$ for finding the appropriate λ ;
- Complexity: $O(nm) + O(m^3)$;

Trust-region LSR Subproblem

Riemannian Version

Local problem in Riemannian trust region:

$$\min_{\|\mathfrak{s}\| \leq \Delta, \mathfrak{s} \in T_x \mathcal{M}} m(\mathfrak{s}) = g_x(\mathfrak{s}, g) + \frac{1}{2} g_x(\mathfrak{s}, \mathcal{B}\mathfrak{s});$$

Turn Riemannian LSR1 subproblem into Euclidean version:

- Q_x : an orthonormal basis of $T_x \mathcal{M}$;
- $T_x \mathcal{M} = \{Q_x c : c \in \mathbb{R}^d\}$;
- Replace $\mathfrak{s}$ by $Q_x c$:

$$\min_{\|Q_x c\| \leq \Delta} c^T Q_x^b g + \frac{1}{2} c^T (Q_x^b \mathcal{B} Q) c = \min_{\|c\| \leq \Delta} c^T \hat{g} + \frac{1}{2} c^T \hat{\mathcal{B}} c;$$

- Solved by Euclidean version algorithm

Trust-region LSR Subproblem

Riemannian Version

Local problem in Riemannian trust region:

$$\min_{\|\mathfrak{s}\| \leq \Delta, \mathfrak{s} \in T_x \mathcal{M}} m(\mathfrak{s}) = g_x(\mathfrak{s}, g) + \frac{1}{2} g_x(\mathfrak{s}, \mathcal{B}\mathfrak{s});$$

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- Solved by Euclidean version algorithm

Orthonormal basis Q_x is crucial!

Intrinsic Representation of Tangent Vector

- Given $\eta_x \in T_x \mathcal{M}$, compute $c_x = Q_x^b \eta_x$, i.e., find $c_x \in \mathbb{R}^d$ such that $\eta_x = Q_x c_x$;
- Given $c_x \in \mathbb{R}^d$, compute $\eta_x = Q_x c_x$;

-
- Stiefel manifold as an example: $\text{St}(p, n) = \{X \in \mathbb{R}^{n \times p} : X^T X = I_p\}$.
 - Tangent space at $X \in T_X \mathcal{M}$ is
 $T_X \text{St}(p, n) = \{X\Omega + X_\perp K \mid \Omega^T = -\Omega, X^T X_\perp = 0\}$;

$$Q_x = \left\{ \begin{bmatrix} X & X_\perp \end{bmatrix} \begin{bmatrix} 0 & 1 & \dots & 0 \\ -1 & 0 & \dots & 0 \\ \dots & \dots & \dots & \dots \\ 0 & 0 & \dots & 0 \\ 0 & 0 & \dots & 0 \\ 0 & 0 & \dots & 0 \\ \dots & \dots & \dots & \dots \\ 0 & 0 & \dots & 0 \end{bmatrix}, \dots, \begin{bmatrix} X & X_\perp \end{bmatrix} \begin{bmatrix} 0 & 0 & \dots & 0 \\ 0 & 0 & \dots & 0 \\ \dots & \dots & \dots & \dots \\ 0 & 0 & \dots & 0 \\ 1 & 0 & \dots & 0 \\ 0 & 0 & \dots & 0 \\ \dots & \dots & \dots & \dots \\ 0 & 0 & \dots & 0 \end{bmatrix} \right\}$$

Intrinsic Representation of Tangent Vector

$$T_X \text{St}(p, n) = \{X\Omega + X_\perp K \mid \Omega^T = -\Omega, X^T X_\perp = 0\};$$

Extrinsic η_X :

$$\eta_X = [X \quad X_\perp] \begin{bmatrix} \Omega \\ K \end{bmatrix}$$

$$= [X \quad X_\perp] \begin{bmatrix} 0 & a_{12} & \dots & a_{1p} \\ -a_{12} & 0 & \dots & a_{2p} \\ \dots & \dots & \dots & \dots \\ -a_{1p} & -a_{2p} & \dots & 0 \\ b_{11} & b_{12} & \dots & b_{1p} \\ b_{21} & b_{22} & \dots & b_{2p} \\ \dots & \dots & \dots & \dots \\ b_{(n-p)1} & b_{(n-p)2} & \dots & b_{(n-p)p} \end{bmatrix}$$

Intrinsic c_X :

$$c_X = \begin{bmatrix} a_{12} \\ a_{13} \\ a_{23} \\ \vdots \\ a_{(p-1)p} \\ b_{11} \\ b_{21} \\ \vdots \\ b_{(n-p)p} \end{bmatrix}$$

Intrinsic Representation of Tangent Vector

Question

Extrinsic representation $\eta_X \iff$ Intrinsic representation c_X

$$\bullet \eta_X = \begin{bmatrix} X & X_{\perp} \end{bmatrix} \begin{bmatrix} \Omega \\ K \end{bmatrix} \iff \begin{bmatrix} \Omega \\ K \end{bmatrix} \iff c_X$$

Intrinsic Representation of Tangent Vector

Question

Extrinsic representation $\eta_X \iff$ Intrinsic representation c_X

- $\eta_X = [X \quad X_\perp] \begin{bmatrix} \Omega \\ K \end{bmatrix} \iff \begin{bmatrix} \Omega \\ K \end{bmatrix} \iff c_X$
- Apply Householder transformation to X , (Done in retraction)

$$Q_p^T Q_{p-1}^T \dots Q_1^T X = R = I_{n \times p}.$$

Intrinsic Representation of Tangent Vector

Question

Extrinsic representation $\eta_X \iff$ Intrinsic representation c_X

- $\eta_X = [X \quad X_\perp] \begin{bmatrix} \Omega \\ K \end{bmatrix} \iff \begin{bmatrix} \Omega \\ K \end{bmatrix} \iff c_X$

- Apply Householder transformation to X , (Done in retraction)

$$Q_p^T Q_{p-1}^T \dots Q_1^T X = R = I_{n \times p}.$$

- $[X \quad X_\perp] = Q_1 Q_2 \dots Q_p$ (Do not compute)

Intrinsic Representation of Tangent Vector

Question

Extrinsic representation $\eta_X \iff$ Intrinsic representation c_X

- $\eta_X = [X \quad X_\perp] \begin{bmatrix} \Omega \\ K \end{bmatrix} \iff \begin{bmatrix} \Omega \\ K \end{bmatrix} \iff c_X$
- Apply Householder transformation to X , (Done in retraction)

$$Q_p^T Q_{p-1}^T \dots Q_1^T X = R = I_{n \times p}.$$
- $[X \quad X_\perp] = Q_1 Q_2 \dots Q_p$ (Do not compute)
- Extrinsic to Intrinsic: $Q_p^T Q_{p-1}^T \dots Q_1^T \eta_X = \begin{bmatrix} \Omega \\ K \end{bmatrix}$ and reshape to c_X ;
($4np^2 - 2p^3$) flops
- Intrinsic to Extrinsic: reshape c_X and $\eta_X = Q_1 Q_2 \dots Q_p \begin{bmatrix} \Omega \\ K \end{bmatrix}$;
($4np^2 - 2p^3$) flops

Intrinsic Representation of Tangent Vector

- The basis Q_x is orthonormal with respect to the canonical metric:

$$g_X(\eta_X, \xi_X) = \text{trace} \left(\eta_X^T \left(I_n - \frac{1}{2} X X^T \right) \xi_X \right), \forall \eta_X, \xi_X \in T_X \text{St}(p, n).$$

- The above idea can be applied to
 - Fixed-rank real matrices $\mathbb{R}_r^{m \times n} = \{X \in \mathbb{R}^{m \times n} : \text{rank}(X) = r\}$;
 - Fixed-rank symmetric positive semidefinite matrices:
 $S_r^n = \{X \in \mathbb{R}^{n \times n} : X = X^T, X \succeq 0, \text{rank}(X) = r\}$;
 - Fixed-rank complex matrices;
 - Fixed-rank Hermitian positive semidefinite matrices;

Riemannian LSR1 Update

$$\min_{\|Q_x c\| \leq \Delta} c^T Q_x^b g + \frac{1}{2} c^T (Q_x^b \mathcal{B} Q_x) c = \min_{\|c\| \leq \Delta} c^T \hat{g} + \frac{1}{2} c^T \hat{\mathcal{B}} c;$$

$$\begin{aligned} \tilde{\mathcal{B}}_k &= \mathcal{B}_{k-1} + \frac{(\eta_{k-1} - \mathcal{B}_{k-1} \mathfrak{s}_{k-1})(\eta_{k-1} - \mathcal{B}_{k-1} \mathfrak{s}_{k-1})^b}{g(\mathfrak{s}_{k-1}, \eta_{k-1} - \mathcal{B}_{k-1} \mathfrak{s}_{k-1})}, \\ \mathcal{B}_k &= \mathcal{T}_{S_{\mathfrak{s}_{k-1}}} \circ \tilde{\mathcal{B}}_k \circ \mathcal{T}_{S_{\mathfrak{s}_{k-1}}}^{-1}, \end{aligned}$$

where $\eta_{k-1} = \mathcal{T}_{S_{\mathfrak{s}_{k-1}}}^{-1} \text{grad } f(R_{x_{k-1}}(\mathfrak{s}_{k-1})) - \text{grad } f(x_{k-1}) \in T_{x_{k-1}} \mathcal{M}$,
 $\mathfrak{s}_{k-1} \in T_{x_{k-1}} \mathcal{M}$, $\mathcal{B}_k$ is a linear operator on $T_{x_k} \mathcal{M}$ and $\mathcal{T}_S$ is an isometric vector transport.

Riemannian LSR1 Update

Compact representation

Theorem

Let $\mathcal{B}_{k-\ell}$ be a symmetric linear operator on $T_{x_{k-\ell}} \mathcal{M}$ and $\mathcal{T}_S$ is an isometric vector transport. Suppose $g(s_{k-i}^{(k)}, \eta_{k-i}^{(k)} - \mathcal{B}_{k-i}^{(k)} s_{k-i}^{(k)}) \neq 0$ for $i = 1, \dots, \ell$. Then the linear operator $\mathcal{B}_k^{(k)}$ is given by

$$\mathcal{B}_k^{(k)} = \mathcal{B}_{k-\ell}^{(k)} + \left(\mathcal{Y}_\ell^{(k)} - \mathcal{B}_{k-\ell}^{(k)} \mathcal{S}_\ell^{(k)} \right) \left(M_\ell^{(k)} \right)^{-1} \left(\mathcal{Y}_\ell^{(k)} - \mathcal{B}_{k-\ell}^{(k)} \mathcal{S}_\ell^{(k)} \right)^b,$$

where $D_\ell^{(k)} = \text{diag}((s_{k-\ell}^{(k)})^b \eta_{k-\ell}^{(k)}, \dots, (s_{k-1}^{(k)})^b \eta_{k-1}^{(k)})$, $\mathcal{S}_\ell^{(k)} = [s_{k-\ell}^{(k)}, \dots, s_{k-1}^{(k)}]$, $\mathcal{Y}_\ell^{(k)} = [\eta_{k-\ell}^{(k)}, \dots, \eta_{k-1}^{(k)}]$, $(L_\ell)_{i,j}^{(k)} = (s_{k-\ell-1+i}^{(k)})^b \eta_{k-\ell-1+i}^{(k)}$ if $i > j$, $(L_\ell)_{i,j}^{(k)} = 0$ otherwise, and $M_\ell^{(k)} = D_\ell^{(k)} + L_\ell^{(k)} + (L_\ell^{(k)})^T - (\mathcal{S}_\ell^{(k)})^b \mathcal{B}_{k-\ell}^{(k)} \mathcal{S}_\ell^{(k)}$ is nonsingular.

- $\mathcal{B}_{k-i}^{(k)}$ denotes $\mathcal{T}_{S_{s_{k-1}}} \circ \dots \circ \mathcal{T}_{S_{s_{k-i}}} \circ \mathcal{B}_{k-i} \circ \mathcal{T}_{S_{s_{k-i}}}^{-1} \circ \dots \circ \mathcal{T}_{S_{s_{k-1}}}^{-1}$;
- $\eta_{k-i}^{(k)}$ denotes $\mathcal{T}_{S_{s_{k-1}}} \circ \dots \circ \mathcal{T}_{S_{s_{k-i}}} \eta_{k-i}$;
- $s_{k-i}^{(k)}$ denotes $\mathcal{T}_{S_{s_{k-1}}} \circ \dots \circ \mathcal{T}_{S_{s_{k-i}}} s_{k-i}$.

Riemannian LSR1 Update

Compact representation

Vector transport by parallelization:

$$\mathcal{T}_{S_{\eta_x}} \xi_x = Q_y Q_x^b \xi_x$$

where $y = R_x(\eta_x)$;

- We have

$$\begin{aligned} \mathcal{B}_{k-i}^{(k)} &= \mathcal{T}_{S_{\mathfrak{s}_{k-1}}} \circ \dots \circ \mathcal{T}_{S_{\mathfrak{s}_{k-i}}} \circ \mathcal{B}_{k-i} \circ \mathcal{T}_{S_{\mathfrak{s}_{k-i}}}^{-1} \circ \dots \circ \mathcal{T}_{S_{\mathfrak{s}_{k-1}}}^{-1} \\ &= Q_{x_k} Q_{x_{k-1}}^b \dots Q_{x_{k-i+1}}^b Q_{x_{k-i}}^b \mathcal{B}_{k-i} Q_{x_{k-i}} Q_{x_{k-i+1}}^b \dots Q_{x_{k-1}}^b Q_{x_k}^b \\ &= Q_{x_k} Q_{x_{k-i}}^b \mathcal{B}_{k-i} Q_{x_{k-i}} Q_{x_k}^b; \end{aligned}$$

- Similarly, $\eta_{k-i}^{(k)} = Q_{x_k} Q_{x_{k-i}}^b \eta_{k-i}$ and $\mathfrak{s}_{k-i}^{(k)} = Q_{x_k} Q_{x_{k-i}}^b \mathfrak{s}_{k-i}$.

Riemannian LSR1 Update

Compact representation

Compact representation:

$$\mathcal{B}_{k+1} = Q_{x_{k+1}} \left(\gamma_{k+1} I + \Phi_k (M_m^{(k+1)})^{-1} \Phi_k^T \right) Q_{x_{k+1}}^b$$

where $\Phi_k = [Q_{x_{k-m+1}}^b (\eta_{k-m+1} - \gamma_{k+1} s_{k-m+1}), \dots, Q_{x_k}^b (\eta_k - \gamma_{k+1} s_k)]$.

$$\min_{\|Q_x c\| \leq \Delta} c^T Q_x^b g + \frac{1}{2} c^T (Q_x^b \mathcal{B} Q_x) c = \min_{\|c\| \leq \Delta} c^T \hat{g} + \frac{1}{2} c^T (\gamma I + \Phi M^{-1} \Phi^T) c$$

- $s_k = Q_{x_k}^b s_k$;
- $g_k = Q_{x_k}^b \text{grad } f(x_k)$;
- $y_k = g_{k+1} - g_k$;
- Compute B_k as the Euclidean version using s_k, y_k ;

Riemannian LSR1 Update

Initial Hessian Approximation γ

Set

$$\gamma = \text{sign}(g(\eta, \varsigma)) \min \left(\tau_\gamma \frac{|g(\eta, \varsigma)|}{g(\varsigma, \varsigma)}, \frac{g(\eta, \eta)}{|g(\eta, \varsigma)|} \right),$$

where τ_γ is a prescribed parameter;

- LBFGS: Initial inverse Hessian approximation is $\frac{g(\varsigma, \eta)}{g(\eta, \eta)}$;
- LSR1: Initial Hessian approximation is $\frac{g(\eta, \eta)}{g(\varsigma, \eta)}$;
- If $\|\eta\| \leq L\|\varsigma\|$, then
 - $\left| \frac{g(\eta, \varsigma)}{g(\varsigma, \varsigma)} \right|$ bounded above;
 - $\left| \frac{g(\eta, \eta)}{g(\varsigma, \eta)} \right|$ may not be bounded above;
- Bounded above: for theoretical analysis

LRTRSR1-R (Part I)

Input: Riemannian manifold $\mathcal{M}$ with Riemannian metric g ; retraction R ; isometric vector transport by parallelization $\mathcal{T}_S$; smooth function f on $\mathcal{M}$; initial iterate $x_0 \in \mathcal{M}$;

- 1: Choose an integer $m > 0$ and real numbers $\bar{\Delta} > 0$, $\Delta_0 \in (0, \bar{\Delta})$, $\nu \in (0, 1)$, $c \in (0, 0.25)$, $\tau_1 \in (0, 1)$, and $\tau_2 > 1$;
- 2: Compute $\text{grad}^d f(x_0) = Q_{x_0}^b \text{grad} f(x_0) \in \mathbb{R}^d$; Set $k \leftarrow 0$, $\ell \leftarrow 0$, $\gamma_0 \leftarrow 1$;
- 3: Obtain $s_k \in \mathbb{R}^d$ by (approximately) solving

$$s_k \approx \arg \min_{\|s\|_2 \leq \Delta_k} f(x_k) + s^T \text{grad}^d f(x_k) + \frac{1}{2} s^T B_k s.$$

- 4: Compute $s_k = Q_{x_k} s_k$;
- 5: Set $\rho_k \leftarrow \frac{f(x_k) - f(R_{x_k}(s_k))}{m_k(0) - m_k(s_k)}$;
- 6: Set $y_k \leftarrow Q_{R_{x_k}(s_k)}^b \text{grad} f(R_{x_k}(s_k)) - \text{grad}^d f(x_k)$;
- 7: **if** $|s_k^T (y_k - B_k s_k)| \geq \nu \|s_k\|_2 \|y_k - B_k s_k\|_2$ **then**
- 8: **if** $\ell = m$ **then**
- 9: $\gamma_{k+1} \leftarrow \text{sign}(y_k^T s_k \min \left(\tau_\gamma \frac{|y_k^T s_k|}{s_k^T s_k}, \frac{y_k^T y_k}{|y_k^T s_k|} \right))$ **if** $|y_k^T s_k| \leq \epsilon$, **and** $\gamma_{k+1} \leftarrow 0$ **otherwise**,
- where** ϵ **denotes the machine epsilon**; Set $\Psi_{k,\ell}$ **and** $M_{k,\ell}$ **to be empty matrices**;
- 10: **else**
- 11: $\gamma_{k+1} \leftarrow \gamma_k$;
- 12: **end if**
- 13: Set $\tilde{\Psi}_{k+1,\ell+1} = [\tilde{\Psi}_{k,\ell}, y_k - \gamma_{k+1} s_k]$; Computing $M_{k+1,\ell+1}$ **by updating** $M_{k,\ell}$; Set $\ell \leftarrow \ell + 1$;

LRTRSR1-R (Part II)

```

14: else
15:   Set  $\gamma_{k+1} \leftarrow \gamma_k$ ,  $M_{k+1,\ell} \leftarrow M_{k,\ell}$ ;
16: end if
17: if  $\rho_k > c$  then
18:    $x_{k+1} \leftarrow R_{x_k}(\mathfrak{s}_k)$ ;
19: else
20:    $x_{k+1} \leftarrow x_k$ ;
21: end if
22: if  $\rho_k > \frac{3}{4}$  then
23:   If  $\|\eta_k\| \geq 0.8\Delta_k$ , then  $\Delta_{k+1} \leftarrow \min(\tau_2\Delta_k, \overline{\Delta})$ ; otherwise  $\Delta_{k+1} \leftarrow \Delta_k$ ;
24: else if  $\rho_k < 0.1$  then
25:    $\Delta_{k+1} \leftarrow \tau_1\Delta_k$ ;
26: else
27:    $\Delta_{k+1} \leftarrow \Delta_k$ ;
28: end if
29:  $k \leftarrow k + 1$ , goto 2 until convergence.
    
```

Complexity:

$$f_f + g_f + r_f + \theta_f + \vartheta_f + 10md + O(m^3),$$

where f_f , g_f , r_f , θ_f , ϑ_f respectively denote the computational cost of one function evaluation, one gradient evaluation, one retraction evaluation, one $Q_x c$ evaluation and one $Q_x^b \eta_x$ evaluation.

Theoretical Results

Assumption

There exists a positive constant δ_r , such that the function f is Lipschitz continuously differentiable with respect to the vector transport $\mathcal{T}_S$ in Ω_{δ_r} , where the Lipschitz continuous differentiability is defined by

$$\|\mathcal{T}_{\xi_x}^{-1}\eta_y - \eta_x\|_x \leq L_v \|\xi_x\|_x,$$

and $\Omega_{\delta_r} = \{R_x(\eta_x) : x \in \Omega, \|\eta_x\| \leq \delta_r\}$.

Theorem

Suppose above Assumption holds and the parameter $\bar{\Delta}$ in LRTRSR1-R satisfies $\bar{\Delta} \leq \delta_r$. Then there exists a constant β such that the sequence $\{\mathcal{B}_k\}$ generated by LRTRSR1-R satisfies that $\|\mathcal{B}_k\| \leq \beta$ for all k .

Theoretical Results

Lemma

Suppose that $\mathcal{T}_S \in C^1$, $f \in C^2$ in $\mathcal{U}$, and $\mathcal{U}$ is compact. Further suppose that there exists a positive constant M_r such that for any $x \in \mathcal{U}$, it holds that $\{R_x(\eta_x) : \|\eta_x\| > M_r\} \cap \mathcal{U} = \emptyset$ or $\{R_x(\eta_x) : \|\eta_x\| \leq M_r\} \supseteq \mathcal{U}$. Then there exists a constant L_v such that the inequality

$$\|\mathcal{T}_{\eta_x}^{-1} \text{grad } f(y) - \text{grad } f(x)\|_x \leq L_v \|\eta_x\|_x,$$

holds for all $x \in \mathcal{U}$, $\eta_x \in T_x \mathcal{M}$ satisfying $R_x(t\eta_x) \in \mathcal{U}$ for all $t \in [0, 1]$, where $y = R_x(\eta_x)$.

Theoretical Results

The global convergence follows from [HAG15, Theorem 1].

Corollary

Under reasonable assumptions and the parameter $\bar{\Delta}$ in LRTRSR1-R satisfies $\bar{\Delta} \leq \delta_r$. Then (i) if $f \in C^2$ is bounded below on the sublevel set Ω , then $\lim_{k \rightarrow \infty} \text{grad } f(x_k) = 0$; (ii) if $f \in C^2$, $\mathcal{M}$ is compact, then $\lim_{k \rightarrow \infty} \text{grad } f(x_k) = 0$, $\{x_k\}$ has at least one limit point, and every limit point of $\{x_k\}$ is a stationary point of f ; and (iii) if $f \in C^2$, the sublevel set Ω is compact, f has a unique stationary point x^ in Ω , then $\{x_k\}$ converges to x^* .*

Numerical Experiments

- Convex quadratic problem
- Matrix completion
- Joint diagonalization
- Phase retrieval

Numerical Experiments

Convex Quadratic Problem

Problem:

$$\min_{x \in \mathbb{R}^n} \frac{1}{2} x^T A x,$$

where $A = A^T$ and $A \succ 0$;

- Compare with LRBFGS, LRBFGS-R, LRTRSR1, LRTRSR1-R;
- Two types of matrices A ;
 - Type 1: $A \in \mathbb{R}^{200 \times 200}$, 4 different eigenvalues, condition number ≈ 100 ;
 - Type 2: $A \in \mathbb{R}^{1000 \times 1000}$, random matrices, condition number $\approx 10^6$;

Numerical Experiments

Convex Quadratic Problem

Table: An average result of 20 random runs for the convex quadratic problem with the two types of matrices A .

		LRBFGS		LRBFGS-R		LRTRSR1		LRTRSR1-R	
	m	2	4	2	4	2	4	2	4
Type 1	<i>iter</i>	27	23	30	24	37	5	37	5
	<i>nf</i>	31	26	34	28	38	6	38	6
	<i>ng</i>	28	24	31	25	38	6	38	6
	$\frac{\ g_f^f\ }{\ g_0^f\ }$	5.76 ₋₇	5.19 ₋₇	6.01 ₋₇	5.26 ₋₇	5.05 ₋₇	3.25 ₋₉	4.67 ₋₇	2.83 ₋₉
	<i>f</i>	6.97 ₋₁₀	2.26 ₋₉	4.07 ₋₉	1.84 ₋₉	4.59 ₋₉	7.96 ₋₁₂	4.08 ₋₉	4.34 ₋₁₂
Type 2	<i>iter</i>	830	733	837	815	1323	1475	747	680
	<i>nf</i>	887	772	920	883	1324	1476	748	681
	<i>ng</i>	831	734	838	816	1324	1476	748	681
	$\frac{\ g_f^f\ }{\ g_0^f\ }$	9.25 ₋₇	9.18 ₋₇	9.09 ₋₇	9.14 ₋₇	9.53 ₋₇	9.54 ₋₇	9.17 ₋₇	9.05 ₋₇
	<i>f</i>	3.12 ₋₄	1.73 ₋₄	3.49 ₋₄	2.92 ₋₄	7.13 ₋₄	5.77 ₋₄	5.51 ₋₄	6.33 ₋₄

Numerical Experiments

Matrix Completion

Optimization formulation in [Van13]:

$$\min_{X \in \mathbb{R}_r^{m \times n}} f(X) := \frac{1}{2} \|P_\Omega(X - A)\|_F^2,$$

where $\mathbb{R}_r^{m \times n} = \{X \in \mathbb{R}^{m \times n} : \text{rank}(X) = r\}$ is the manifold of fixed-rank matrices, Ω is a given index set, $P_\Omega(A)$ is the given sparse matrix, and $P_\Omega(Z_{i,j}) = Z_{i,j}$ if $(i,j) \in \Omega$; $P_\Omega(Z_{i,j}) = 0$ otherwise,

- Compare with LRBFGS, RCGR [Van13], and LMAFit [WYZ12];

Numerical Experiments

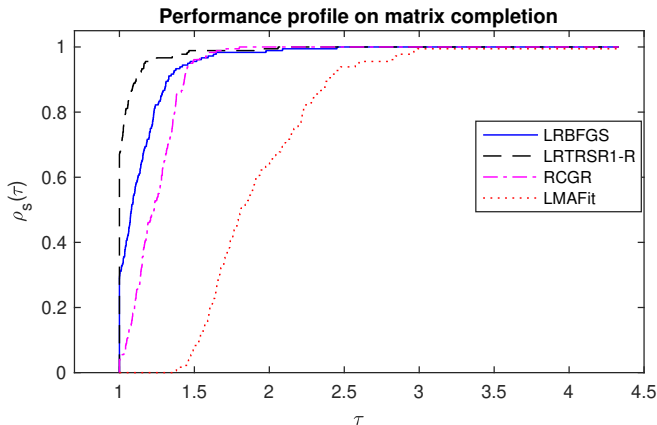
Matrix Completion

Table: An average result of 20 random runs of LRBFGS (i), LRTRSR1-R (ii), RCGR (iii), and LMAFit (iv) with multiple sizes and ranks are reported. *rel* denotes the relative residual $rel = \|P_{\Omega}(X - A)\|_F / \|P_{\Omega}(A)\|_F$.

	r	n = m = 2000			n = m = 4000			n = m = 8000		
		10	20	30	10	20	30	10	20	30
(i)	<i>iter</i>	132	97	101	145	99	93	165	107	93
	<i>nf</i>	159	112	115	172	118	109	195	125	111
	<i>rel</i>	8.11 ₋₁₁	7.13 ₋₁₁	6.67 ₋₁₁	8.71 ₋₁₁	7.30 ₋₁₁	7.10 ₋₁₁	8.59 ₋₁₁	7.69 ₋₁₁	7.16 ₋₁₁
	<i>time</i>	9.47 ₋₁	1.50	2.74	1.99	3.63	6.18	5.20	9.09	1.43 ₁
(ii)	<i>iter</i>	124	106	96	143	110	105	169	133	112
	<i>nf</i>	125	107	97	144	111	106	170	134	113
	<i>rel</i>	6.13 ₋₁₁	6.66 ₋₁₁	7.53 ₋₁₁	7.30 ₋₁₁	6.60 ₋₁₁	6.53 ₋₁₁	6.96 ₋₁₁	6.66 ₋₁₁	7.51 ₋₁₁
	<i>time</i>	7.18 ₋₁	1.41	2.29	1.62	3.39	6.26	4.77	9.76	1.45 ₁
(iii)	<i>iter</i>	80	62	57	89	66	61	101	72	64
	<i>nf</i>	81	63	58	90	67	62	102	73	65
	<i>rel</i>	9.17 ₋₁₁	8.46 ₋₁₁	7.94 ₋₁₁	9.34 ₋₁₁	8.92 ₋₁₁	8.19 ₋₁₁	9.51 ₋₁₁	9.09 ₋₁₁	8.49 ₋₁₁
	<i>time</i>	9.06 ₋₁	1.55	2.40	2.02	4.35	7.02	5.96	1.10 ₁	1.81 ₁
(iv)	<i>iter</i>	338	197	164	402	228	185	491	252	206
	<i>nf</i>	339	198	165	403	229	186	492	253	207
	<i>rel</i>	9.61 ₋₁₁	9.42 ₋₁₁	9.05 ₋₁₁	9.68 ₋₁₁	9.44 ₋₁₁	9.19 ₋₁₁	9.72 ₋₁₁	9.41 ₋₁₁	9.31 ₋₁₁
	<i>time</i>	1.54	2.24	3.46	3.82	5.94	9.62	1.02 ₁	1.89 ₁	2.43 ₁

Numerical Experiments

Matrix Completion



Numerical Experiments

Joint Diagonalization

Joint diagonalization problem [TCA09] is

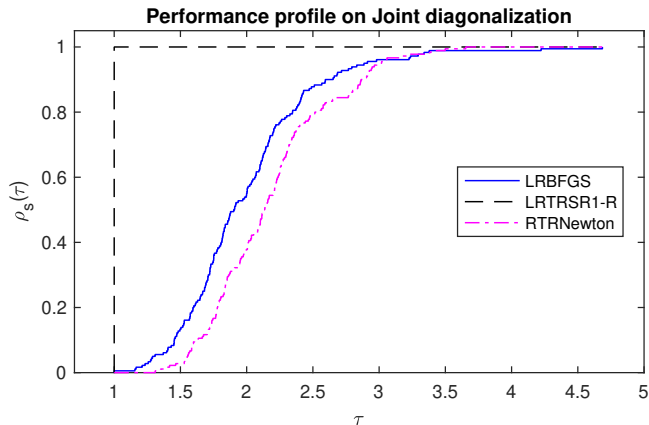
$$f : \text{St}(p, n) \rightarrow \mathbb{R} : X \mapsto - \sum_{i=1}^N \|\text{diag}(X^T C_i X)\|_2^2,$$

where $C_1, \dots, C_N$ are given symmetric matrices, $\text{diag}(M)$ denotes the vector formed by the diagonal entries of M .

- Compare with LRBFGS and RTRNewton [TCA09]

Numerical Experiments

Joint Diagonalization



Numerical Experiments

Phase Retrieval

The phase retrieval problem in [CLS16] is

$$\min_{x \in \mathbb{C}^n} f(x) := \frac{1}{\|b\|_2^2} \|b - \text{diag}(Zxx^*Z^*)\|_2^2,$$

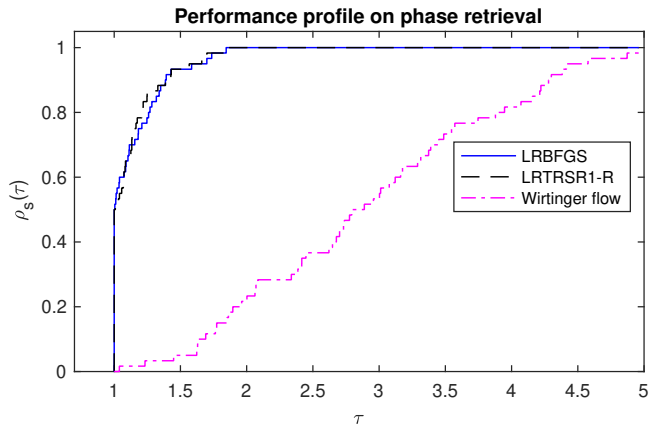
where $b \in \mathbb{R}^m$ is the observation measurements, $Z \in \mathbb{C}^{m \times n}$ is

$$Z = \begin{bmatrix} (\mathcal{F}_s \otimes \mathcal{F}_{s-1} \otimes \dots \mathcal{F}_1) \text{ddiag}(w_1) \\ \vdots \\ (\mathcal{F}_s \otimes \mathcal{F}_{s-1} \otimes \dots \mathcal{F}_1) \text{ddiag}(w_l) \end{bmatrix},$$

$w_i \in \mathbb{C}^n, i = 1, \dots, l$ denotes masks, $\mathcal{F}_i \in \mathbb{C}^{n_i \times n_i}, i = 1, \dots, s$ denotes the one-dimensional DFT, $\otimes$ denotes the Kronecker product, $m = ln$, and $\text{ddiag}(v)$ denotes a diagonal matrix whose diagonal entries are v .

Numerical Experiments

Phase Retrieval



Numerical Experiments

Conclusion

- Smooth optimization on Riemannian manifold
- Limited-memory symmetric rank-one trust-region subproblem
- Intrinsic representation of tangent vectors
- Global convergence analysis
- Numerical experiments

Thank you

Thank you!

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